Dr. Dobb's Journal

For the Experienced in Microcomputing

A Public Key
Cryptography
System, Part II

Optimizing
Strings in C

BASICFMT

Dr. Dobb's Journal

For the Experienced in Microcomputing

CONTENTS

ARTICLES

DEPARTMENTS

EDITORIAL

Dr. Dobb's Journal began as a forum for sharing information and ideas about pro-
gramming and computers. It continues to be a place to present new languages, utilities,
tools, applications, algorithms, discoveries, and techniques to the microcomputing
community. That community may then use, comment upon, improve, and expand these
items. Many of the things explored within our pages have had major influence in the
industry. *DDJ* readers have always been at the forefront of the microcomputer industry,
providing a wealth of expertise to fuel discussion. Our authors have come primarily from
within the readership, and it is this reader involvement that has sustained and guided
the *Journal* through the years.

The material currently discussed in our pages is far more complex than it was eight
years ago, and this has forced some adjustments. We continue to print complete
listings of useful and innovative software, and to explore developments in areas such as
languages and operating systems. Maintaining the community discussion of these larger,
more complex ideas, however, has become increasingly cumbersome. We have periodic-
ally instituted reader-driven columns to provide more intimate exchange in areas of
particular reader interest, the latest addition dealing with C and Unix. Most recently,
we formed a board of referees (again, from within the readership) to help ensure that
the articles published are as appropriate and accurate as possible.

From our homespun beginnings with tiny BASIC and tiny computers, we have
grown with the microcomputing industry and community. Cosmetic changes have
occurred to improve appearance and marketability. Fiscal considerations have prompted
us to accept advertising, though we still work to keep the percentage down and to main-
tain our objectivity. And it was advertising revenue that helped us to begin paying
authors within the last few months. Each move has been made because we felt it would
allow us to better perform our function without compromising our traditions.

Now we have had a new opportunity presented to us – one which can open the
door for many of the projects that we have longed to pursue, and one which will provide
added resources with which to continue improving the quality of material we bring to
you. This new opportunity will allow the tradition of *DDJ* to continue and allow
exploration of new vistas.

We are joining forces with M&T Publishing, Inc., a company dedicated to quality
and integrity, a company which likes what *DDJ* has been doing and would like to see it
done even better. Laird Foshay, the president of M&T Publishing, was part of the *DDJ*
staff in the past, and has always been interested in our objectives. People's Computer
Company will still own *DDJ* while M&T will be licensed to publish the magazine. PCC's
Executive Director will chair an advisory board set up to ensure that *DDJ* maintains its
focus and direction. The current staff, slightly expanded, will continue to produce the
magazine as usual.

Among the new faces will be Michael Swaine, who is joining the staff as editor-in-
chief. Many of you may recognize the name from his three-year tour at *InfoWorld* as a
senior editor. A long-time reader of *DDJ*, Mike is dedicated to the traditions of the
Journal. A professional journalist and computer enthusiast, Mike is interested in con-
tinuing to improve the quality and scope of the service we provide. We are excited to
have him on the team.

So what does this all mean? Well, it means that we will have a new address. It means
new projects (perhaps a bulletin board here at the *DDJ* office) and new opportunities.
It means better service to our subscribers and better communication links to our readers.
It means more quality, more information, more of what you buy *DDJ* for. We will con-
tinue to explore ways to facilitate the dialogue among the microcomputing community.
Don't look for radical changes in the *Journal;* we have no intention of diverting from
our long-chosen path. Expect instead a continuation of our gentle evolution.

Reynold Wiggins

THIS MONTH'S REFEREES

LETTERS

NBASIC, Postprocessed

Mr. Shammas' NBASIC program should not be confused with the NBASIC program marketed by the Computer Toolbox, Inc. – Ed.

Dear Sirs:

In reference to the article about NBASIC on page 24 of your 1/84 issue, I suspect the program has the following errors:

(1) Line 1220. If a variable has 'REM' in it and it appears on a line before a Remark, it will be changed to a single quote.

(2) Line 1120. If a filename is entered without a period in it, then line 1120 never gets executed and the output filename FOUT $ never gets defined.

I have several preprocessors and use them a lot. I think they should be promoted more in all computer journals.

Sincerely yours,
Peter Ansbacher
918 South Rustic Road
Columbia, MO 65201

Dear *DDJ*:

I read the article "NBASIC: A Structured Preprocessor for MBASIC" in the January 1984 issue. In running the program on a test sample of code I discovered a problem with the routine for converting remark statements to single quote marks. On page 27 the following lines are given for the conversion:

```
1210 FOR I=1 TO N
1220     P=INSTR (L$ (I), "REM")
1230     IF P < > 0 THEN L$ (I) =
         MID$ (L$ (I) ,1,P–1) +
         " ' " + MID$ (L$ (I) ,P+3)
1240 NEXT I
```

This will replace the first occurrence of the string "REM" in each line of text. However, the character may appear within quoted strings or as part of a variable name. It will cause execution problems if the variables and strings are changed. Mr. Shammas' program will change the test code in Listing One (at right) into the text in Listing Two (at right).

Making a few changes to his code will intelligently change the remark statements. The speed for the conversion is about the same as Mr. Shammas'. Delete lines 1210 and 1220 and add the lines in Listing Three (at right) to the program. Listing Four (at right) shows the

output from the corrected program. This should handle most programs written in MBASIC.

This code may be improved upon, however it does work. Thank you for such a fine journal.

Sincerely,
Keith L. Terrill
603 Court Street
Syracuse, New York 13208

Listing One

```
10  REM
20  REM                 THIS IS A TEST FOR THE REM PROGRAM
30  REM
40  A=1 : B=2 : C=3 : REM THIS IS ANOTHER REM LINE
50  INPUT "THIS REM SHOULD NOT BE CHANGED";A
60  PRINT "THIS REM SHOULD REM BE CHANGED" :
        REM BUT THIS ONE SHOULD
70  DATA ONE, 1, TWO, 2, REM, REM1, REM2
80  END
```

Listing Two

```
10  '
20  '                   THIS IS A TEST FOR THE REM PROGRAM
30  '
40  A=1 : B=2 : C=3 : ' THIS IS ANOTHER REM LINE
50  INPUT "THIS ' SHOULD NOT BE CHANGED";A
60  PRINT "THIS ' SHOULD REM BE CHANGED" :
        REM BUT THIS ONE SHOULD
70  DATA ONE, 1, TWO, 2, ' , REM1, REM2
80  END
```

Listing Three

```
1209  QT$=CHR$ (34)
1210  FOR I=1 TO N
1213      P=INSTR (L$ (I) , "REM") : IF P=0 THEN 1240
1216      P1=INSTR (L$ (I) , "DATA") : IF P1=0 THEN 1222
1219      IF P>P1 THEN 1240
1222      P2=1
1225      P1=INSTR (P2, L$ (I) ,QT$) : IF P1>P OR P1=0 THEN 1237
1228      P2=INSTR (P1+1,L$ (I),QTS) : IF P2=0 THEN 1240
1231      P2=P2+1 : IF P2<P THEN 1225
1234      P=INSTR (P2,L$ (I) ,"REM") : IF P=0 THEN 1240 ELSE 1225
1237      L$ (I)=MID$ (L$ (I) ,1,P–1) +" ' "+MID$ (L$ (I) ,P+3)
1240  NEXT I
```

Listing Four

```
10  '
20  '       THIS IS A TEST FOR THE REM PROGRAM
30  '
40  A=1 : B=2 : C=3 : ' THIS IS ANOTHER REM LINE
50  INPUT "THIS REM SHOULD NOT BE CHANGED";A
60  PRINT "THIS REM SHOULD REM BE CHANGED" :
        ' BUT THIS ONE SHOULD
70  DATA ONE, 1, TWO, 2, REM, REM1, REM2
80  END
```

Fussin' over Sal/80

Dear Doctor:

This is to express my heartfelt appreciation for the very thorough and objective review of our SAL/80 package in the February issue by Jim Kronman. I'm embarrassed to take *any* exceptions to what was written, *but . . .*

One thing I do feel obliged to remark on: Jim's example of code-density compares two printer drivers, one in naked assembler and one in SAL/80. The SAL/80 driver is 400 bytes larger than the naked assembler, which might seem to imply that SAL/80 has a run-time package of 400 bytes. This is *not* the case. SAL/80 has *no* run-time package (and *no* royalties for commercial distribution of code generated using SAL/80).

But then it might seem to imply that the code emitted by SAL/80 is 40% larger than that obtained by coding in naked assembler, which is even *worse!* I haven't seen the two drivers in question, but I am quite certain that there are differences between the two which would account for all except a maximum of about 100 bytes plus perhaps 5% of the remaining difference.

There are several UTILITIES in the packages which, all taken together, might contribute perhaps 100 bytes of excess code because of their generality. And the test/branch code can typically be squeezed by about 5% in the average case. The 5% results from the fact that jumps to jumps do occur when a forward branch is emitted in the last group of a loop-body (hence, three bytes may be squeezable), and that the HL register is saved in certain memory-reference comparisons (which gives two possibly squeezable bytes if HL was inactive at the time).

The reason I'm making such a fuss about this is that SAL/80 is the only structured assembly package I know of that actually does *optimize* the test/branch code! There is a huge decision-tree in the compiler which does all the tricks dear to the hearts of assembly-language programmers to avoid writing two conditional jumps to implement LE ($<=$) and GT ($>$) comparisions. On a branch-by-branch basis, SAL/80 emits the tightest code possible if the double-registers are to be preserved over the comparison.

Cordially,
Steve Newberry
Protools
24225 Summerhill Avenue
Los Altos, CA 94022

Ackermann Uncovered

Dear Editor,

In response to Paul Condon's question (*DDJ* No. 88, February 1984, p. 9) "Who is Ackermann and where did he publish a definition of this function?" the following:

Wilhelm Ackermann died about twenty years ago. A student of David Hilbert's, he was one of the giants in research in the Foundations of Mathematics movement of the early part of this century. The function now associated with his name first appeared in a research monograph entitled "Zum Hilbertschen Aufbau der reelen Zalen" ("On Hilbert's Reconstruction of the Real Numbers") in *Math. Annalen,* Vol. 99 (1928), pp. 118-133.

The function was used to disprove (by counter-example) a conjecture by Hilbert to the effect that the set of real numbers definable by infinite sequences of what are now called "primitive-recursive" functions were, in an intuitive sense, all the real numbers that exist. Part of the argument of supposing this to be true was another conjecture that all recursive functions are primitive-recursive. The Ackermann function is *not* a primitive-recursive function, because it can be shown to "grow faster" than any primitive-recursive function. So, by showing that there are recursive functions which are not primitive-recursive, Ackermann disproved the first conjecture. Later, in the '30s, it was shown by Alan Turing that there are real numbers which are not definable even in terms of the most broadly generalized form of recursive function.

The original formulation of Ackermann's was slightly more complex than that now used, which is:

$$Ack(0, n) = n+1$$
$$Ack(m+1, 0) = Ack(m, 1)$$
$$Ack(m+1, n+1) = Ack(m, Ack(m+1,n))$$

A fuller discussion will be found in Hans Hermes' *Enumerability, Decidability, Computability,* Academic Press (1965), p. 84, and in Rozsa Peter's *Recursive Functions,* 3rd Revised Edition, Academic Press (1967), p. 106.

Didactically,
Steve Newberry
Newberry Microsystems
24225 Summerhill Ave.
Los Altos, CA 94022

DDJ

DR. DOBB'S CLINIC

by D.E. Cortesi, Resident Intern

Super Directory Fix

A while back we told how David McLanahan found that the public-domain program SD (one of the extended DIR programs) failed under CP/M Plus. Ted Shapin, of Orange, CA, writes to say that "The SD program can easily be fixed to give the correct free space under CP/M 3.0. I fixed my copy with information from D. Sternlight. At the label TMOVE there is an instruction, LDAX D. A few lines further there is a second LDAX D. After this instruction insert the line

ANI 0FH ;DROP ARCHIVE FLAGS

Meanwhile, McLanahan had found another solution. A later version of SD, SD48, works just fine on CP/M Plus.

RAM-Drive Needed

Malcom Fordham, of Bowie, MD, has "an IBM PC on which I use PL/I-86 under CP/M-86. I have sufficient memory to run the compiler and linker from a RAM disk. However, so far I haven't been able to find the software to do this under CP/M-86. Do you know of such a program?"

We only know that Concurrent CP/M for the PC contains RAM-drive support out of the box. Does anyone know of RAM-drive support for plain old CP/M-86? Maybe one of the makers of add-on RAM cards supplies one; they usually supply the support code for PC-DOS.

The Great American Merchandiser

We grumbled loudly about the merchandising tactics used in our area to promote subscriptions to *PC* magazine: a come-on mailed in an envelope that emulated an official government mailing. Dave Sullivan of Palo Alto, CA, writes to say that he got a second one of those foolers. "The latest one was from *PC Tech Journal*. It appears that Ziff-Davis is the one exhibiting all the brass.

"But I have another case of gall and guts for your scrapbook," he continues. "Ziff-Davis publishes any number of magazines that touch on computers. One of the cheaper ones is called *Computers and Electronics*, and sells for $1.75 at the newsstand. The January 1984 issue has what appears to be the first of a new series of articles on assembly language programming on the PC. It looks promising, and I'll take all the help I can get to learn assembly language.

"But the last paragraph of the article reads as follows: 'If you want to pursue learning assembly language for the IBM PC, additional articles that build on this one can be read in upcoming issues of our sister publication, *PC* magazine.'

"It takes one's breath away."

What's So Hard About CP/M Plus?

We're very fond of CP/M Plus now that we have it working well. So fond that we are trying it out on other computers besides our workhorse S-100 system. Or we would be, if anybody could get it working.

There's a Tandy Model 4P, an adorable little machine, that is patiently gathering dust while it awaits the long-delayed arrival of its CP/M Plus support from Fort Worth. Tandy promised CP/M Plus when they announced the Model 4, many months back, and they still haven't delivered. They keep promising it Real Soon Now, and they keep not shipping.

Then consider our brother-in-law Bill, whom we had all primed to buy a Morrow MD-11 as his first computer, largely because it was supposed to come with CP/M Plus. Only it won't, it seems. At least the first units of the MD-11 will be shipped with plain old CP/M 2. The dealer claims "it's real hard to make CP/M Plus work with the hard disk." To which we are inclined to say, piffle, but then what do we know?

What goes on here? Bob Blum got CP/M Plus working; we got it working. David McLanahan has had it for months on his Micro Mate (a system that, he says, deserves much wider exposure than it gets). The pubs are clear and complete; the generic parts of the system are reliable; the payoff in extra function is large. So what's the holdup?

CP/M Plus, Apple, and ALS

Advanced Logic Systems is a company that managed to make CP/M Plus work quite early. We have been trying out an Apple IIe equipped with the ALS CP/M card and CP/M Plus, and it works. It isn't as nice as CP/M Plus on a *real* computer (that should bring some mail), but to a large extent that's the fault of the Apple disk drives, not the fault of ALS or the software. There just isn't a whole lot you can do with only 170K per drive.

ALS did skimp on their implementation in one area. They completely omitted implementing the CP/M Plus character

I/O table and its character I/O redirection features. We can sympathize, because character I/O in the Apple is a zoo; there are so many possible permutations of serial cards, parallel cards, ROM versions, and slot numbers that it's hard to blame them for throwing up their hands. However, they didn't help the situation when they omitted documenting character I/O in any way whatsoever. The ALS bios contains some kind of support for the AUX device, for instance, but the ALS manuals contain not even a hint as to what AUX connects to.

Since our plans for the IIe involve getting MODEM to run on it so we can exchange files between it and the big system, we had to learn something about Apple I/O under the ALS bios. The ALS hot-line seems to be overloaded; they returned one of our calls out of six. The one contact we managed to make with them didn't produce any useful technical information, but it did give us the pointer we needed.

The ALS support person pointed us to a new CP/M Plus User Group being formed around the ALS card and the Apple. The group is CP/PLUG; its president is Jim Scardelis; and you can reach the group at P. O. Box 295, Little Falls, NJ 07424. Dues are $25 a year (we think, our notes are a mess) and the first quarterly newsletter was in preparation in January.

Scardelis gave us some valuable technical hints, but told us that if we wanted more we should dial into a certain RCP/M system run by one of the club members. We did and found a gold mine of Apple/ALS material. The system is the Central New York Technical RCP/M system, Mark Howard SYSOP. It is up from 1700 to 0800 hours EST, on (315) 437-4890, and is an excellent source of technical data. Thanks to Scardelis and Howard, we got our IIe running MODEM712 in fine shape.

How It's Done

Oh, you wanted to know how you do serial I/O on the Apple? Thought you'd never ask. When you are running under the ALS CP/M card, bank 1 (the bank containing the transient program area) consists of RAM chips on the ALS card itself. Global RAM starts from E000h and is also physically on the ALS card. The original Apple address space is bank 0. That Apple address space, as seen by a Z80 program, is remapped. Where in a real Apple the various I/O cards are mem-

ory mapped beginning at C000h, the Z80 sees them at 6000h in bank 0. So what you have to do is to find out which memory-mapped addresses your Apple serial card is supposed to appear at (somewhere in C0xxh). Under CP/M Plus, it will appear at 60xxh in bank 0.

Now all you have to do is find a way to read and write those bytes of bank 0 from a program running in bank 1. Mark Howard has written a Resident System Extension (RSX) that will do that for you. If it has been installed, you will be able to call it with a bank-0 address and get back the value of, say, the status port of your serial card. You write a MODEM7 overlay to make these calls and, voila, you are on the air. Telescoped into that "voila," of course, are about thirteen hours of sweat and hair-pulling.

Our April Item

We are in receipt of a letter from one T. Heintern of Milpitas, CA. Despite his slightly stilted English (Heintern is an immigrant from Islandia), we think his story will intrigue you.

"Your readers might be interested in the activities of two of my co-workers, or really ex-co-workers. I am, and they were, employed by a well-known maker of magnetic media here in Silicon Valley.

"At any rate, my ex-colleagues recently left our beneficent company and struck out on their own. They were remarkably close-mouthed about what they meant to do, but everyone in our department was aware that they had been spending long lunch hours in the bar at the Red Lion Inn in San Jose, in the company of various folk who wore the kind of urban-western gear that replaces a venture capitalist's ivy-league suit when he aims to blend into the crowd 'out west.' I presume my peers found the money they needed, or at any rate the strong hint of it, because they tendered their resignations just before Comdex.

"I next saw them in Las Vegas, where they were holding some fairly lush parties in a private suite. It was also at Comdex that I began to pick up rumors of a new, even-smaller, floppy disk format that (everyone told me) was about to burst on the scene. 'Forget the three-inchers,' I heard one fellow tell another in the midst of the crowd at the Microsoft Windows demo. 'It's almost standardized, so it's dead, right? The big thing is gonna be the CPS . . . .' Here the crowd pressed me away from them. Had they said CPS? I wasn't sure. Later I heard similar references to a 'Colt.'

"Back home I was swapping rumors with the other folk who'd attended the convention, and they too had heard rumors of a new diskette — and that it was to be of domestic, not Japanese, manufacture. Best of all, one had picked up an indication of its size: it was to be of 3.8 cm diameter.

"Well, I was mulling all this over when it all fell into place in my mind: CPS, Colt, 3.8 cm — a Colt Police Special is a 38-calibre pistol! Just as 'Winchester' was the code name for the IBM 3030, 'Colt' was a code name for a 3.8 cm diskette. Three point eight centimeters comes to about one and a half inches, I thought. And the light dawned!

"There flashed before my mind's eye the picture of our production floor: The great sheets of mylar spinning off the drums and under the punches . . . the thousands of 3.8 cm mylar circles fluttering down in drifts from the 8-inch floppies-to-be . . . our janitors sweeping them up . . . our teams of diligent house-wives bundling them into tidy stacks to be piled in building-high columns in the insulation space of our great Eastern wall.

"I dashed outside and rapped on the wall — hollow! Our insulation, years of 3.8 cm floppy-punchings, was gone!

"I have, of course, notified corporate management, who are taking legal advice. However, I wish to warn your readers. Shun the new diskette format! After all, you've already paid for the doughnuts. Will you now buy the holes?" ▪DDJ

CP/M EXCHANGE

by Robert Blum

Over the past few months it has been increasingly difficult to obtain DRI's permission to publish the application notes for CP/M. Apparently a shift in personnel at DRI had slowed the approval process, and I misunderstood how freely I could use the notes for CP/M Plus. Nonetheless, DRI has given us permission to continue publishing the application notes for CP/M V2.2 and CP/M Plus. This includes the CP/M Plus BDOS patch 02, published in the February issue of *DDJ*, which dealt with the enhancement of the LRU buffering logic of CP/M Plus. It has now been approved without changes.

When I was told of DRI's decision to allow us to continue publishing the application notes I breathed a deep sigh of relief. From the onset of our problems I was concerned that DRI was beginning to develop an attitude, or attitude problem depending on how you look at it, of pulling back from directly supporting their products. Time and again I have seen the customer support function inadequately treated by sales organizations and otherwise excellent products given undeserving tarnished reputations. I am glad that we can continue to depend on DRI to directly support their products. Starting next month we will begin publishing more application notes for CP/M.

Compatibility

Manufacturers dream of being able to move from one product generation to the next without any severe software incompatibility. And in most cases they are almost completely successful. This issue must have been a painful one at DRI while forming a product strategy for CP/M Plus. If their product marketing department is subject to the type of pressures that I have seen in similar industries they were undoubtedly presented with a difficult decision.

DRI was obligated to maintain compatibility with CP/M V2.2. It would have been incomprehensible for them to even consider breaking the tie with the hundreds of thousands of V2.2 systems currently in use. Any attempt to entice the software developers, at this time, to convert their packages to a new operating system or just to make a few modifications would have been out of the question. Being the incumbent in a time when cries are heard from every street corner that this or that new machine or software package is better than the others is worth

more in advertising dollars than most companies can muster.

By having available an upward migration path the current user base would be able to easily purchase, or upgrade to, the latest version of the product. Even though not prevalent in the microcomputer industry today, we will soon find software companies demanding that customers maintain their software at current levels or pay the price of nonsupport. From the manufacturers' standpoint this is the least costly method of introducing a new product because one immediately has a qualified customer base to sell to. Support costs are reduced as well.

CP/M Plus's compatibility with older versions of CP/M is maintained at a system level. All the former BDOS calls were maintained, although in many cases extra data or status information is made available for the programmer's use. To expand the system many new calls were implemented, but again, their use is entirely up to the programmer. This is as far as DRI will go in saying that CP/M Plus is compatible with V2.2. In the System Guide they refer to CP/M Plus as simply being "upward-compatible."

In V2.2 it was common practice to "hook" into the BIOS table to enable a program to trap status codes and other information before the BDOS had a chance to act upon it. Under CP/M Plus this practice is frowned upon and is largely unnecessary due to the new BDOS error modes and the additional status information returned from each function. Programs that do make use of the old error trapping technique will continue to run without any problems provided that system alteration does not go beyond the console and list BIOS vector points. Since these were the most-used points of alteration, DRI specified that they remain available as before.

The main reason for this restriction is the address extracted from the BIOS

jump table may, and probably will not, point to a real memory address. Keep in mind that in virtual memory systems (a banked CP/M Plus system falls into this category) any memory address may or may not currently reside in the processor's physical address space. Depending on the hardware manufacturer the determination of whether a memory address is real or virtual at any moment is a combined effort of both hardware and software.

In the case of CP/M Plus, some logic sections and table structures are required to be in real (common) memory at all times due to performance and memory conflict considerations. The majority of the system, however, will be virtual. When an application program makes a call to the BDOS to request a system service, a check is made to determine whether the requested portion of the system is in real memory. If not, a portion of the application program is made virtual and the missing section of the operating system is switched to occupy the real memory address space. As complicated or confusing as this may seem, it is handled quite simply in the CP/M Plus software and should not be intimidating.

To better visualize the BIOS routines that are in virtual memory refer to Listing One (page 14). All the routines pointed to by an address constant marked by a double quote are in virtual memory. You may wish to also refer back to this column in the July and August 1983 issues. In those columns I talked more about virtual memory systems and how they are generally implemented on microcomputers and particularly under CP/M Plus.

Most of the programs I have tested under CP/M Plus ran without any problems. I have found a couple of very popular programs that refused to cooperate and caused some totally unexpected and very difficult-to-find system problems until I discontinued their use. On the other hand, some programs that I would have

+0	+1	+2	+3	+4		+n
MX	NC	C1	C2	C3	. . . .	CN

Figure

Function 10 Console Input Buffer

bet wouldn't work, do. The only suggestion that I can make to you is to ask your dealer or the manufacturer whether their programs will work under CP/M Plus. Listen to their answers very closely for key phrases such as "it should" or "I don't see why it wouldn't." If you are given an answer of this type pursue the subject further by asking whether a test was actually run and if so on what machine. I think you will find that, when backed to the wall, most of the manufacturers' representatives will admit that no direct experience is available. This is not to say that what you are looking to buy won't work; chances are it will. Just protect yourself by making sure that you can return the package if some problem does exist.

BDOS Function 10: Vastly Improved

The read console buffer function, BDOS function 10, has been vastly improved in CP/M Plus. It will now not only automatically input and edit a string from the console, but will also display the string buffer before going into edit mode.

The format of the input buffer pointed to by the DE register pair is unchanged from V2.2 (refer to Figure, page 12). The single byte, MX, is the maximum number of characters that can be placed into the buffer. And the single byte, NC, is the count of the number of characters that have been input to the buffer. The actual data that has been entered into the buffer follows the NC byte. During the data entry operation the full keyboard editing features of CP/M Plus can be used to adjust the input string.

If register pair DE is set to zero prior to calling the BDOS, function 10 assumes that an initialized input buffer is located at the current DMA address. This allows a program to put a string on the console for the user to edit. To initialize the input buffer, place the default string into the buffer following the NC byte terminated by a binary zero. When initializing the input buffer with a default string it is not necessary to set the NC byte to the number of characters in the buffer.

When a program calls the function 10 routine with an initialized buffer, function 10 operates as if the user had typed in the string. A search is made through the buffer for the binary zero terminator byte in order to count the number of active characters. This count is then placed into the NC field. The routine's attention is then tuned to allow more data to be entered from the console. At this point the user is in full keyboard editing mode and can enter more data or edit the default string. In any case hitting return will immediately send the string to the program.

As now programmed, function 10 could greatly assist in those programs that offer a default value to the user as the first choice. Refer to Listing Two (page 16) for an example of a program that makes full use of function 10's expanded capabilities.

■■J

CP/M Exchange (Text begins on page 12)
Listing One

```
                    ; bios jump vector.

                    ; all bios routines are invoked by calling these
                    ;       entry points.

C3 0000"  ?boot:    jp    boot      ; initial entry on cold start
C3 006C'  ?wboot:   jp    wboot     ; reentry on program exit, warm start

C3 0177'  ?const:   jp    const     ; return console input status
C3 0192'  ?conin:   jp    conin     ; return console input character
C3 00DA'  ?cono:    jp    conout    ; send console output character
C3 00E6'  ?list:    jp    list      ; send list output character
C3 00E0'  ?auxo:    jp    auxout    ; send auxiliary output character
C3 0198'  ?auxi:    jp    auxin     ; return auxiliary input character

C3 006E"  ?home:    jp    home      ; set disks to logical home
C3 003F"  ?sldsk:   jp    seldsk    ; select disk drive
C3 0071"  ?sttrk:   jp    settrk    ; set disk track
C3 0077"  ?stsec:   jp    setsec    ; set disk sector
C3 007D"  ?stdma:   jp    setdma    ; set disk i/o memory address
C3 0094"  ?read:    jp    read      ; read physical block(s)
C3 00AA"  ?write:   jp    write     ; write physical block(s)

C3 0112'  ?lists:   jp    listst    ; return list device status
C3 0089"  ?sctrn:   jp    sectrn    ; translate logical to physical sector
```

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CP/M Exchange   (Listing Continued, text begins on page 12)
Listing One

```
C3 0106'    ?conos:     jp      conost      ; return console output status
C3 017D'    ?auxis:     jp      auxist      ; return aux input status
C3 010C'    ?auxos:     jp      auxost      ; return aux output status
C3 00D2'    ?dvtbl:     jp      devtbl      ; return address of device def table
C3 0000*    ?devin:     jp      ?cinit      ; change baud rate of device

C3 00D6'    ?drtbl:     jp      getdrv      ; return address of disk drive table
C3 00CB"    ?mltio:     jp      multio      ; set multiple record count for disk i/o
C3 00CF"    ?flush:     jp      flush       ; flush bios maintained disk caching

C3 0000*    ?mov:       jp      ?move       ; block move memory to memory
C3 0000*    ?tim:       jp      ?time       ; signal time and date operation
C3 0250'    ?bnksl:     jp      bnksel      ; select bank for code execution and default dma
C3 0085"    ?stbnk:     jp      setbnk      ; select different bank for disk i/o dma operations
C3 0000*    ?xmov:      jp      ?xmove      ; set source and destination banks
```

End Listing One

Listing Two

```
            .z80

bdos        equ     5                   ;BDOS function call entry point

set_dma     equ     26                  ;set DMA address function code
read_buffer equ     10                  ;read/edit console buffer function code

            ld      sp,stack            ;set a local stack

            ld      c,set_dma           ;set DMA address to input buffer
            ld      de,buffer           ;*
            call    bdos                ;*

            ld      c,read_buffer       ;edit string in buffer
            ld      de,0                ;*
            call    bdos                ;*

            rst     0                   ;return to CP/M

buffer      equ     $
            defb    10                  ;maximum number of characters allowed
            defb    5                   ;number of active characters
            defb    'string',0          ;initial string - zero delimited
            defb    '     '             ;room for maximum input length

            defs    30
stack       equ     $

            end
```

End Listings

Optimizing Strings in C

The C programming language offers many advantages: flexibility, speed, and portability. Routines that could take months in assembler can be written and tested much more easily in C. It supports true recursion and produces assembler language source, which allows users to optimize portions of their code.

Optimization is the process of improving the efficiency of a program; i.e., getting more bang for your buck, whether the buck is time or memory. Basically, you try to make the program smaller, faster, or both, although you sometimes have to trade size for speed or vice versa. Optimization can apply to one particular program or to a group of programs. One problem with optimization is that it takes both time and effort.

The following optimization approaches have certain advantages and disadvantages:

Improving the compiler output by improving the compiler

Improving the compiler output by tinkering

Improving the individual program's algorithm

Rewriting entire routines in assembler

To improve the compiler output, you must have the source for the compiler and must know enough to improve the code generation logic. If you have those skills, please try it and publish the results. For the rest of us mere mortals, however, this is a foolhardy exercise that can bring disaster. Besides, it takes a lot of time.

C programmers have another alternative: to recognize a commonly used pattern of code that can be improved and to write a program that tinkers with the compiler's output to improve that part of the code. This approach is very similar to the optimization portion of the Small-C compiler.

Improving an individual program by rewriting all or part of it is a perfectly valid approach, except that the payback doesn't go any farther. The rest of your programs are still the same.

The final approach is write some commonly used routines in assembler. This

by Edward McDermott

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will speed things up and you can generalize this approach to all your programs by recompiling them. It improves both execution speed and size (to some extent). The problem is that these routines are no longer portable from machine to machine.

If you choose this last approach, its success depends on which functions you select for optimization. The ideal candidates would have the following characteristics:

Contains tight loops

Is used by almost every program

Is small enough to easily rework

C is an ideal language for editors, compilers, operating systems and so on; that was what C was designed for. Although every one of these applications involves a good deal of string handling, C has no string manipulation verbs. All string handling is done on a character-by-character basis, often by very small functions, which contain some very tight loops.

Another aspect to consider is any special commands of the CPU you are using. Can any functions built into the chip be taken advantage of by another means? The Z80 chip has a few commands that are unavailable on the 8080, specifically: LDIR, LDDR, LDI, LDD, CPIR, CPDR, CPI, and CPD. Each of these instructions is the equivalent of a number of 8080 instructions and, properly used, can be substituted for a number of instructions in C.

For example, the LDIR command takes the contents of the memory position pointed at by the HL register pair, transfers those contents to the memory position specified by the DE register pair, increments HL and DE, and decrements BC until the BC register pair is zero. Perhaps this makes it a little clearer:

while (−−BC) *DE++ = *HL++;

Does that seem familiar? It is, in fact, almost identical to the STRCPY function in *The C Programming Language* by Kernighan and Ritchie. The principal difference lies in the delimitation of the loop.

If you run on a Z80 chip, string manipulation is a natural choice. But the question that haunts all optimization efforts still remains: Is the payback worth the effort? To try and answer that question we chose to optimize the following four functions:

STRCPY − copy a string from one place to another

STRINIT − initialize a string to a specified value

STRLEN − find the length of a string

CMATCH − find a character within a string

Since the Small-C compiler produces assembler code, our first approach was to review its output. In small loops the compiler produces code that does not optimally use the registers, so we completely rewrote the functions; most of the speed gain can be attributed to this rewrite. This optimization should apply to 8080-based machines as well.

The accompanying Listing (page 20) gives the new routines in Z80 code. We chose Z80 both to take advantage of its special OP codes and to emphasize the major disadvantage of this type of optimization: the loss of portability. Readers, however, can adapt these routines to the 8080 environment fairly easily, for use with C or other languages.

Usually you judge the success of an optimization on the following criteria:

Is it faster and if so how much?

Is it smaller and if so how much?

Does it suggest other methods of improvement?

Was it worth the time and effort to do?

Faster

We acquired the timings here by executing each command 10,000 times on a string of 30 characters. The first set of timings is for the driver program, without any string manipulation at all, to see the impact of the code used to handle the test. The second set is using C routines, compiled under the Small-C compiler, version 2. The third set of timings is for the same program but using the assembler routines. The timings are as follows:

Null test	3 sec.
Small C (v2)	195 sec.
Assembler	16 sec.

We also compared the CMATCH routine to the rest, since this routine is significantly longer and more involved in both languages:

Null test	1 sec.
Small C (v2)	90 sec.
Assembler	6 sec.

If you remove the constant load of the code to drive the test program, you

find the new routines are 14.8 times faster. Skeptics can ignore the impact of the loops and still end up with new routines that are 12.2 times faster.

Smaller

The memory saving were illustrated by the following:

Language	Size
Small C(v2)	254 bytes
Assembler	116 bytes

The assembler code for the four routines was less than half the size of the comparable C code. However, saving 138 bytes is not going to make anyone jump for joy. With fine tuning, this figure could be improved, but once again it is a matter of cost versus gain.

Other Improvements

Our first thought was to optimize other small routines for handling strings. But part of the savings had come from using all three register pairs and keeping the values in the registers throughout the loop. Where the number of parameters exceeds three (one for each of BC, DE, and HL register pairs), the potential savings drop and the effort begins to rise.

Another improvement would be using these optimized routines as heavily as possible. That has some major implications for the style of C code a programmer uses. You would have to build up your functions out of smaller optimized functions instead of writing complex character-by-character processes.

Some other potential candidates for optimization are as follows:

Convert to upper case and convert to lower case

Find string

Compare string

Extract part of a string

All string manipulation functions should not necessarily be optimized. The gain in speed through improving screen and printer routines is hidden by the limited transmission speed to these devices for most small systems; you won't see any improvement. Most likely your disk I/O routines are already in assembler to interface the language to your particular operating system.

By the way, the STRCPY and STRINIT functions were designed to return the address of the resulting string. This may seem an unnecessary complication, but it does allow the programmer to nest functions within each other when writing the C code.

Time and Effort

Whether the optimization was worth the time and effort depends on how you value your time and on your requirements for speed and/or memory. My personal expectation for speed improvement was somewhere in the order of five times faster. The improvement, however, was between 12 and 15 times faster. Since creating these routines, testing them, correcting them, and correcting them again resulted in more speed than I expected, the optimization proved successful. You have to decide if it's worth your time to type these routines in and use them in your code.

The improvement you can achieve by converting small routines into assembler is significant, if not startling. Leaving the original C routine as a com- ment in the assembler code lets you combine the best of both worlds: the most efficient routine and the original portable routine. Any small, tight looping function that is concerned with three or fewer variables and is commonly used is an ideal candidate. However, to get the most out of such an optimization, you must go on to build more complex functions out of these smaller ones. **DDJ**

(Listing begins on page 20)

Editor's note on some enhancements: Any comparison of the accumulator to 0 can be replaced with an OR A command, and any initialization of the accumulator can be replaced with an XOR A command. This improves speed and reduces memory requirements.

Some places where one could make the code tighter are perhaps less obvious. In lines 50–54 of the listing, the following command sequence could replace the existing ones:

```
LD    A,B
OR    B
JR    Z,STRIN2
```

Likewise, lines 97–101 could be re-placed with

```
LD    A,B
OR    C
JR    Z,CMA1
```

and lines 103–106 could be replaced with

```
LD    A,B
OR    C
JR    NZ,CMAX1
```

In lines 104 and 106, the author used JP instructions instead of JRs. While JR is probably the better instruction, JP is safer in the case of the author's assembler and if one is uncertain about the distance to the branch location. In this case, JR should work quite well.

Optimizing Strings in C (Text begins on page 18)

```
 1: #asm
 2: ;/*
 3: .Z80
 4:
 5: ;strcpy(s,t) char *s, *t;
 6: ;          { int ret; ret = *s; while (*s++ = *t++); return(ret)  ; }
 7:
 8: STRCPY::
 9:          POP     BC                  ; pop return address
10:          POP     HL                  ; pop address of t
11:          POP     DE                  ; pop address of s
12:          PUSH    DE                  ; restore stack
13:          PUSH    HL
14:          PUSH    BC
15:          PUSH    DE                  ; save address of start of s
16:
17: STRLP:   LD      A,(HL)              ; transfer loop
18:          CP      00                  ; test for zero in t
19:          JR      Z,STREXT            ; if so exit
20:          LDI                         ; move incrementing
21:                                      ; s++ = t++
22:          JR      STRLP               ; continue loop
23:
24: STREXT:  LD      A,00                ; zero final bite in s
25:          LD      (DE),A
26:          POP     HL                  ; return original address of s
27:          RET
28:
29: ;strinit(s,c,i ) char *s; char c ; int i ;
30: ;          { int ret; ret = *s;
31: ;          while (i-- ) *s++=c;
32: ;          return(ret) ; }
33:
34: STRINIT::
35:          POP     IX                  ; pop return address
36:          POP     BC                  ; pop i (len for init)
37:          POP     DE                  ; pop c  ( init character
38:          POP     HL                  ; POP ADDRESS OF STRING
39:          PUSH    HL                  ; restore stack
40:          PUSH    DE
41:          PUSH    BC
42:          PUSH    IX
43:          PUSH    HL                  ; save address of start of s
44:          LD      A,E                 ; LOAD A  with c character
45:          LD      (HL),A              ; store character in begin of s
46:          POP     DE                  ; SET DE to begin of s
47:          PUSH    DE
48:          INC     DE                  ; point DE to next bye
49:          DEC     BC                  ; REDUCE FOR CHAR TAKEN
50:          LD      A,00
51:          CP      C                   ; MAKE SURE BC > 00
```

Optimizing Strings in C

```
52:             JR      NZ,STRIN1
53:             CP      B
54:             JR      Z,STRIN2
55:
56: STRIN1: LDIR                            ; propogate it through string
57: STRIN2: POP     HL                      ; get address of s to return
58:             RET
59: ;/*   strlen --- return length of string s (page 98) */
60: ;strlen(s) char *s;
61: ;       < int p ;  p =s;
62: ;       while (*s) ++s  ;
63: ;       return (s-p);>
64:
65: STRLEN::
66:             POP     BC                  ; pop return address
67:             POP     HL                  ; pop address of s
68:             PUSH    HL                  ; restore stack
69:             PUSH    BC
70:             PUSH    HL                  ; save address of start of s
71:             LD      BC,0FFFFH           ; BC IS BYTE COUNT (decremented)
72:             LD      A,00                ; searching for a  00
73:             CPIR                        ; equivalent of while (HL++)
74:             POP     DE                  ; DE = s
75:             SBC     HL,DE               ; HL = HL - DE
76:             DEC     HL                  ; correction for count last char
77:             RET                         ; RETURN (HL);
78:
79: ;cmatch(s,p,i )                                    /* find first p   in str s */
80: ;       char s[] ; int i, p ;
81: ;       < if (i > strlen(s)) return(0);
82: ;       while (s[i] != 0 )
83: ;               < if (p == s[i++]) return  (i); >
84: ;       return 0; >
85:
86: CMATCH::
87:             POP     IX                  ; pop return address
88:             POP     BC                  ; pop i  (indent for init)
89:             POP     DE                  ; pop p  (search character
90:             POP     HL                  ; POP ADDRESS OF STRING s
91:             PUSH    HL                  ; restore stack
92:             PUSH    DE
93:             PUSH    BC
94:             PUSH    IX
95:
96:             PUSH    BC                  ; save BC for count of bytes
97:             LD      A,0                 ; LOAD A  with 0
98:             CP      B                   ; IF BC == 0   THEN GO TO CMA1
99:             JR      NZ,CMA0             ;
100:            CP      C
101:            JR      Z,CMA1
102: CMA0:      CPIR                        ; check 00  before end of BC
103:            CP      B                   ; IF (BC != 0 ) return 0
104:            JP      NZ,CMAX1
105:            CP      C
```

```
106:           JP      NZ,CMAX1
107:                                        ; else continue
108: CMA1:     POP     BC                   ; restor BC = i for offest count
109:                                        ; i is within  the string ?
110: CMA2:     INC     BC                   ; count bytes
111:           LD      A,(HL)
112:           CP      00                   ; end of string?
113:           JR      Z,CMAX
114:           CP      E                    ; check search character
115:           JR      Z,CMAE               ; if (HL = p) contin
116:           INC     HL                   ; HL ++
117:           JR      CMA2
118:
119: CMAE:     LD      H,B                  ; return (HL = BC)
120:           LD      L,C
121:           RET
122:
123:           ; i is beyond the string end
124: CMAX1:    POP     BC                   ; restore stack
125: CMAX:     LD      HL,00H               ;RETURN (NULL);
126:           RET
127: #endasm
```

End Listing

Expert Systems and the Weather

S omebody famous once said: "This is gonna get me in trouble." Since I am about to discuss a program that is capable of predicting the weather (sometimes), I figure I'll get in trouble, but here goes.

I chose predicting the weather as the subject of an early effort at knowledge engineering using EXPERT-2, a tiny fifth-generation type language.[1] My friends in the business of weather prediction think this is somewhat amusing, but the results are quite interesting and possibly useful. I deliberately kept the knowledge engineering exercise simple since I intended to run this program in a 48K Apple II, using a Forth system with EXPERT-2 on top and the application program, the weather predictor.

Expert systems are computer programs that are specially designed to perform inferences: to prove things. EXPERT-2 is a program that allows individuals to write task programs much like BASIC lets users write programs. It is a high-level language written for performing logic inferences using rules written into the user's program. These rules may be used to solve important problems, like predicting the stock market . . . or the weather.

EXPERT-2 allows two methods for encoding knowledge and information into the task program — in this case, the weather predictor:

IF-THEN statements

Analytic subroutines

IF-THEN statements, or production rules, allow the user to direct the inference process by entering into the computer specific statements that say, "if something is true, then something else is true." The EXPERT-2 syntax is an extension of the IF-THEN syntax used to develop an animals game.[2,3] The extensions include allowing statements to be entered in a negative context (IFNOT, ANDNOT) and calling analytic subroutines.

Analytic subroutines were added to give the programmer access to the full power of the underlying Forth system. If EXPERT-2 had been written in, say,

Pascal, then the analytic subroutine calls (IFRUN, IFNOTRUN, ANDRUN, ANDNOTRUN) would permit access to that environment. As it is, EXPERT-2 was written in Forth, primarily because that's the program environment in which I work best. Users of EXPERT-2, however, need not know or understand Forth; if analytic subroutines are not needed, no Forth coding is required.

The weather predictor presented here does use analytic subroutines. These routines are used in two ways:

To prompt the user to input data

To process data and return truth values to the rules

The IF-THEN rules call the appropriate analytic subroutines when answers about weather data are needed.

The IF-THEN rules contain knowledge about the weather. This knowledge is encoded as antecedents (e.g., IF barometric pressure is falling rapidly) and consequents (e.g., THEN the weather is turning bad). Consequents follow appropriate antecedents. Designing rules to correctly encode knowledge is called knowledge engineering.

Knowledge Engineering

Someone once asked me if EXPERT-2 ever adds to its own knowledge base or makes original statements. As it turns out, it often makes original statements – especially during system debugging. Early versions of EXPERT-2, in fact, were self-modifying (not by design!). Test runs were incredibly mystifying. Knowledge engineering, as I use the term here, assumes one has available a working inference machine: a program that permits inferences to be performed on a set of rules. EXPERT-2 is one of the early microcomputer-based inference machines available; I expect there eventually will be a slew of them.

Knowledge engineering is the process of defining a problem to solve (something for the expert to do, especially something useful), encoding into a program the expertise required to solve the problem, testing, and finally validating the expert program.

To encode knowledge, one needs an expert whose brains (expertise) are accessible for encoding. As it turns out, that's not a generally available commodity; experts may be around, but it takes charm, wit, and experience to get at their brains. Trying to determine what tools an expert

brings to bear on a problem often gets you a textbook set of answers. Textbook answers are not necessarily what you need to solve a problem. Factors including years of experience, intuition, and a host of other issues must be explored before a knowledge engineer fully exhausts all the elements needed to solve a particular problem.

For the weather predictor, I chose to explore only books and magazines, partly because I didn't want to call in any chits I may have with the weather service and partly because I discovered some interesting prospects in the literature for home-brew weather prediction that could be translated to the EXPERT-2 system. A set of programs have been presented for computers that run BASIC.[4] These programs ask certain questions and issue weather predictions based on the answers. Because the algorithm used in the programs relates well to the general literature on weather prediction, I selected it as a representative "expert" set of rules.

For additions and changes to the rules derived from the magazine article I drew on the Heathkit Weather Training Manual, which has a section on weather forecasting.[5] Selected rules on cloud motion, for example, were used to enhance the rule base.

Analytic subroutines are then used to acquire weather data from the user and to process that data as the inference machine needs answers.

Weather patterns are quite sensitive to local terrain influences. Therefore, the responses the weather predictor offers may not be appropriate to, say, a canyon region where you live. It seems to work well where I live. I leave it as an exercise for the reader to validate the weather predictor program in his or her own area.

System Operation

Rules are typed into a disk file using the editor available with the underlying Forth system. EXPERT-2 is then loaded and compiled on top of Forth as a task. Finally, the user task (in this case, the weather predictor) is loaded on top of EXPERT-2. First Forth compiles the analytic subroutines and any variables or constants. EXPERT-2 then compiles the IF-THEN rules once the word RULES is encountered. Rule compiling ends with the word DONE.

Forth, in its usual fashion, tells you that compiling has ended by printing OK on the screen. At this time it should be

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safe to type RUN and start the forecasting program. RUN initializes the data base by asking appropriate questions then calls EXPERT-2 to solve the problems. The problem is to prove one of four possible hypotheses, identified in the rule base by THENHYP. The four hypotheses are:

Weather is OK

Weather is improving

Weather is turning bad

Insufficient data to perform a forecast

EXPERT-2 takes the first hypothesis and begins the task of proving it. The routine that performs this task is DIAGNOSE, which is called by RUN. DIAGNOSE selects the hypothesis and passes it deeper into the EXPERT-2 program. EXPERT-2 collects all the rules that might support a proof.

If the hypothesis is "weather is OK," then all rules that have as one of their consequent fields "weather is OK" are tagged and drawn into the inference process. One such rule is:

IFRUN BP>30.2

ANDRUN BP-SLOW-FALL

ANDRUN WDIR

BECAUSE WINDS FROM SW, W, OR NW

THENHYP WEATHER OK

ANDTHEN FAIR AND WARM NEXT 48 HRS

One of the consequent fields (in this case, the first consequent) is the same as our hypothesis; in fact, this rule happens to be the one selected to define an hypothesis with THENHYP. EXPERT-2 finds this rule and tries to prove it.

The inference machine takes each of the antecedents, one at a time, and tests them. If all pass with a truth value of TRUE, the consequents are taken as true, and the proof is complete. If any antecedent fails with a truth value of FALSE, the rest of that rule is not tested. Anything that is "learned" along the way, however, is saved as a fact. If the system asks you which way the clouds are moving, for example, your answer becomes a fact so that EXPERT-2 does not have to ask you that question again.

If EXPERT-2 proves a rule, DIAGNOSE tells you about it, and the program ends with I CONCLUDE . . . . If, on the other hand, a given hypothesis cannot be proven, then the next available hypothesis is selected, and a new proof is started. The new proof has available to it everything that the system learned during all previous proof attempts (all attempts, that is, since RUN was typed). If no proof is available on any hypothesis, EXPERT-2 typically states, CANNOT PROVE ANYTHING.

The Weather Predictor

You may wonder how such a simple-looking program can do any useful weather prediction; most people ask about this. As it turns out, the program concerns itself only with large-scale frontal movements. These fronts are best visualized (but not correctly described) as large bubbles of air hundreds of miles across, sliding around over the ground, bumping into each other, and generally moving from the west toward the east. Some of these bubbles are huge low-pressure masses of air, and some are large high-pressure masses of air.

When a low-pressure mass of air approaches your area, the barometer will spot the frontal motion and indicate that the barometric pressure is falling. As that air mass leaves your area, the barometer will begin to rise. Thus, by using readings of barometric pressure and asking questions about which way the winds are blowing, the weather predictor attempts to determine which of its "known" frontal-air-mass patterns best describes the conditions present in the atmosphere. If it finds a pattern that fits, it issues a prediction. The patterns it looks for, however, must be calibrated for your area.

As a simple exercise in knowledge engineering, the weather predictor demonstrates nearly the full capabilities of EXPERT-2. The program uses information on barometric pressure, surface wind direction, and upper air wind direction (as evidenced in upper cloud motion).

The rules begin with WALL, a word that lets you clear this rule base from memory if you want to load another program (by typing FORGET WALL). Following WALL, all analytic subroutines are loaded, written in Forth. This version (Listing One, page 27) is written in an extended Forth-79 dialect and based on the 40-character line width of a standard Apple II. Readers who type this in on other systems are encouraged to explore the enormous advantages of larger line widths; some awkward statements often are created when one is constrained to a narrow screen.

Following the analytic subroutines, RULES starts EXPERT-2's rule compiler. This invokes the IF-THEN syntax that the listing illustrates.

This weather predictor tests for 19 different possible combinations of weather conditions, using 19 basic weather prediction rules. If the requested weather information you type in fits within one of the 19 possible data combinations, a prediction will be issued. If the data falls outside all possible tests, the system will issue the statement INSUFFICIENT DATA FOR A FORECAST.

Since EXPERT-2 would rather say CANNOT PROVE ANYTHING, the last rule in the set stops that statement and issues a custom proof-collapse statement. That way, if EXPERT-2 is unable to prove any of the three main hypotheses, it is offered the last hypothesis as a consolation prize. This allows you to end an unsuccessful session with any statement you wish.

Three sets of rules are available for checking on upper level clouds. These rules, which support some of the 19 main rules by asking special questions about such clouds, underscore one of the weaknesses of EXPERT-2: its inability to handle an "I don't know" answer. That answer, in fact, does not exist in the syntax. An astute knowledge engineer, however, is still burdened with the likely event that this particular answer will be needed. Weather predictor handles the "don't know" case by allowing you to answer N for no if either you cannot see the upper clouds (there may not be any clouds to see) or you haven't gone outside to look as instructed. If you answer no, the system assumes the cloud motion supports whatever rule it happens to be trying to prove at the time.

This derives from the second rule in each of the three rule groups. The second rule, in effect, says: After dealing with the instructions (MESSAGE1), IFNOT you can see the clouds THEN clouds indicate . . . . One of these exists for each of the three possible conditions. A true condition is forced if you answer no. I am quite sure I'm gonna get in trouble for this.

A typical terminal session with this weather predictor follows:

1. System starts by asking barometric pressure; you enter 30.2.
2. System next asks what the barometric pressure is doing; you enter the code for falling slowly.
3. System asks for the direction from which the wind blows; you enter the code for south.
4. The terminal response is:
   NOW, WE CAN PREDICT THE WEATHER.
   I MIGHT BE ASKING A FEW MORE QUESTIONS.
   I DEDUCE
       RAIN WITHIN 24 HRS
   I DEDUCE
       WEATHER TURNING BAD
   I CONCLUDE
       WEATHER TURNING BAD
   OK (Forth's usual prompt)

No further questions were asked in that session. This particular session included data that triggered (successfully) the following rule:

IFRUN BP>30.1
    (30.2)
ANDRUN BP-SLOW-FALL
    (pressure falling slowly)
ANDRUN SDIR
    (southerly direction winds)
BECAUSE WINDS FROM SE OR S
THEN RAIN WITHIN 24 HRS
ANDTHEN WEATHER TURNING BAD

This rule illustrates a number of the design methods that were used in this knowledge engineering exercise. First, the final consequent, "weather turning bad," helps direct the proof search. If you read the rules, especially the consequent fields, you will see that a search could be conducted on dozens of different possible

consequents: "rain within 24 hrs," for example. If I had set as many hypotheses as there are different consequents in this program, the hypothesis field would have been enormous. Instead I chose three primary consequents as "search pruners" to limit the search requirement. All possible rules that have a particular hypothesis are collected at the same time for proof. The first rule to pass the test terminates all further tests. Other than that, "weather turning bad" is a pretty useless consequent.

A BECAUSE field is included in this rule to allow the programmer to tack explanations in where the code might otherwise be incomprehensible. For example, if the name chosen for an analytic subroutine that is called to check wind direction is not particularly descriptive, a BECAUSE statement is added. If the system asks a question in the form IS THIS TRUE?, you can answer Y for yes, N for no, or W for

"Why did you ask me that question?" It allows you to get an explanation, which is where the BECAUSE clause is handy.

The BECAUSE clause is also handy when passing around source code for programs. One of the advantages of the EXPERT-2 style program is that it is largely self-commenting. Where it is not self-commenting, you can add BECAUSE clauses.

This EXPERT-2 program's heavy use of analytic subroutines reduces the number of questions the program might be inclined to ask at the terminal. By contrast, the animals game mentioned earlier uses absolutely no analytic subroutines: whatever it learns in the proof process is gained from the terminal. Expert programs are much more interesting if they ask something at the terminal.

Going Beyond

EXPERT-2 is a high-level language for experimenters in early fifth-generation machines. It is made available in source code form, loadable on a variety of Forth systems. Using the source, this weather predictor program, and a lot of sweat equity, I expect some pretty strong weather predictor programs will emerge.

Weather prediction is largely a data processing and pattern recognition effort, one well suited to an expert system. By adding some simplified models related to atmospheric physics and coupling the program (through analytic subroutine functions) to an automatic weather station, you could perhaps create a new program for predicting weather hazards for farmers or sporting events.

Weather forecasters rate each other by a value they call "skill." Flipping a coin usually gets a skill of 0.5 or 50% right. It will be interesting to see what skill this program and its derivatives achieve. Since local weather prediction is sensitive to local geographic terrain effects, this program will need some fiddling to raise its skill, perhaps even to the 50% level.

Some thought might go into letting the program run on a pseudo-real-time basis, keeping track of its own skill. It also might be written to fiddle with its own inferences in an attempt to raise its skill: a self-learning system. You will have to make sure, however, that it does not learn how to fiddle with its skill algorithm.

References

[1] *EXPERT-2 - A Knowledge-Based Inference Program*. Available as a document or document with source code on disk from Mountain View Press, Inc. (P. O. Box 4656, Mt. View, CA 94040), Miller Microcomputer Services (61 Lake Shore Rd., Natick, MA 01760), or Helion, Inc., Box 445, Brownsville, CA 95919.

[2] Richard Duda and John Gaschnig. "Knowledge-Based Expert Systems Come of Age." *Byte*. September 1981.

[3] P. H. Winston and B. K. P. Horn. *LISP*. Addison-Wesley.

[4] George W. Miller. "Weather Forecaster." *COMPUTE!* August 1983. See also Tom Fox, *Unique Electronic Weather Projects*. Howard Sams & Co., 1978.

[5] *Weather Training Manual*. Benton Harbor, MI: Heath Company.

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Weather Predictor (Text begins on page 24)

```
Screen 1
( WEATHER PREDICTOR - EXPERT-2        JP)
: WALL ; FORTH
VARIABLE MESS1            ( toggle flag )
2VARIABLE BP        ( barometric press)
VARIABLE DIR 2 ALLOT        ( dir code )
VARIABLE BP?      ( pressure trend code )

: RUN 0 MESS1 ! DIR 4 ERASE HOME CR ."
     NIMBLE WEATHER PROGRAM" CR CR   ."
THIS PROGRAM ATTEMPTS TO PROVE ONE OF
THE FOLLOWING : " CR                 ."
   THE WEATHER IS OK" CR             ."
   THE WEATHER IS IMPROVING" CR      ."
   THE WEATHER IS DETERIORATING" CR  ."
NOT ALL WEATHER CASES ARE INCLUDED IN
THE FORECAST ALGORITHM." CR CR       ."
BEFORE WE ATTEMPT A FORECAST,
I NEED SOME DATA." CR CR             ."
WHAT IS THE BAROMETRIC PRESSURE READING?
< EXAMPLE - TYPE 30.2 > " INPUTD# CR CR
   BP 2!
-->
   CONTINUED ON NEXT SCREEN


Screen 2
( WEATHER PREDICTOR - EXPERT-2       JP)
                                     ."
SELECT FROM THE FOLLOWING : " CR     ."
   1 - BP IS STEADY" CR              ."
   2 - BP IS RISING SLOWLY" CR       ."
   3 - BP IS RISING RAPIDLY" CR      ."
   4 - BP IS FALLING SLOWLY" CR      ."
   5 - BP IS FALLING RAPIDLY " INPUT#
```

```
     BP? ! CR                          ."
THE WIND IS BLOWING FROM WHICH OF THE
FOLLOWING DIRECTIONS ? " CR           ."
< N, NE, E, SE, S SW, W, NW > "
   DIR 3 EXPECT CR CR                 ."
NOW WE CAN PREDICT THE WEATHER." CR ."
I MIGHT BE ASKING A FEW MORE QUESTIONS"
   DIAGNOSE ( start EXPERT-2 ) ;
-->


Screen 3
( WEATHER PREDICTOR - EXPERT-2       JP)
( analytical subroutines follow )
: BP>30.2 ( -- TF )
   BP 2@ 30.2 D> ;

: BP>30.1 ( -- TF )
   BP 2@ 30.1 D> ;

: BP<30.1 ( -- TF )
   BP 2@ 30.1 D< ;

: BP>29.9 ( -- TF )
   BP 2@ 2DUP 29.8 D< >R 29.8 D= R> OR ;

HEX ( now set-up direction tests )
: WDIR ( -- TF )
   DIR @ 5753 ( SW ) OVER = SWAP 57 ( W )
   OVER = SWAP 574E ( NW ) = OR OR ;
: SDIR ( -- TF )
   DIR @ 4553 ( SE ) OVER = SWAP 53 ( S )
   = OR ;
-->
```

(Continued on next page)

Weather Predictor (Listing Continued, text begins on page 24)

```
Screen 4
( WEATHER PREDICTOR - EXPERT-2        JP)
: NEDIR DIR @ 454E ( NE ) OVER = SWAP
  45 ( E ) = OR ;
: NDIR NEDIR DIR @ 4E ( N ) = OR ;
: EDIR NEDIR DIR @ 4553 ( SE ) = OR ;
: SEDIR SDIR DIR @ 45 ( E ) = OR ;
: SWDIR DIR @ 53 ( S ) OVER = SWAP
  5753 ( SW ) = OR ;
: BP-SLOW-FALL BP? @ 4 = ;
: BP-RAPID-FALL BP? @ 5 = ;
: BP-SLOW-RISE BP? @ 2 = ;
: BP-RAPID-RISE BP? @ 3 = ;
: BP-STEADY BP? @ 1 = ;
DECIMAL
-->


Screen 5
( WEATHER PREDICTOR - EXPERT-2        JP)

: MESSAGE1 MESS1 @ NOT
  IF CR                                 ."
FOR THE NEXT QUESTION, YOU SHOULD STAND
OUTSIDE WITH YOUR BACK TO THE SURFACE
WIND." CR                               ."
YOU NOW MUST OBSERVE THE DIRECTION UPPER
LEVEL CLOUDS ARE MOVING." CR            ."
YOU OBSERVE THEM TO BE MOVING FROM YOUR
RIGHT, FROM YOUR LEFT, OR IN A DIRECTION
PARALLEL TO THAT WHICH YOU ARE FACING."
        CR                              ."
USE THIS INFORMATION TO ANSWER THE
FOLLOWING QUESTION. IF YOU ARE UNABLE TO
SEE THE UPPER LEVEL CLOUDS, ANSWER NO TO
THE FOLLOWING QUESTION" CR
    THEN 1 MESS1 ! TRUE ;
-->


Screen 6
( WEATHER PREDICTOR - EXPERT-2        JP)
RULES   ( start the rule compiler )
( deteriorate test using clouds )
  IFRUN MESSAGE1
  ANDIF YOU CAN SEE UPPER LEVEL CLOUDS
  ANDIF CLOUDS MOVING FROM YOUR LEFT
THEN CLOUDS INDICATE BAD WEATHER
  IFRUN MESSAGE1
  ANDNOT YOU CAN SEE UPPER LEVEL CLOUDS
THEN CLOUDS INDICATE BAD WEATHER
( improving test using clouds )
  IFRUN MESSAGE1
  ANDIF YOU CAN SEE UPPER LEVEL CLOUDS
  ANDIF CLOUDS MOVING FROM YOUR RIGHT
THEN CLOUDS INDICATE WEATHER IMPROVING
  IFRUN MESSAGE1
  ANDNOT YOU CAN SEE UPPER LEVEL CLOUDS
THEN CLOUDS INDICATE WEATHER IMPROVING
-->


Screen 7
( WEATHER PREDICTOR - EXPERT-2        JP)
( steady test using clouds )
  IFRUN MESSAGE1
  ANDIF YOU CAN SEE UPPER LEVEL CLOUDS
  AND CLOUDS MOVE PARALLEL TO DIRECTION
THEN CLOUDS INDICATE STEADY WEATHER
  IFRUN MESSAGE1
  ANDNOT YOU CAN SEE UPPER LEVEL CLOUDS
THEN CLOUDS INDICATE STEADY WEATHER
-->


Screen 8
( WEATHER PREDICTOR - EXPERT-2        JP)
  IFRUN BP>30.2
  ANDRUN BP-SLOW-FALL
  ANDRUN WDIR
  BECAUSE WINDS FROM SW, W, OR NW
THENHYP WEATHER OK
ANDTHEN FAIR AND WARM NEXT 48 HRS
  IFRUN BP>30.2
  ANDRUN BP-STEADY
  ANDRUN WDIR
  BECAUSE WINDS FROM SW, W, OR NW
THEN CONTINUED FAIR
ANDTHEN LITTLE TEMPERATURE CHANGE
ANDTHEN WEATHER OK
  IFRUN BP>30.1
  ANDRUN BP-STEADY
  ANDRUN WDIR
  BECAUSE WINDS FROM SW, W, OR NW
THEN FAIR WEATHER
ANDTHEN LITTLE CHANGE NEXT 48 HRS
ANDTHEN WEATHER OK
-->


Screen 9
( WEATHER PREDICTOR - EXPERT-2        JP)
  IFRUN BP>30.1
  ANDRUN BP-SLOW-FALL
  ANDRUN NEDIR
  BECAUSE WINDS FROM NE OR E
  AND SEASON IS SUMMER
  BECAUSE NOT WINTER
  AND CLOUDS INDICATE STEADY WEATHER
THEN RAIN MAY NOT FALL
ANDTHEN STEADY FOR SEVERAL DAYS
ANDTHEN WEATHER OK
-->


Screen 10
( WEATHER PREDICTOR - EXPERT-2        JP)
  IFRUN BP<29.8
  ANDRUN BP-RAPID-RISE
THENHYP WEATHER IMPROVING
ANDTHEN CLEARING AND COLDER
  IFRUN BP<30.1
  ANDRUN BP-SLOW-RISE
  ANDRUN SWDIR
  BECAUSE WINDS FROM S OR SW
  AND CLOUDS INDICATE WEATHER IMPROVING
THEN CLEARING WITHIN A FEW HRS
ANDTHEN FAIR NEXT SEVERAL DAYS
ANDTHEN WEATHER IMPROVING
-->


Screen 11
( WEATHER PREDICTOR - EXPERT-2        JP)
  IFRUN BP>30.1
  ANDRUN BP-SLOW-FALL
  ANDRUN EDIR
```

Weather Predictor

(Listing Continued, text begins on page 24)

```
BECAUSE WINDS FROM NE, E, OR SE
AND CLOUDS INDICATE BAD WEATHER
THEN RAIN IN 12 - 18 HRS
ANDTHEN WEATHER TURNING BAD
  IFRUN BP>30.1
AND RUN BP-RAPID-FALL
AND RUN EDIR
BECAUSE WINDS FROM NE, E, OR SE
THEN WINDY, RAIN WITHIN 12 HRS
ANDTHEN WEATHER TURNING BAD
-->


Screen 12
( WEATHER PREDICTOR - EXPERT-2         JP)
  IFRUN BP>30.1
  AND RUN BP-SLOW-FALL
  AND RUN NEDIR
  BECAUSE WINDS FROM NE OR E
  AND SEASON IS WINTER
  BECAUSE NOT SUMMER
  AND CLOUDS INDICATE BAD WEATHER
THEN RAIN WITHIN 24 HRS
ANDTHEN WEATHER TURNING BAD
  IFRUN BP>30.1
  AND RUN BP-RAPID-FALL
  AND RUN NEDIR
  BECAUSE WINDS FROM NE OR E
  AND SEASON IS SUMMER
  BECAUSE NOT WINTER
  AND CLOUDS INDICATE BAD WEATHER
THEN RAIN WITHIN 12 - 24 HRS
ANDTHEN WEATHER TURNING BAD
-->


Screen 13
( WEATHER PREDICTOR - EXPERT-2         JP)
  IFRUN BP<29.8
  AND RUN BP-RAPID-FALL
  AND RUN NDIR
  BECAUSE WINDS FROM N, NE, OR E
THEN SEVERE STORM WARNING
ANDTHEN SEVERE NORTHEAST GALES
ANDTHEN WEATHER TURNING BAD
  IFRUN BP<29.8
  AND RUN BP-RAPID-FALL
  AND RUN SEDIR
  BECAUSE WINDS FROM E, SE, OR S
THEN SEVERE STORM WARNING
ANDTHEN RAIN OR SNOW IMMINENT
ANDTHEN WEATHER TURNING BAD
-->

Screen 14
( WEATHER PREDICTOR - EXPERT-2         JP)
  IFRUN BP<30.1
  AND RUN BP-SLOW-FALL
  AND RUN EDIR
  BECAUSE WINDS FROM SE, E, OR NE
  AND CLOUDS INDICATE BAD WEATHER
THEN RAIN FOR NEXT DAY OR TWO
ANDTHEN WEATHER TURNING BAD
-->
```

```
Screen 15
( WEATHER PREDICTOR - EXPERT-2          JP)
  IFRUN BP<30.1
  ANDRUN BP-RAPID-FALL
  ANDRUN EDIR
  BECAUSE WINDS FROM SE, E, OR NE
  AND CLOUDS INDICATE BAD WEATHER
THEN RAIN WITH HIGH CLOUDS
ANDTHEN CLEARING WITHIN 24 HRS
ANDTHEN COOLER TEMPERATURES
ANDTHEN WEATHER TURNING BAD
-->

Screen 16
( WEATHER PREDICTOR - EXPERT-2          JP)
  IFRUN BP>30.1
  ANDRUN BP-RAPID-RISE
  ANDRUN WDIR
  BECAUSE WINDS FROM SW, W, OR NW
THEN FAIR TODAY
ANDTHEN RAIN & WARMER NEXT 48 HRS
ANDTHEN WEATHER TURNING BAD
  IFRUN BP>30.1
  ANDRUN BP-SLOW-FALL
  ANDRUN WDIR
  BECAUSE WINDS FROM SW, W, OR NW
THEN WARMER
ANDTHEN RAIN WITHIN 24 - 36 HRS
ANDTHEN WEATHER TURNING BAD
-->

Screen 17
( WEATHER PREDICTOR - EXPERT-2          JP)
  IFRUN BP>30.1
  ANDRUN BP-RAPID-FALL
  BECAUSE RAPID FALL > .06" PER HOUR
  ANDRUN WDIR
  BECAUSE WINDS FROM SW, W, OR NW
  AND CLOUDS INDICATE BAD WEATHER
THEN WARMER
ANDTHEN RAIN WITHIN 18 - 24 HRS
ANDTHEN WEATHER TURNING BAD
  IFRUN BP>30.1
  ANDRUN BP-SLOW-FALL
  ANDRUN SDIR
  BECAUSE WINDS FROM SE OR S
THEN RAIN WITHIN 24 HRS
ANDTHEN WEATHER TURNING BAD
  IFRUN BP>30.1
  ANDRUN BP-RAPID-FALL
  ANDRUN SDIR
  BECAUSE WINDS FROM SE OR S
THENHYP WEATHER TURNING BAD
ANDTHEN WINDY, RAIN WITHIN 12 HRS
-->

Screen 18
( WEATHER PREDICTOR - EXPERT-2          JP)
( closing statement generator )
  IFNOT WEATHER OK
  ANDNOT WEATHER TURNING BAD
  ANDNOT WEATHER IMPROVING
THENHYP INSUFFICIENT DATA FOR A FORECAST

DONE
;S
```

End Listing

RSA: A Public Key Cryptography System, Part II

Welcome back for the second pass. If you are like I am, you have read or scanned between twenty and forty magazines in the month that has passed since you read the first part of this article. (It keeps getting harder to keep up with technology.) Therefore, we should briefly review what we presented in Part I.

Part I Revisited

We first learned that Diffie and Hellman introduced the Public Key System (PKS) in 1976[1] and expanded on it in 1979.[2] In 1978, Rivest, Shamir, and Adleman (RSA) proposed a method of PKS implementation and introduced signing messages with digital signatures.[3]

Then we ran through a brief introduction to RATFOR (RATional FORtran), the implementation language. RATFOR, which is similar in many ways to C and Pascal, adds several structured constructs to Fortran, and the output of the RATFOR precompiler is Fortran. Sample compilation of a RATFOR "program" to source was provided in Listings One to Three.

Next we reviewed modulo arithmetic. We saw that modulo arithmetic has some advantages and disadvantages. It limits the size of the numbers that are generated, but it requires considerable extra computation. As we will see, modulo arithmetic is an integral part of the RSA cryptosystem.

Next we looked at multiple-precision arithmetic. The multiple-precision arithmetic algorithms implemented in Part I form the core of the RSA system that we will present here. Knuth's algorithms[4] were stated and Listing Four showed the RATFOR source that implemented these algorithms. Four algorithms were presented: addition, subtraction, multiplication, and division. We saw that these algorithms were very computation intensive and that optimization efforts should probably be directed at them.

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Finally, the Russian Peasant exponentiation algorithm was presented. The basic algorithm found in Knuth's book had to be expanded to include modulo operations (a requirement of RSA). The Russian Peasant algorithm is most suitable for use with large exponents, which is why it was selected. Listing Five showed the RATFOR implementation of the algorithm.

In this part we will explore the generation and testing of large prime numbers, the generation of public/private keys, the encryption/decryption process, and the use of digital signatures. However, before we get to the meat of things, we must revisit RATFOR.

RATFOR Revisited

One of the advantages of RATFOR is that it allows character and string manipulation, a shortcoming of Fortran. Most of the library functions are associated with character manipulation. Since several of these library routines are used in the programs presented in this part, we will look at them briefly.

Also, RATFOR can have a global definition file, which I have called RATDEF.RAT. RATDEF contains definitions for most of the common ASCII characters, including control characters. For example:

```
define(BIGD,68)  # "D"
define(LETD,100) # "d"
define(BACKSPACE,8)
         # <BACKSPACE> character
```

These are defined because Microsoft's Fortran converts all lower-case characters to upper case, i.e., "d" translates to "D." In addition, RATDEF contains some key word definitions. For example:

```
define(character,byte)
         # make CHARACTER type equi-
valent to BYTE
define(EOS,-2)
         # End of String character
define(EOF,-1) # End of File character
```

The library functions that are used in the RSA-PKS routines include:

CLOSER – closes a file opened by OPENR

EQUAL – compares two strings for equality

ERROR – prints a string then exits to the operating system

GETC – waits for and gets a character from the console

GETCH – gets a character from a Logical Unit specified in OPENR

OPENR – opens a character type file in the specified mode and sets aside a specified buffer size

PUTCH – sends a character to the Logical Unit specified in OPENR

PUTDEC – converts an integer to a character string of specified length and sends the string to the console

REMARK – sends a string to the console without a carriage return; a period terminates the string

RMRKLN – sends a string to the console terminated by a carriage return; a period terminates the string

TYPE – determines if a character is a DIGIT or a LETTER

So much for our revisit to RATFOR. Now, let's talk a little bit about the RSA system and what makes it so advantageous.

Public Key Systems

PKS cryptography systems are different from most other cryptography systems in that a key is made public. Other cryptography systems require that the key be kept secret; otherwise, the system is compromised. Private key methods require that both the sending and the receiving sites have copies of the key so that they can encrypt/decrypt messages. The logistics of key distribution can be very awkward and costly. For communication networks, this process is almost useless.

PKS systems do not have these problems, since everyone has access to the public (encrypting) key; however, only the receiver knows the private (decrypting) key. This dual key approach allows anyone to use the public key to encrypt a message and send it to the holder of the private key, who is the only individual that can decrypt the message. For this system to work, knowledge of the public key cannot provide any knowledge about the private key, which would compromise the cryptography system.

The RSA system generates keys that have 200+ decimal digits by using prime numbers that have 100+ decimal digits. What makes the public key so secure is that efficient methods of factoring a very large number into its prime factors is beyond current mathematical feasibility, even using today's supercomputers. Estimates of factoring such numbers range from hundreds to billions of years, essentially an impossible task.

That is not to say that a mathematical

breakthrough cannot happen tomorrow and render the RSA system useless, similar to the recent breaking of the Knapsack System (a rival PKS). Knuth presents at least four methods of factoring numbers into their prime factors, including the sieve method that is used as a standard benchmark program. To factor a number, N, all we have to do is check all prime numbers less than or equal to the square root of N. If N is less than a million, there are only 168 primes to search. As numbers become larger, a prime number occurs about once every 2.3 times the number of decimal digits in the number. Thus, for 100-digit numbers, a prime number will be found once every 2300 numbers or so.

The problem then becomes: How do we find primes having 100+ decimal digits?

Large Prime Numbers

We are not really interested in factoring numbers into all their prime factors; we simply want to find a pair of prime numbers. Luckily, a theorem exists for this, although it does not say that a number is prime, only that it is not prime. The algorithm is based on Fermat's theorem, which states that if P is prime and X is not a multiple of P, then $(X ** (P-1))$ mod $P = 1$. Unfortunately, the theorem is not an if-and-only-if condition; i.e., it is not necessarily true that if the result is one, then P is prime. However, we *can* say that if the result is not one, then P is definitely not prime. (By the way, now we can see one of the reasons that we were interested in modifying the Russian Peasant algorithm to include modulo arithmetic.)

Knuth points out that some non-prime numbers can give prime testing routines great difficulty, such as $(X ** 560)$ mod $561 = 1$ for all X relatively prime to 561. To overcome this difficulty, Knuth provides a probabilistic prime test. If the algorithm is invoked K times, choosing X independently and randomly, and if the algorithm finds that the number is probably prime in all cases, then there is only one chance in $(4 ** K)$ that the number is not prime. On the other hand, if during any pass the number is found to be non-prime by the algorithm, then the number is definitely not prime.

In Listing Six (page 40) you will find the prime test routine (PRIMTEST.RAT). It is not an implementation of Knuth's Algorithm P but is a simpler offshoot of the Fermat theorem, which we will call algorithm P'. Algorithm P' proceeds as follows:

Prime Test Algorithm

Algorithm P' (probabilistic primality test). Given an odd, positive integer, N, this algorithm attempts to decide whether or not N is prime.

P'1. [Check divisibility by 5] If the least significant digit of N is divisible by 5 (1sd(N) mod 5 = 0), then terminate the algorithm and say N is not prime.

P'2. [Generate exponent] Set $M = N-1$.

P'3. [Initialize] Set K = 1.

P'4. [Generate X] Let X be a random, positive integer in the range $1 < X < N$.

P'5. [Exponentiate] Set $Y = (X ** M)$ mod N.

P'6. [Done?] If K = KTIMES and Y = 1, then terminate the algorithm and say N is probably prime.

P'7. [Not prime] If Y != 1, then terminate the algorithm and say N is not prime.

P'8. [Increase K] Set $K = K+1$ and return to Step P'4.

As with Knuth's Algorithm P, if this algorithm reports that N is not prime, then N is definitely not prime. If N is reported to be probably prime, a finite probability remains that N is not prime! I cannot say with any certainty what the probability is. I checked 561 with the above algorithm and, after fifty attempts to determine the primality of 561, the algorithm only reached five passes (K = 5) once; usually it only reached one or two passes. The algorithm could probably be made more efficient at Step P'1 by using some of the fast divisibility algorithms presented by H. T. Gordon.[5] It should be noted that KTIMES does not have to be 30, as is the case in Listing Six (see the *define* near the beginning of the listing), but reducing the value increases the probability that N will not be prime. The source code in Listing Six provides a test number count as the algorithm is executing (PUTDEC call) so that you can see how the operation is progressing.

Now that we can test for and generate large prime numbers, we are in a position to generate the public and private keys. So, let's look at the key generation procedure.

Key Generation

Before we dig into the generation of the keys, let's look at the RSA–PKS cryptography system in some detail to get a better feel for where we are heading. The RSA system is based on the use of large prime numbers, but in fact, we need only two large prime numbers, P and Q (at least 100 decimal digits each is preferable). The public (encryption) key is formed by the product of P and Q, i.e., $N = P * Q$. As you can imagine, factoring N is nearly impossible since N is a number composed of at least 200 decimal digits.

Next we now take the clear text and break it up into blocks; each block is formed by translating the characters into numbers and concatenating the digits to form a number, X, such that $0 < X < N$. To encrypt the message, we take each block, X, in turn, and form $Y = (X ** 3)$ mod N. Each encrypted block, Y, is then transmitted to the intended receiver. The person receiving the message has a private (decryption) key, generated from P and Q. The private key, D, is formed by picking D such that $D < N$ and $(3 * D)$ mod $((P-1) * (Q-1)) = 1$. We find that $((X ** 3) ** D)$ mod $N = X$; i.e., D is the cube root operator modulo N.

There is one further restriction on X. X must be greater than the cube root of N, otherwise we could easily decrypt the encrypted message using a standard cube root operation, since $(X ** 3)$ mod $N = X ** 3$! What happens if one or more of the blocks does not satisfy this restriction? Well, we just pad the end of the block with random characters until the restriction is satisfied. Also, you might ask at this point, why can't I just take the cube root of Y and retrieve the clear text message, X? Well, think about what happened when you performed the modulo operation. The RSA–PKS operation is called a one-way, trap-door function, because it is easy to go one direction but not the other direction.

A couple of conditions on P and Q should be observed. First, (P-1) and (Q-1) cannot be divisible by 3; if (P-1) or (Q-1) is divisible by 3, then D may not decode the encrypted message because it does not produce a unique cube root. Second, (P-1) and (Q-1) should each have at least one large prime factor, otherwise the potential exists that N could be easily factored (see Knuth). Finally, P/Q should not be close to a simple fraction, or, again, N could be easily factored.

Now let's see how we can generate P and Q to get the keys. If you review the two previous paragraphs, you will agree that generating the keys is our most complex task (and the code will reflect this fact). Listing Seven (page 46) is the key generation routine (GENPKEYS.RAT). It uses every routine that we have discussed so far. Near the beginning of the program, some *defines* specify the KEYLENGTH and the BYTEMODULUS. These values are not set in concrete and can be reduced subject to some caution. BYTEMODULUS can be any value between 2 and 128, at the potential cost of increasing computation time in the prime testing routine. KEYLENGTH can be reduced at the risk of increasing the key's vulnerability. The length of 90 is equivalent to a 630-bit key length $((2 ** 7) ** 90)$ or a 190-decimal-digit key.

Our first task is to seed the random number generator. We must do this because the Microsoft random number routine always generates the same sequence of numbers if not seeded, thus giving everyone the same set of keys. We can enter any number from 1 to 9999999. The length of the seed is probably the main weakness of this RSA implementation. Given enough resources and the

Microsoft random number generation routine algorithm, I theoretically could go to any supercomputer or use a room full of microcomputers to generate all 10 million key sets. Then when you publish your public key, I could look it up, get the corresponding private key, and decrypt any of your messages. You could make my task more difficult by changing the value of KEYLENGTH and the associated values for the length of P and Q.

The program randomly selects the length of P and Q while assuring that their ratio is not a simple fraction. Then both P and Q are generated by the same routine, GENPRM. The algorithm used by GENPRM follows.

Generate Prime Algorithm

Algorithm GENPRM (generate a prime number – after Knuth). This algorithm generates a prime number of length Len(PRIME).

G1. [Get length of PO and K] Set Len(PO) = 2 * Len(PRIME) / 3 and Len(K) = Len(PRIME) – Len(PO). Here, PO will end up being a prime number used to generate PRIME, and K will be a multiplying factor used to generate PRIME.

G2. [Generate random PO and K] Generate a random value of PO of length Len(PO). Generate a random value of K of length Len(K).

G3. [Make PO odd and K even] If PO is even, set PO = PO + 1. If K is odd, set K = K – 1. (PO must be odd to be a prime.)

G4. [PO prime?] Test primality of PO. If PO is probably prime, go to Step G6. (Approximately Ln(PO)/2 numbers will be tested before finding a prime.)

G5. [PO not prime] Set PO = PO + 2 and repeat Step G4.

G6. [Adjust K] Find K such that K is even and K mod 3 = PO mod 3.

G7. [Get trial PRIME] Set PRIME = K * PO + 1.

G8. [PRIME prime?] Test primality of PRIME. If PRIME is probably prime, go to Step G10. (Approximately Ln(K * PO)/6 numbers will be tested before finding a prime.)

G9. [PRIME not prime] Set PRIME = PRIME + (6 * PO) and repeat Step G8. (Keeping K mod 3 = PO mod 3, K even and PRIME = K * PO + 1, implies K = K + 6 at this iteration. Substituting the new K into the equation for PRIME produces the above PRIME adjustment.)

G10. [Check PRIME + 1] Find greatest common divisor of (PRIME + 1). If the greatest prime factor of (PRIME + 1) is less than 100,000, then go to Step G9. Otherwise, terminate the

algorithm and return PRIME.

With P and Q determined, N and D can be calculated. The public key, N, is found by N = P * Q and is of length Len(P) + Len(Q) = KEYLENGTH. The private key, D, is found by picking D such that D < N and (3 * D) mod (P−1) * (Q−1) = 1. This statement is equivalent to D = 2 * ((P−1) * (Q−1) + 1)/3. We now have our keys. After all of this trouble, we certainly want to save them so that they can be used in the encryption/decryption process.

Finally, we are ready to look at the encryption/decryption process. This process is the topic of the next section.

Encryption/Decryption

This section will be fairly short because we have already taken a look at the encryption/decryption process in the beginning of the last section. Listing Eight (page 56) depicts the encryption/decryption process. Again, the definitions of BYTEMODULUS and KEYLENGTH appear near the start of EDCRYPT.RAT. Note that they must be the same values used in the key generation process!

We start with a clear text message X = {x1, x2, . . . , xn} where $0 < xk < N$. If xn, the final block, is less than the cube root of N, terminate the block with an End of File character and pad xn by concatenating random characters to the end of the block. I use ASCII control characters except ∧Z, which is the CP/M End of File character. To encrypt each block, xk, we look at each character of xk and strip off the most significant bit; this ensures that $0 < char(xk) < BYTEMODULUS$. Next we form yk = (xk ** 3) mod N. Finally, we scan each character of yk and set the most significant bit if it is an ASCII control character. This action prevents a ∧Z or some other character from giving the RATFOR library character I/O routines a problem. The encrypted blocks Y = {y1, y2, . . . , yn} are sent to the intended receiver. Note that each encrypted block, yk, is of length Len(N).

As the authorized receiver, we can now apply our private key to decode the received message. We reblock Y into blocks of length Len(N). We take each block, yk, in turn. Each character (byte)

is scanned and the most significant bit is stripped, returning the character to its original encrypted form. We now apply the inversion operation, xk = (yk ** D) mod N. The final block, xn, is scanned from the last character to the first character, looking for the End of File character to find the end of the original text. The blocks X = {x1, x2, . . . , xn} are assembled into the decrypted, clear text. EDCRYPT.RAT performs these basic operations, as well as selecting the files to encrypt/decrypt and the necessary file I/O.

Figure 1 (below) shows an example of generated keys, clear text, encrypted text, and decrypted text. The values of KEYLENGTH and KTIMES have been reduced to permit a reasonable amount of execution time: KEYLENGTH was set to 20 instead of 90 and KTIMES was set to 5 instead of 30. It took almost an hour to find the keys. The encryption of CLEARO.TXT into ENCRYPTO.TXT took about 8 minutes and the decryption of ENCRYPTO.TXT to DECRYPTO.TXT took nearly 3 hours and 19 minutes. When I used the original values of KEYLENGTH and KTIMES, I found that it took almost 11 days to find the keys.

To test how much an arithmetic coprocessor might increase performance, I used an AMD 9511 and the Redding Group APU Fortran library. Unfortunately, the library did not include a random number generator and the Microsoft module would not work with the APU library. Therefore, I had to substitute the random number routine found in Listing Nine (page 59). With these changes and using the reduced example above, I found that the routines executed in about half the time (key generation in 1:20, encryption in 0:04, and decryption in 1:33). The key generation took longer because the random number generator was not the same and different keys were produced. Encryption/decryption, however, used the same keys produced by the non-APU version to keep comparisons as close as possible. As I have pointed out before, optimizing the multiple-precision arithmetic and Russian Peasant routines could also have a great impact on the performance. However, I will leave that as an exercise for the reader.

One final topic remains. How can we sign our messages so that the receiver knows without a doubt that we sent the message? Digital signatures provide the answer.

Digital Signatures

Now, let's consider a little "Spy vs. Spy," a la *Mad Magazine*. White 1 wants to send secure messages to White 2. However, Black is always capable of intercepting the transmitted messages and substituting his own messages. The ability of Black to intercept and substitute causes some significant problems for White,

---

```
Encryption Key - Len(N) & N
20
26 26 45 90 98 127 6 31  90 33 86 86 13 99 35 121 110 14 44 77

Decryption Key - Len(D) & D
20
17 60 30 60 65 127 46 49 0 90 98 46 40 95 60 13 54 79 10 83

First Prime - Len(P) & P
7
38 117 16 61 54 24 123

Second Prime - Len(Q) & Q
13
86 25 25 67 16 80 45 20 85 30 97 67 87
```

Figure 1A.

Keys and Primes generated by GENPKEYS. KEYLENGTH = 20 and KTIMES = 5.

This is an example of th RSA encryption/decryption algorithm. This example uses a KEYLENGTH of 20 (140 binary digits), a P length of 49 binary digits, and a Q length of 91 binary digits. The random number seed was 1111111. The number of prime tests was reduced from 30 to 5. It took approximately 00:52:38 to generate the keys. It took about 00:08:02 to encrypt this file and 03:18:46 to decrypt the encrypted file.

For a RSA with a KEYLENGTH of 90 (630 binary digits), it takes approximately 11 days to generate the keys. The computer used to perform the algorithm was a Z80-based machine running at 4Mhz under MP/M 2.0 (single user).

Figure 1B.

Clear (Plain) Text to be encrypted.

---

e.g., risk of dirty tricks and crosses over the eyes. The problem is: How can White 2 be sure that a received message from White 1 is the real thing and not a setup?

A plain text signature is no good, since Black could sign the message with White 1's name before encrypting the message with White 2's public key. How can White use the RSA–PKS system to confirm that a message was sent from a friend? Well, assume that there are public keys that are published in something like the phone book. Anyone is able to take his plain text message and pass it through the encryption algorithm using another's published, public key.

Let's go back to the section on mod-ulo arithmetic and look at the commuta-tive property of exponentiation. Eureka! We find that:

$$X = ((X ** 3) ** D) \bmod N$$
$$= (((X ** 3) \bmod N) ** D) \bmod N$$
$$= (((X ** D) \bmod N) ** 3) \bmod N$$

That is, the order of encryption and decryption algorithm application does not matter. Thus, White 1 can take his clear text message, X, and encrypt it with his private key, Dm, then encrypt the result further with White 2's public key, Ni, and finally send the resultant message, Y, to White 2.

$$Ysig = (X ** Dm) \bmod Nm$$
<– Signature Step

$$Y = (Ysig ** 3) \bmod Ni$$
<– Encryption Step

Black cannot read the message! To read the message, White 2 reverses the process. The message, Y, is passed through White 2's private algorithm, Di, and a yet garbled message is found. White 2 suspects that White 1 sent the message, so White 2 passes the resultant message through White 1's encryption algorithm, Nm, and, lo and behold, White 1's plain text mes-sage, X, appears.

$$Ysig = (Y ** Di) \bmod Ni$$
<– Decryption Step

$$X = (Ysig ** 3) \bmod Nm$$
<– Signature Verification Step

```
81 7F 6B 85 22 85 21 66    8E 40 99 3A 5A 35 8D 40    .~k.".!f.@.:Z5.@
65 9D 9D 6B 82 3C 94 74    25 8E 72 47 68 2B 50 30    e..k.<.t%.rGh+PO
68 62 6A 4D 2D 30 7A 87    8A 9C 59 28 4A 6F 5C 2C    hbjM-Oz...Y(Jo\,
8E 99 20 8E 4D 45 63 28    83 7E 32 6A 98 43 2A 8D    .. .MEc(.~2j.C*.
45 70 5C 4A 61 3A 4F 63    92 6E 5F 90 44 9A 69 2E    Ep\Ja:Oc.n_.D.i.
89 82 87 9D 6F 82 6E 7F    52 61 29 4A 95 75 9A 37    ....o.n~Ra)J.u.7
49 8D 2B 70 82 85 29 98    5E 67 9E 40 25 62 9B 37    I.+p..).^g.@%b.7
9B 60 91 44 4A 73 7C 6C    83 90 4E 94 58 36 78 70    .`.DJs|l..N.X6xp
46 56 95 75 2B 82 52 7A    34 9C 46 8B 88 7A 45 3B    FV.u+.Rz4.F..zE;
59 98 38 29 8B 3C 9C 71    44 64 2A 48 7E 26 6C 7D    Y.8).<.qDd*H~&l}
8D 6B 72 3D 52 55 4B 28    4A 3A 48 6B 49 39 81 9B    .kr=RUK(J:HkI9..
43 57 42 67 8C 4C 9D 5F    58 43 2E 39 8E 34 48 36    CWBg.L._XC.9.4H6
3F 89 2F 49 97 62 86 67    93 85 42 8C 55 38 58 49    ?./I.b.g..B.U8XI
9E 46 58 78 21 6C 63 42    38 25 9C 59 8F 6C 25 7B    .FXx!lcB8%.Y.1%{
77 4E 82 33 51 65 86 98    32 88 5A 2E 83 42 6C 33    wN.3Qe..2.Z..B13
94 5E 56 6B 8D 33 83 51    8A 25 87 4F 99 64 99 83    .^Vk.3.Q.%.O.d..
7F 43 5A 58 8F 23 9C 8A    75 68 70 73 7F 3E 9F 80    ~CZX.#..uhps~>..
5E 75 8F 5F 3E 7B 63 43    8B 87 6F 44 83 24 5E 48    ^u._>{cC..oD.$^H
89 7F 89 81 71 5C 84 37    61 5F 63 41 8D 34 62 8C    .~..q\.7a_cA.4b.
3E 3C 83 42 58 29 64 72    56 46 60 63 7E 42 9D 79    ><.BX)drVF`c~B.y
8B 22 93 20 2A 71 37 6D    70 9E 63 99 2F 2D 61 3A    ." .*q7mp.c./-a:
57 4A 85 25 80 64 87 52    39 7E 4D 38 7B 36 89 2F    WJ.%.d.R9~M8{6./
59 6F 41 6A 76 7B 91 57    8B 8B 52 64 31 92 65 4F    YoAjv{.W..Rd1.eO
50 35 90 90 91 6B 57 7C    51 87 9F 9E 97 89 40 77    P5...kW|Q.....@w
67 74 4E 2A 34 31 65 7D    3D 97 63 9C 67 61 99 52    gtN*41e}=.c.ga.R
92 8A 99 43 97 6A 4E 44    57 91 26 7F 8A 78 2A 6C    ...C.jNDW.&~.x*l
67 54 3E 6B 84 75 94 88    35 71 56 90 55 51 41 30    gT>k.u..5qV.UQAO
61 2B 8F 4D 3D 43 6C 81    89 2F 22 69 9B 5E 6A 8B    a+.M=Cl../"i.^j.
61 77 95 36 75 3E 6A 67    7F 21 5C 4C 84 63 20 93    aw.6u>jg~!\L.c .
77 7E 6D 86 21 7A 8B 59    64 3C 5E 4F 87 7F 39 29    w~m.!z.Yd<^O.~9)
8C 8C 5A 45 8C 36 9E 4E    29 61 73 63 2D 7F 7C 33    ..ZE.6.N)asc-~|3
6D 86 25 94 86 62 7F 8E    3F 44 7D 35 5B 7D 66 27    m.%..b~.?D}5[}f`
7E 7F 99 9F 71 5C 30 34    88 51 25 80 78 43 90 8D    ~~..q\04.Q%.xC..
49 8A 9D 99 48 95 94 8C    63 38 2D 2E 96 9E 26 70    I...H...c8-...&p
2B 46 4E 50 3D 7D 81 59    66 52 87 45 3B 40 26 93    +FNP=}.YfR.E;@&.
80 82 4C 8A 3E 37 4A 81    88 6D 24 2D 25 8F 65 85    ..L.>7J..m$-%.e.
20 67 45 90 97 66 7D 67    39 6F 20 23 29 46 72 7F    gE..f}g9o #)Fr~
97 85 70 47 63 9B 3B 41    8D 27 44 87 56 50 2A 41    ..pGc.;A.'D.VP*A
2B 45 45 6E 5F 74 79 44    85 2A 23 8B 91 2E 6C 89    +EEn_tyD.*#...l.
8B 4E 8F 76 8E 24 94 9F    42 85 6D 78 9A 63 48 22    .N.v.$..B.mx.cH"
92 2B 6F 5C 71 34 43 59    6A 56 96 30 47 9F 25 30    .+o\q4CYjV.OG.%O
9D 35 85 8A 87 3C 88 68    35 62 81 93 84 5D 97 61    .5...<.h5b...].a
51 7C 87 5D 67 93 66 7C    1A                         Q|.]g.f|.
```

Figure 1C.

Encrypted Text using Encryption Key in Figure 1A and Clear Text in Figure 1B.

White 2 is now sure that the message came from White 1 because no one else could know Dm! Black could not make a substitution because he cannot know White 1's private key, unless White 1 has sloppy security procedures.

To review the process, the encryption of a message using digital signatures consists of:

$$Y = (((X ** Dm) \bmod Nm) ** 3) \bmod Ni$$

The decryption process consists of:

$$X = (((Y ** Di) \bmod Ni) ** 3) \bmod Nm$$

That is all there is to it. For further information on the use of digital signatures, you should read the articles in IEEE *Computer Magazine,* February 1983.[6-9]

Summary

We have covered a lot of material in these two parts. In the first part we looked at RATFOR, modulo arithmetic, multiple-precision arithmetic, and the Russian Peasant algorithm. We saw that these routines form the core of the RSA–PKS cryptography system.

In this second part we revisited RATFOR to look at the RATDEF file and some of the RATFOR library routines. The RATDEF file contains some basic keyword and character definitions. The library routines add character I/O to the basic Fortran library.

Next, we looked at Public Key Systems and found what made them advantageous. We do not have to worry about key distribution, because the public key can be made accessible to anyone without worrying about compromising the encryption/decryption process. The public key does not provide any information about the private key.

Third, we looked at the generation of large prime numbers. We used Fermat's theorem to test odd, positive numbers for primality. The test was a probabilistic test that has a finite probability that a reported prime number is not prime, even though the probability is very small. However, if a number is said to be not prime, then it is definitely not prime.

Fourth, we looked at the generation of the public and private keys. We used all of the previously developed routines to generate the prime factors of the keys,

```
54 68 69 73 20 69 73 20   61 6E 20 65 78 61 6D 70    This is an examp
6C 65 20 6F 66 20 74 68   65 20 52 53 41 20 65 6E    le of the RSA en
63 72 79 70 74 69 6F 6E   2F 64 65 63 72 79 70 74    cryption/decrypt
69 6F 6E 0D 0A 61 6C 67   6F 72 69 74 68 6D 2E 20    ion..algorithm.
20 54 68 69 73 20 65 78   61 6D 70 6C 65 20 75 73     This example us
65 73 20 61 20 4B 45 59   4C 45 4E 47 54 48 20 6F    es a KEYLENGTH o
66 20 32 30 20 28 31 34   30 0D 0A 62 69 6E 61 72    f 20 (140..binar
79 20 64 69 67 69 74 73   29 2C 20 61 20 50 20 6C    y digits), a P l
65 6E 67 74 68 20 6F 66   20 34 39 20 62 69 6E 61    ength of 49 bina
72 79 20 64 69 67 69 74   73 2C 20 61 6E 64 0D 0A    ry digits, and..
61 20 51 20 6C 65 6E 67   74 68 20 6F 66 20 39 31    a Q length of 91
20 62 69 6E 61 72 79 20   64 69 67 69 74 73 2E 20     binary digits.
20 54 68 65 20 72 61 6E   64 6F 6D 20 6E 75 6D 62     The random numb
65 72 0D 0A 73 65 65 64   20 77 61 73 20 31 31 31    er..seed was 111
31 31 31 31 2E 20 20 54   68 65 20 6E 75 6D 62 65    1111.  The numbe
72 20 6F 66 20 70 72 69   6D 65 20 74 65 73 74 73    r of prime tests
20 77 61 73 0D 0A 72 65   64 75 63 65 64 20 66 72     was..reduced fr
6F 6D 20 33 30 20 74 6F   20 35 2E 20 20 49 74 20    om 30 to 5.  It
74 6F 6F 6B 20 61 70 70   72 6F 78 69 6D 61 74 65    took approximate
6C 79 20 30 30 3A 35 32   3A 33 38 0D 0A 74 6F 20    ly 00:52:38..to
67 65 6E 65 72 61 74 65   20 74 68 65 20 6B 65 79    generate the key
73 2E 20 20 49 74 20 74   6F 6F 6B 20 61 62 6F 75    s.  It took abou
74 20 30 30 3A 30 38 3A   30 32 20 74 6F 0D 0A 65    t 00:08:02 to..e
6E 63 72 79 70 74 20 74   68 69 73 20 66 69 6C 65    ncrypt this file
20 61 6E 64 20 30 33 3A   31 38 3A 34 36 20 74 6F     and 03:18:46 to
20 64 65 63 72 79 70 74   20 74 68 65 20 65 6E 63     decrypt the enc
72 79 70 74 65 64 0D 0A   66 69 6C 65 2E 0D 0A 0D    rypted..file....
0A 46 6F 72 20 61 20 52   53 41 20 77 69 74 68 20    .For a RSA with
61 20 4B 45 59 4C 45 4E   47 54 48 20 6F 66 20 39    a KEYLENGTH of 9
30 20 28 36 33 30 20 62   69 6E 61 72 79 20 64 69    0 (630 binary di
67 69 74 73 29 2C 20 69   74 0D 0A 74 61 6B 65 73    gits), it..takes
20 61 70 70 72 6F 78 69   6D 61 74 65 6C 79 20 31     approximately 1
31 20 64 61 79 73 20 74   6F 20 67 65 6E 65 72 61    1 days to genera
74 65 20 74 68 65 20 6B   65 79 73 2E 20 20 54 68    te the keys.  Th
65 0D 0A 63 6F 6D 70 75   74 65 72 20 75 73 65 64    e..computer used
20 74 6F 20 70 65 72 66   6F 72 6D 20 74 68 65 20    to perform the
61 6C 67 6F 72 69 74 68   6D 20 77 61 73 20 61 20    algorithm was a
5A 38 30 2D 62 61 73 65   64 0D 0A 6D 61 63 68 69    Z80-based..machi
6E 65 20 72 75 6E 6E 69   6E 67 20 61 74 20 34 4D    ne running at 4M
68 7A 20 75 6E 64 65 72   20 4D 50 2F 4D 20 32 2E    hz under MP/M 2.
30 20 28 73 69 6E 67 6C   65 20 75 73 65 72 29 2E    0 (single user).
0D 0A 1A                                             ...
```

Figure 1D.

Decrypted Text using Decryption Key in Figure 1A and Encrypted Text in Figure 1C.

P and Q. These prime numbers were then used to generate the public and private keys, N and D, respectively.

Next, we used the keys to encrypt and decrypt plain text messages. We saw how to block out the messages to encode or decode them. We found that we must pad the last block if it is too short to keep someone from easily decrypting it. Also, we found that using an APU can halve the time required to perform the encryption/decryption process.

Finally, we looked at digital signatures. We saw that another person could sign a message digitally by encrypting the message initially with their private key and then encrypting the result with the intended receiver's public key. The receiver reverses the process to confirm that the message was sent by the signee.

I hope you have enjoyed this article and found it understandable and educational. I owe Donald Knuth most of the credit for this article, since his book was the major contributor to my understanding of the RSA–PKS cryptography system. Of course, this article would not have been possible without the basic papers of Diffie and Hellman and of Rivest,

Shamir, and Adleman. As I stated previously, I am sure there are better ways to implement the algorithms and even better and more efficient algorithms that could be used to improve the performance. I invite you to improve any part of what I have presented.

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RSA (Text begins on page 30)
Listing Six

```
# PROGRAM NAME:  PRIMTEST.RAT
# PURPOSE:    This routine is a probabilistic test for the primality of a number.
#             If the test is true, there is a probability of (1/4)**K that the
#             number is not prime.  The test is based on Fermat's Theorem.
#             (re. D.E. Knuth, The Art of Computer Programming, V. 2
#             (Semi-Numerical Algorithms), 2nd Ed.,(Addison-Wesley, Reading, MA),
#             pp. 374-380.)
#
# LANGUAGE:  RATFOR
# AUTHOR:  CEB
# USAGE:   LOGCL=PRMTST(PRIMEQ,LENP,WORK,LENW,MODULO)
#              <PRMTST -- Logical function, indicates that the number in
#                      question (PRIMEQ) is probably prime with probability
#                      1-(1/4)**K.  PRMTST = .TRUE. if probably prime,
#                      and .FALSE. if it is definately not prime.
#              >PRIMEQ -- Byte array, contains the number (byte modulus MODULO)
#                      to be tested.  The MSDigit(s) are in PRIMEQ(1) and
#                      the LSDigit(s) are in PRIMEQ(LENP).
#                      CAUTION:  PRIMEQ > 2
#              >LENP -- Integer variable, defines length of PRIMEQ array.
#              <WORK -- Byte array, a working array needed to do test.
#              >LENW -- Integer variable, defines length of WORK array.
#                      CAUTION:  LENW = 7 * LENP + 3
#              >MODULO -- Integer variable, defines the arithmetic modulus that
#                      is to be used.  MODULO has a byte-wide effect and
#                      should be between 2 and 128 (e.g. 100 for Decimal
#                      Numbers and 128 for ASCII Characters).
#
```

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RSA (Listing Continued, text begins on page 30)
Listing Six

```
# ARRAYS USED: PRIMEQ(*)
# EXTERNALS:  RPEXP,MPSUBT
# UPDATE HISTORY:  INITIAL RELEASE -- 01/25/83 CEB
#
define(KTIMES,30) # number of times, K, to test PRIMEQ

logical function prmtst(primeq,lenp,work,lenw,modulo)

     character back3(4)
     byte primeq(1),work(1),one,notprm
     logical xflag
     integer lenp,lenw,modulo

     data back3/BACKSPACE,BACKSPACE,BACKSPACE,EOS/

     data one/1/
     data rndnxt/.5/ # make RAN(*) generate the next random number
                  #  (any positive number will do)
                  #  *** could seed it with a prior call to RAN(-.xxxxxx)
                  #       where <xxxxxx> is entered by the user ***

     prmtst=.false. # initialize to indicate not prime
     #######
     #
     # The following test assumes PRIMEQ != 5 (i.e. the prime number 5 is not
     #  sought).  If assumption is invalid them remove this test.

     lsdig=primeq(lenp) # get the LSDigit of PRIMEQ
     if ((modulo > 5) & (mod(lsdig,5) == 0)) # PRIMEQ divisible by 5?
          return # no need to check further so Exit !!!

     #
     #######
     idxprm=1 # initialize pointer to 1st digit of PRIMEQ
     while (primeq(idxprm) == 0 & idxprm <= lenp) # digit is zero
          idxprm=idxprm+1 # move index to next position
     lenp0=lenp-idxprm+1 # get length of right justified prime
     if ((lenw >= 7*lenp0+3) & (lenp0 >= lenp) &
         ((lenp0 != 1) | (primeq(1) > 2))) # work array has sufficient length
                                       #  and PRIMEQ > 2?

          {
          idxexp=lenp0+1 # get index to exponent
          call mpsubt(primeq(idxprm),lenp0,one,1,work(idxexp),lenp0,modulo)
                                    # generate (PRIMEQ - 1), i.e. exponent
          idxrlt=idxexp+lenp0 # get index to RESULT: X ** (PRIMEQ-1) mod PRIMEQ
          idxscr=idxrlt+lenp0 # get index to scratch work area for calculations
          fmodul=float(modulo) # float MODULO
          notprm=NO # initialize Not Prime flag
          call remark('    Test Number:  .') # indicate test running
          do i=1,KTIMES # test for primality
               {
               call putdec(i,3) # show test number
               call remark(back3) # backspace 3
               idxx=1 # initialize index to X
               idigit=primeq(idxprm) # get 1st digit of PRIMEQ
               if (idigit == 1) # first digit a 1?
                    {
                    xflag=.true. # indicate 1st digit of X is zero
                    ix=0 # force X < PRIMEQ
                    }

               else # first digit > 1
                    {
```

```
                    xflag=.false. # indicate 1st digit of X is non-zero
                    repeat # get first digit of X
                         {
                         ix=mod(int(fmodul*ran(rndnxt)),idigit)
                                                           # force X < PRIMEQ
                         }
                    until ((lenp0 != 1 & ix != 0) | (lenp0 == 1 & ix > 1))
                                    # 1st digit of X is valid
                    }
            work(idxx)=ix # initialize 1st digit of X
            idxx=idxx+1 # move index to next position
            while (idxx <= lenp0) # generate a random integer
                            #  1 < X < PRIMEQ
                 {
                 if (xflag) # second digit of X needs special handling?
                      {
                      xflag=.false. # indicate 2nd digit of X is non-zero
                      repeat # get second digit of X
                           {
                           ix=mod(int(fmodul*ran(rndnxt)),modulo)
                                                       # force X >= 2
                           }
                      until ((lenp0 != 2 & ix != 0) | (lenp0 == 2 & ix > 1))
                                        # 2nd digit of X is valid
                      }
                 else # PRIMEQ < X < 1

                      ix=mod(int(fmodul*ran(rndnxt)),modulo) # get next
                                                       #  digit
                 work(idxx)=ix # load next digit of X
                 idxx=idxx+1 # move index to next position
                 }
            call rpexp(work,lenp0,work(idxexp),lenp0,work(idxrlt),lenp0,
                   primeq(idxprm),lenp0,work(idxscr),4*lenp0+3,modulo)
                       # generate RESULT = X ** (PRIMEQ -1) mod PRIMEQ
            #
            #   RESULT = (X ** (PRIMEQ-1) mod PRIMEQ) != 1 --> not prime
            #
            idxwrk=idxscr-2 # get index to Digit above LSDigit of RESULT
            if (work(idxwrk+1) != 1) # LSDigit of RESULT --> not prime?
                 notprm=YES # indicate Not Prime
            else # check rest of Digits
                 {
                 while (idxwrk >= idxrlt) # check the rest of RESULT's
                                       #  Digits
                      {
                      if (work(idxwrk) != 0) # not prime?
                           {
                           notprm=YES # indicate Not Prime
                           break # leave WHILE loop
                           }
                      idxwrk=idxwrk-1 # move index to next position
                      }
                 }
            if (notprm == YES) # not prime?
                 break # leave DO loop
            }
       if (notprm == NO) # high probability PRIMEQ is a prime?
            prmtst=.true. # indicate probably prime
       }
  else # cannot test PRIMEQ
       {
```

RSA (Listing Continued, text begins on page 30)
Listing Six

```
        call rmrkln('.')
        call error('Length of [WORK] too small or [PRIMEQ] <= 2 !!!.')
                                        # print message and exit
        }
    return

end
```

End Listing Six

Listing Seven

```
# PROGRAM NAME:  GENPKEYS.RAT
# PURPOSE:   This routine generates the public keys for the RSA
#            (Rivest-Shamir-Adleman) Public Key Encryption/Decryption System.
#            It uses the method described by D.E. Knuth, The Art of Computer
#            Programming, V. 2 (Semi-Numerical Algorithms), 2nd Ed., (Addison-
#            Wesley, Reading, MA), pp. 386-389.
#
# LANGUAGE:  RATFOR
# AUTHOR:  CEB
# USAGE:  ******** (main program)
#
#
#
# ARRAYS USED:   P(KEYLENGTH),Q(KEYLENGTH),N(KEYLENGTH),D(KEYLENGTH),
#                WORK(WORKLENGTH)
# EXTERNALS:  MPADD,MPSUBT,MPMULT,MPDIV,RPEXP,PRMTST
# UPDATE HISTORY:  INITIAL RELEASE -- 01/28/83 CEB
#

define(KEYLENGTH,90) # length of key for characters (assume 80 char. lines)
define(WORKLENGTH,810) # length of WORK buffer for characters
                       #  (9 * KEYLENGTH)
define(BYTEMODULUS,128) # byte modulus for characters (1 character/byte)

define(LUNOUT,6) # output logical unit

program pubkey

    character type
    byte p(KEYLENGTH),q(KEYLENGTH),n(KEYLENGTH),d(KEYLENGTH),work(WORKLENGTH),
         one,two,three
    double precision seed # allow entry of up to 8 characters

    equivalence (seed,work(1))

    data one,two,three/1,2,3/
    data rndnxt/0.5/ # positive number needed to generate next random number
    data modulo/BYTEMODULUS/ # get byte modulus for arithmetic


    write(CONSOLE,100)
100 format(1x,'Enter a seed for the random number generator (XXXXXXX):  ')
    read(CONSOLE,200) seed
200 format(a7)
    rnd1st=0.0 # initialize random number seed
```

```
do i=1,7 # generate random number generator seed
    {
    if (type(work(i)) == DIGIT) # entry is a digit?
        {
        inum=work(i)-DIG0 # get value of seed digit
        rnd1st=10.0*rnd1st+float(inum) # add digit to seed
        }
    else # non-digit found
        break # exit DO loop
    }
rnd1st=-rnd1st/1.0e7 # -1.0 < seed <= 0
call ran(rnd1st) # seed random number generator
len=KEYLENGTH # get key length desired (i.e. length of P * Q)
lendel=4+int(3.0*ran(rndnxt)) # get a delta length between 4 and 7
                              #   (i.e. P and Q length diff. of 8 and 12)
lenp=len/2-lendel # get desired length of prime factor P
lenq=len-lenp # get desired length of prime factor Q
call rmrkln('.')
call rmrkln('Generating P.')
call genprm(p,lenp,work,modulo) # generate prime factor P
call rmrkln('.'); call rmrkln('.')
call rmrkln('Generating Q.')
call genprm(q,lenq,work,modulo) # generate prime factor Q
#
# generate the Public Key encryptor:
#   N = P * Q
#
call rmrkln('.'); call rmrkln('.')
call rmrkln('Generating N.')
call mpmult(p,lenp,q,lenq,n,len,modulo)
idx1=1 # get index to 1st digit of N
```

Listing Seven

```
while (n(idx1) == 0) # find 1st non-zero digit of N
     idx1=idx1+1 # advance to next position
if (idx1 > 1) # need to shift N to make 1st digit non-zero (since modulus
               #   of exponential operation cannot have a leading zero)?
     {
     len=len-(idx1-1) # adjust length of N
     do idx2=1,len # left justify N, stripping zeroes
          {
          idx3=idx1+idx2-1 # get index to digit to move
          n(idx2)=n(idx3) # left justify N
          }
     }
#
# generate Public Key decryptor:
#    D = 2 * ((P - 1) * (Q - 1) + 1) / 3 = 2 * (N - P - Q) / 3 + 1
#
call rmrkln('.')
call rmrkln('Generating D.')
len1=max0(lenp,lenq)+1 # get length of (P + Q)
call mpadd(p,lenp,q,lenq,work,len1,modulo) # generate (P + Q)
idx1=len1+1 # get index to result of N - (P + Q)
call mpsubt(n,len,work,len1,work(idx1),len,modulo) # generate N - (P + Q)
idx2=idx1+len # get index to start of result of 2 * (N - P - Q)
len2=len+1 # get length of result of 2 * (N - P - Q)
     call mpmult(work(idx1),len,two,1,work(idx2),len2,modulo)
                                        # generate 2 * (N - P - Q)
     work(1)=0 # insure Numerator starts with 0 for division normalization
     do idx=1,len2 # copy Numerator <2 * (N - P - Q)>
          {
          work(idx+1)=work(idx2)
          idx2=idx2+1 # advance index to next position
          }
     len1=len2+1 # account for extra digit in numerator
     idx2=len1+1 # get index to quotient of <2 * (N - P - Q) / 3>
     len2=len1-1 # get length of quotient
     idx3=idx2+len2 # get index to remainder of <2 * (N - P - Q) / 3>
     call mpdiv(work,len1,three,1,work(idx2),len2,work(idx3),1,modulo)
                              # generate 2 * (N - P - Q) / 3 (i.e. Quotient)
     len3=len2+1 # get length of result of 2 * (N - P - Q) / 3 + 1
     call mpadd(work(idx2),len2,one,1,work(idx3),len3,modulo)
                              # generate 2 * (N - P - Q) / 3 + 1
     do idx=1,len # copy 2 * (N - P - Q) / 3 + 1 to D from LSDigit to MSDigit
               #  Note: D < N, so Length(D) = Length(N)
               #        LEN3 = LEN + 3
          {
          idx1=len-idx+1 # get index to D
          idx2=idx3+len3-idx # get index to 2 * (N - P - Q) / 3 + 1
          d(idx1)=work(idx2) # move 2 * (N - P - Q) / 3 + 1 to D
          }
     # send Keys to CONSOLE
     call rmrkln('.')
     call rmrkln('Keys and Key Factors.')
     write(CONSOLE,300) len # N
300  format(1x,20(i3,1x))
     write(CONSOLE,300) (n(idx), idx=1,len)
     write(CONSOLE,300) len # D
     write(CONSOLE,300) (d(idx), idx=1,len)
     write(CONSOLE,300) lenp # P
     write(CONSOLE,300) (p(idx), idx=1,lenp)
     write(CONSOLE,300) lenq # Q
     write(CONSOLE,300) (q(idx), idx=1,lenq)
```

```
      # send Keys to LUNOUT file
      call open(LUNOUT,'keys    dat',0) # open Key file on current drive
      write(LUNOUT,400) len # N
400   format(20(i3,1x))
      write(LUNOUT,400) (n(idx), idx=1,len)
      write(LUNOUT,400) len # D
      write(LUNOUT,400) (d(idx), idx=1,len)
      write(LUNOUT,400) lenp # P
      write(LUNOUT,400) (p(idx), idx=1,lenp)
      write(LUNOUT,400) lenq # Q
      write(LUNOUT,400) (q(idx), idx=1,lenq)
      endfile LUNOUT

end

##################################################################################

# GENPRM -- Generate a prime number

subroutine genprm(prime,lenp,work,modulo)

      byte prime(1),work(1),one,two,three,four,six,p0mod3,kadj
      logical prmtst
      integer lenp,modulo

      data one,two,three,four,six/1,2,3,4,6/

      len1=(2*lenp)/3 # get length of prime number needed to generate P or Q
      len2=lenp-len1 # get length of random number needed to generate P or Q
      call rmrkln('     Generating P0.')
      call genrnd(work,len1,modulo) # generate random number P0 to use as
                                    #  initial prime factor
      idx2=len1+1 # get index to start of K
      call rmrkln('     Generating K.')
      call genrnd(work(idx2),len2,modulo) # generate random number K
      if ((work(len1) & 1) == 0) # P0 even?
          work(len1)=work(len1)+1 # make P0 odd
      if ((work(lenp) & 1) != 0) # K odd?
          work(lenp)=work(lenp)-1 # make K even
      idxscr=lenp+1 # get index to work scratch area
      lenscr=7*len1+3 # get length of work scratch area
      call remark('     Making P0 Prime.')
      while (! prmtst(work,len1,work(idxscr),lenscr,modulo)) # P0 not prime
          {
          call mpadd(work,len1,two,1,work(idxscr),len1+1,modulo)
                                                        # P0 = P0 + 2
          if (work(idxscr) >= 1) # P0 overflow, i.e. P0 = 10**(LEN1+1) + 1
              work(idxscr+1)=1 # keep P0 between 10**LEN1 & 10**(LEN1+1),
                               #  i.e. P0 --> 10**LEN1 + 1
          do idxo=1,len1 # move new P0 to old P0
              {
              idxn=idxscr+idxo # get index to new P0
              work(idxo)=work(idxn) # new P0 --> old P0
              }
          call rmrkln('.')
          call remark('                        .')

          }
      do idxo=1,len1 # copy P0 to temporary P0 in preparation for division
          {
          idxn=idxscr+idxo # get index to temporary P0
          work(idxn)=work(idxo) # P0 --> temporary P0
          }
      work(idxscr)=0 # make sure numerator has a leading zero
```

```
    idxquo=idxscr+len1+2 # get index to quotient
    lenquo=len1 # get length of quotient
    idxrem=idxquo+lenquo # get index to remainder
    call mpdiv(work(idxscr),len1+1,three,1,work(idxquo),lenquo,work(idxrem),
            1,modulo) # get PO mod 3
    pOmod3=work(idxrem) # save PO mod 3
do idx0=1,len2 # copy K to temporary K in preparation for division
        {
        idxo=idx2+idx0-1 # get index to K
        idxn=idxscr+idx0 # get index to temporary K
        work(idxn)=work(idxo) # K --> temporary K
        }
    work(idxscr)=0 # make sure numerator has a leading zero
    idxquo=idxscr+len2+2 # get index to quotient
    lenquo=len2 # get length of quotient
    idxrem=idxquo+lenquo # get index to remainder
    call mpdiv(work(idxscr),len2+1,three,1,work(idxquo),lenquo,work(idxrem),
            1,modulo) # get K mod 3
    if (pOmod3 != work(idxrem)) # PO mod 3 != K mod 3?
        {
        # Note: N mod 3 = 0 or 2 or 1 (in that sequence)
        if (pOmod3 == 0) # generate K adjustment for PO mod 3 == 0
            kadj=two*work(idxrem) # adjustment for K mod 3 == 2 or 1
        else if (pOmod3 == 2) # generate K adjustment for PO mod 3 == 2
            kadj=two*(work(idxrem)+one) # adjustment for K mod 3 == 1 or 0
        else # generate K adjustment for PO mod 3 == 1
            kadj=four-work(idxrem) # adjustment for K mod 3 == 0 or 2
        call mpadd(work(idx2),len2,kadj,1,work(idxscr),len2+1,modulo)
                            # force K mod 3 == PO mod 3
        if (work(idxscr) >= 1) # K overflow, i.e. K = 10**(LEN2+1) + 0 or 2
```

Listing Seven

```
            work(idxscr+1)=1 # keep K between 10**LEN2 & 10**(LEN2+1),
                         #  i.e. K --> 10**LEN2 + 0 or 2
      do idx0=1,len2 # move new K to old K
          {
          idxo=idx2+idx0-1 # get index to old K
          idxn=idxscr+idx0 # get index to new K
          work(idxo)=work(idxn) # new K --> old K
          }
      }
call mpmult(work,len1,work(idx2),len2,prime,lenp,modulo) # generate K * P0
lenscr=lenp+1 # get length of work scratch area
call mpadd(prime,lenp,one,1,work(idxscr),lenscr,modulo)
      # generate PRIME = K * P0 + 1 (10**LENP < PRIME < 10**(LENP+1))
do idxo=1,lenp # move new PRIME to old PRIME
      {
      idxn=idxscr+idxo # get index to new PRIME
      prime(idxo)=work(idxn) # new PRIME --> old PRIME
      }
# Note:  Keeping K mod 3 == P0 mod 3 just involves adding 6 to K
#        and (K + 6) * P0 + 1 == (K * P0 + 1) + (6 * P0)
#                             == PRIME + (6 * P0),
#        i.e. we are done with K !!!
idxscr=idx2 # get index to scratch work area
lenscr=len1+1 # get length of scratch area
call mpmult(work,len1,six,1,work(idxscr),lenscr,modulo) # generate 6 * P0
if (work(idxscr) != 0) # extra digit created?
      len1=len1+1 # get new length of 6 * P0
else # no extra digit created
      idxscr=idxscr+1 # advance index to (6 * P0) to skip over zero digit
      do idxo=1,len1 # move (6 * P0) to old P0
          {
          idxn=idxscr+idxo-1 # get index to (6 * P0)
          work(idxo)=work(idxn) # (6 * P0) --> old P0
          }
      idxscr=len1+1 # get index to scratch work area
      lenscr=lenp+1 # get length of work scratch area for adding (6 * P0)
      lenw=7*lenp+3 # get length of work scratch area for prime testing
      call rmrkln('.')
      call remark('        Generating the prime number.')
      while (! prmtst(prime,lenp,work(idxscr),lenw,modulo)) # PRIME not prime
          {
          call mpadd(prime,lenp,work,len1,work(idxscr),lenscr,modulo)
                # PRIME = PRIME + (6 * P0)
          if (work(idxscr) >= 1) # PRIME overflow,
                         #  i.e. PRIME = 10**(LENP+1) + ?
              work(idxscr+1)=1 # keep PRIME between 10**LENP & 10**(LENP+1),
                         #  i.e. PRIME --> 10**LENP + ?
          do idxo=1,lenp # move new PRIME to old PRIME
              {
              idxn=idxscr+idxo # get index to old PRIME
              prime(idxo)=work(idxn) # new PRIME --> old PRIME
              }
          call rmrkln('.')
          call remark('                          .')
          }
      #
      # As a precaution at this point, may want to check that (PRIME + 1) does
      #   not consist entirely of rather small prime factors.  If so, should
      #   return to previous WHILE loop to get next PRIME and try again.
      #
      #     There are 168 primes (PRIMEk) less than 1000.  It is sufficient to
      #     show that if factoring PRIME by PRIMEk (k = 1,168) leaves a number
```

```
    #       which is greater than 10**5 then PRIME okay.
    #
    return

end

###################################################################################

# GENRND -- Generate a random number of specified length

subroutine genrnd(rnd,lenr,modulo)

    byte rnd(1)
    integer lenr,modulo

    data rndnxt/.5/ # make RAN(*) generate the next random number
                    #   (any positive number will do)

    fmodul=float(modulo) # float byte modulus
    idxr=1 # get index to random number
    repeat # force 1st digit to be non-zero
        irnd=mod(int(fmodul*ran(rndnxt)),modulo) # get a random digit
    until (irnd != 0)
    rnd(idxr)=irnd # load 1st digit of RND
    idxr=idxr+1 # point to next digit of RND
    while (idxr <= lenr) # generate the rest of the random digits
        {
        irnd=mod(int(fmodul*ran(rndnxt)),modulo) # get a random digit
        rnd(idxr)=irnd # load random digit
```

RSA (Listing Continued, text begins on page 30)

Listing Seven

```
        idxr=idxr+1 # advance pointer to next digit
        }
    return

end
```

Listing Eight

```
# PROGRAM NAME:  EDCRYPT.RAT
# PURPOSE:  This routine uses the RSA (Rivest-Shamir-Adleman) keys generated
#           by GENPKEYS.RAT to Encrypt/Decrypt a message.
#           It uses the method described by D.E. Knuth, The Art of Computer
#           Programming, V. 2 (Semi-Numerical Algorithms), 2nd Ed., (Addison-
#           Wesley, Reading, MA), pp. 386-389.
#
# LANGUAGE:  RATFOR
# AUTHOR:  CEB
# USAGE:  ******** (main program)
#
#
#
# ARRAYS USED:  N(KEYLENGTH),D(KEYLENGTH),WORK(WORKLENGTH)
# EXTERNALS:  RPEXP
# UPDATE HISTORY:   INITIAL RELEASE -- 02/14/83 CEB
#

define(KEYLENGTH,90) # length of key for characters
define(WORKLENGTH,450) # length of WORK buffer for characters
                       #  (5 * KEYLENGTH)
define(BYTEMODULUS,128) # byte modulus for characters (1 character/byte)

define(LUNIN,6) # Input Logical Unit
define(LUNOUT,7) # Output Logical Unit
define(BLOCKS,8) # Number of 128 byte blocks to use for disk I/O

define(EOFCHAR,26) # text file EOF character (^Z)

program rsa

    character c,c0,getc,getch,oper,type,extn(3),linein(KEYLENGTH),
            lineot(KEYLENGTH)
    byte n(KEYLENGTH),d(KEYLENGTH),work(WORKLENGTH),three,msb
    integer openr,equal,getl
    double precision filnam # allow entry of up to 8 characters

    equivalence (filnam,work(1))

    data extn/BIGT,BIGX,BIGT/ # default file extension
    data modulo/BYTEMODULUS/ # get byte modulus for arithmetic
    data three/3/
    call rmrkln('.') # show options
    call rmrkln('Select operation:.')
    call rmrkln('     <E>ncryption.')
    call rmrkln('     <D>ecryption.')
    call rmrkln('.')
    repeat # get operation to perform
        {
```

```
                    call remark('Operation -- .')
                    oper=getc(c) & 95 # get selection and convert to upper case
                    call rmrkln('.')
                    }
            until ((oper == BIGE) | (oper == BIGD)) # valid entry
            call open(LUNIN,'keys    dat',0) # open key file on current drive
            read(LUNIN,100) lenn # get the length of N
100         format(i3)
            read(LUNIN,200) (n(idx), idx=1,lenn) # get N
200         format(20(i3,1x))
            if (oper == BIGD) # decryption selected?
                    {
                    read(LUNIN,100) lend # get length of D
                    read(LUNIN,200) (d(idx), idx=1,lend) # get D (cube root of text)
                    }
            else # encryption selected
                    {
                    lend=1 # only one digit
                    d(lend)=three # set-up to cube text
                    }
            endfile LUNIN # close key file
            repeat # get name of text file (.TXT extension assumed)
                    {
                    call rmrkln('.')
                    if (oper == BIGD) # decryption selected?
                        call remark('Encrypted.')
                    else # encryption selected
                        call remark('Clear.')
                    call remark(' Text Source File Name (XXXXXXXX) -- .')
                    read(CONSOLE,300) filnam # WORK(1 to 8)
300                 format(a8)
                    do idx=1,8 # convert lower case to upper case
                            if (type(work(idx)) == LETTER) # is character a letter?
                                work(idx)=work(idx) & 95 # to upper case
                    work(9)=EOS # terminate file name
                    }
            until (openr(LUNIN,READONLY,BLOCKS,work,extn) == YES) # successful open
            call scopy(work,1,work,10) # save source file name
            repeat # get name of text file (.TXT extension assumed)
                    {
                    call rmrkln('.')
                    if (oper == BIGD) # decryption selected?
                        call remark('Clear.')
                    else # encryption selected
                        call remark('Encrypted.')
                    call remark(' Text Destination File Name (XXXXXXXX) -- .')
                    read(CONSOLE,300) filnam # WORK(1 to 8)
                do idx=1,8 # convert lower case to upper case
                        if (type(work(idx)) == LETTER) # is character a letter?
                            work(idx)=work(idx) & 95 # to upper case
                work(9)=EOS # terminate file name
                if (equal(work,work(10)) == YES) # Source & Destination names same?
                    {
                    call rmrkln('.')
                    call remark('Source and Destination names cannot be the .')
                    call rmrkln('same !!!.')
                    }
            }
    until (openr(LUNOUT,WRITEONLY,BLOCKS,work,extn) == YES) # successful open
    for (c0=0; c0 != EOF; ) # convert text
        {
        if (oper == BIGE) # encryption?
                {
                linein(1)=0 # make sure LINEIN mod N = LINEIN
                len=2 # point to next character
                }
        else # decryption
                len=1 # initialize line length counter
```

Listing Eight

```
            repeat # get a line of text (encrypted or decrypted)
                {
                linein(len)=getch(c0,LUNIN) & 127 # get a character/strip MSBit
                        # (no effect on Clear Text but MSBit set by encryption)
                len=len+1 # point to next character position
                if ((oper == BIGE) & (c0 == EOF)) # encryption & EOF?
                        {
                        linein(len-1)=EOFCHAR # overwrite EOF with EOFCHAR
                        while (len <= lenn) # clear text line
                            {
                            repeat
                                    num=int(32.0*ran(.5)) # get a number (0 - 31)
                            until (num != EOFCHAR) # not text EOF character
                            linein(len)=num # pad line with a random character
                                            #  (could have previously seeded
                                            #  random number generator !!!)
                            len=len+1 # advance line length counter
                            }
                        }
                }
            until ((c0 == EOF) | (len > lenn)) # ready to encrypt/decrypt
            if ((oper == BIGD) & (c0 == EOF)) # decryption & EOF?
                break # encryption does not leave characters with EOF, so done
            lenin=len-1 # point to last actual character
            lenout=lenn # get length of converted line
            lenw=4*lenn+3 # get length of work area
            call rpexp(linein,lenin,d,lend,lineot,lenout,n,lenn,work,lenw,modulo)
                        # encrypt/decrypt line

    if (oper == BIGD) # decryption selected?
        {
        len0=2 # skip over leading zero from encryption
        msb=0 # MSBit for decrypted text is zero
        if (c0 == EOF) # end of file found on input?
            {
            while ((lineot(lenout) != EOFCHAR) & (lenout >= len0))
                                            # strip pad characters
                    lenout=lenout-1 # backup end of line pointer
            lenout=lenout-1 # backup over EOFCHAR
            }
        }
    else # encryption
        {
        len0=1 # use all characters
        msb=128 # MSBit for encrypted text is one for Control Chars.
        }
    if (lenout >= len0) # at least one character to output?
        {
        do idx=len0,lenout # save encrypted/decrypted text
            {
            c=lineot(idx) # get character to save
            if (c < BLANK) # Control Character (only used during
                            #  encryption)?
                c=c | msb # move Control Characters above printable
                        #  ASCII range, especially EOFCHAR !!!
            call putch(c,LUNOUT)
            }
        }
    }
call closer(LUNIN) # close files
call closer(LUNOUT)
```

end

End Listing Eight

Listing Nine

```
#       Random Number Generator for use with APULIB
#
#             Modeled after W.J. Cody & W. Waite, "Software Manual for the
#                  Elementary Functions"
#

real function ran(x)

      real x,y,const,zero

      data y/100001.0/,const/2796203.0/,zero/0.0/

      if (x == zero)
            ran=y
      else
            {
            if (x < zero)
                  y=-x*const
            y=amod(125.0*y,const)
            ran=y/const
            if (y == zero)
                  y=100001.0
            }
      return

end
```

End Listings

BASICFMT for TRS-80

(A BASIC Formatter)

by Davy Crockett

Davy Crockett, 5807 Cherrywood Ln., Apt. 104, Greenbelt, MD 20770

OK, programmer! It's time to play "Find the Variable"! The object of the game is to determine what the variables are in line 10 of my program. The winner gets a Rubik's Cube!

```
10   IFBASICANDCOMPRESSEDTHEN
SANITY=CRAZYANDYESELSEIFNOT
BASICANDCOMPRESSEDTHENSANIT
Y=CRAZYANDNO
```

The Game

Many programmers like to play games with their programs. I do not mean after you type RUN. These games are played when RUN would only elicit a READY prompt from the BASIC interpreter. They choose variable names like DO. That way, when you use it in an IF statement it can look like this:

100 IF DOOR KNOB THEN 300

Of course, judicious use of spacing could clarify the issue, but that is contrary to the rules of the game.

Now I know what you are thinking: "This guy is leading up to yet another program to format a BASIC ASCII file." You are correct on one count: this is a formatting program. However, it is not just another program. BASICFMT reads a BASIC program in *compressed* format from disk and prints a readable listing. In order to use BASICFMT, all you need is a BASIC program stored on disk without using the ",A" option. No longer is it necessary to load the program and then save it again with the ASCII option.

The Game Pieces

The first byte of these BASIC files is always 0FFH. This byte tells BASIC that it is all right to continue loading this file as a BASIC program. If this byte is not 0FFH then you get the old "Direct Statement in File" error from BASIC. Following this byte, each line is stored in

```
10 '
          ********************** Version 1.0
          *****   DEMO/BAS   ***** 29 October 1982
          ********************** Davy * Crockett

1510 '
          Davy * Crockett, B.S., Esquire
          5807 Cherrywood Lane, Apt. 104
          Greenbelt, Maryland        20770

3010 '
          Program to demonstrate BASICFMT

4510 CLS:CLEAR1000:DEFINTA-Z:DIMNM$(35),AD$(35),TL$(2):GOSUB13510:ONER
RORGOTO12010
6010 DEFFNLN$=LEFT$(NM$(I),INSTR(1,NM$(I),",")-1)
7510 GOTO10510:REM - remark statement
9010 CLS:PRINTTL$(1);:PRINTTL$(2):RETURN
10510 GOSUB9010:I=35:PRINT"Please enter your name (last, first)":INPUT
NM$(I),NM$(I-1):NM$(I)=NM$(I)+", "+NM$(I-1):LN$=FNLN$:PRINT:PRINT"Hell
o, ";LN$:FORK=0TO200:NEXT:GOTO10510
12010 PRINT"ERL =";ERL;"    ERR =";ERR:ONERRORGOTO0
13510 FORK=1TO2:FORM=1TO21:READY:TL$(K)=TL$(K)+CHR$(Y):NEXT:NEXT:RETUR
N
15010 DATA191,131,171,148,32,191,179,179,131,32,191,140,176,140,191,32
,191,131,131,191,10,143,140,142,129,32,143,140,140,140,32,143,32,32,32
,143,32,143,140,140,143,10
```

Figure 1.

BASIC'S LLIST command

a specific format.

The first two bytes of each line contain an address which points to the line number of the following line. The second pair of bytes is the line number in binary format. After this is the actual statement with all of the reserved words compressed into a single byte "token." A binary zero (00H) is the end-of-line marker. For example, the preceding line 100 would, in compressed format on disk, appear as in Figure 1 (page 60). The end of the program is indicated by a next-line pointer consisting of two binary zeros.

Playing The Game

The BASIC tokens are all hex codes which cannot be generated by simply pressing keys on the keyboard. This makes them easy to identify in a section of compressed BASIC code. The character strings associated with each token are stored in the ROM at location 1650H

for both the Model I and the Model III. Once a token is found, it is a simple matter to extract the character string associated with that token from the table stored in the ROM and separate it by blanks before printing.

The first thing I noticed is that BASIC does not always compress a program the way I expected it to, particularly the "remark" commands. The characters REM get translated to a token (93H) as you would expect. However, when I used an apostrophe to indicate a remark, three bytes were inserted instead of a single byte token. The first byte was a colon, the second byte was the remark token, and the third byte was the hex value 0FBH. The 0FBH apparently tells BASIC to back up to the preceding colon and replace it with an apostrophe when listing the program.

The ELSE command contained in an IF statement is another reserved word that is not treated as you would expect.

The ELSE command is always converted to a two-byte sequence with a colon preceding the ELSE token.

Winning At The Game

BASICFMT can help beat these gamesters at their own game by separating the variables and the command words by blanks. The program opens by asking for the filespec of the BASIC file to be formatted. It will then open the file and check to make sure that it is a BASIC program in compressed format. If the file is not in the expected format, an error message is printed and you are returned to the DOS.

When the file checks out, you will be asked for printer specifications. These include: maximum line length, number of lines to print on a page, and number of lines to skip between pages. Printing begins immediately upon entry of the last printer specification. If the printer is not ready, the program will print a mes-

```
  10 '
                  *********************   Version 1.0
                  *****   DEMO/BAS   *****   29 October 1982
                  *********************   Davy * Crockett

1510 '
                  Davy * Crockett, B.S., Esquire
                  5807 Cherrywood Lane, Apt. 104
                  Greenbelt, Maryland        20770

3010 '
                  Program to demonstrate BASICFMT

4510 CLS : CLEAR 1000 : DEFINT A - Z : DIM NM$(35), AD$(35), TL$(2)
     : GOSUB 13510 : ON ERROR GOTO 12010
6010 DEF FN LN$ = LEFT$ (NM$(I), INSTR (1, NM$(I), ",") - 1)
7510 GOTO 10510 : REM - remark statement
9010 CLS : PRINT TL$(1); : PRINT TL$(2) : RETURN
10510 GOSUB 9010 : I = 35 : PRINT "Please enter your name (last,
      first)" : INPUT NM$(I), NM$(I - 1) : NM$(I) = NM$(I) + ", " +
      NM$(I - 1) : LN$ = FN LN$ : PRINT : PRINT "Hello, "; LN$ : FOR
      K = 0 TO 200 : NEXT : GOTO 10510
12010 PRINT "ERL ="; ERL ; "    ERR ="; ERR : ON ERROR GOTO 0
13510 FOR K = 1 TO 2 : FOR M = 1 TO 21 : READ Y : TL$(K) = TL$(K) +
      CHR$ (Y) : NEXT : NEXT : RETURN
15010 DATA 191, 131, 171, 148, 32, 191, 179, 179, 131, 32, 191, 140,
      176, 140, 191, 32, 191, 131, 131, 191, 10, 143, 140, 142, 129,
      32, 143, 140, 140, 140, 32, 143, 32, 32, 32, 143, 32, 143, 140,
      140, 143, 10
```

Figure 2.

BASICFMT output

sage for you to ready the printer and then wait for you to press ENTER.

The target program is read by a routine (GETBYT) which returns a single byte from the disk buffer each time it is called. Some programming languages, notably PL/I, call this technique "character stream" input. One line at a time is read this way, decompressed, and stored in a primary buffer. Extra blanks separating the command words, colons, semi colons, and commas are inserted in the statement at this point. All remark statements and things enclosed within quotation marks are not changed in any way.

Phase two consists of moving the contents of the primary buffer to a secondary buffer. While this is being done, the line breaks are inserted in the statement. Each line is broken at a blank if possible. This would not be possible if there were a remark or quoted string which had no blanks and completely filled the last 10 characters on the line. In this case, the line is broken at the maximum line length specified. Subsequent lines of the same statement are indented past the line number to make the line number easy to read.

Finally, each line in the secondary buffer is printed. The page length is checked before each line to make sure it will not exceed the maximum number of lines specified as page length. Processing continues with the next line until the end of the file is reached. You are then returned to the DOS READY prompt.

Game Strategy

This program will run under any DOS which uses the standard entry points to open an existing file (4424H), read a sector (4436H), and close a file (4428H). It will also run on either the Model I or the Model III. This is due to several factors. First, the BASIC token table is located at the same address in ROM for both models. Second, I used the assembler instruction LD A,(37E8H) to check the printer status instead of IN A,(0FBH), which only works on the Model III. Third, many of the DOS/ROM routines are incorporated in the program so that you do not have to worry about them when switching from one DOS to another. These include keyboard input, screen and printer output, ASCII characters to binary value, binary value to ASCII decimal characters, and even a routine to simulate the single-byte read from disk available with NEWDOS/80.

Some of you may wish to add a routine that prompts for a title line to be printed at the top of each page. I did not feel that this was necessary because I always include an identification section for each program at the beginning. You could also modify the PAGER routine to print a page number when it ejects a page. If your printer requires something separating carriage returns or line feed characters, the LFEED routine will have to be modified to transmit the required codes.

Before I forget, this statement may be easier to read:

10 IF BASIC AND COMPRESSED THEN SANITY = CRAZY AND YES ELSE IF NOT BASIC AND COMPRESSED THEN SANITY = CRAZY AND NO

or in other words:

10 IF BA AND CO THEN SA = CR AND YE ELSE IF NOT BA AND CO THEN SA = CR AND NO

■■J

BASICFMT (Text begins on page 60)

```
00010 ;============================================================*
00020 ;                                                            *
00030 ;          ***********************  Version 3.4              *
00040 ;          *****  BASICFMT  *****  01 November 1983          *
00050 ;          ***********************  Davy * Crockett          *
00060 ;                                                            *
00070 ;        Davy * Crockett, B.S., Esquire                      *
00080 ;        5807 Cherrywood Lane, Apt. 104                      *
00090 ;        Greenbelt, Maryland      20770                      *
00100 ;                                                            *
00110 ;        This program will print a formatted listing         *
00120 ;        of a basic program from disk.  The program          *
00130 ;        must not, I repeat "NOT", be stored in ascii        *
00140 ;        format.  BASICFMT prompts for basic filespec,       *
00150 ;        printer line length, lines per page, and           *
00160 ;        spacing between pages for the output.               *
00170 ;                                                            *
00180 ;============================================================*
00190 ;
00200 ;>>==>   System equates section
00210 ;
00220 TOKTAB  EQU     1650H        ;basic's token table address
00230 LPPORT  EQU     37E8H        ;line printer status port
00240 CTLKEY  EQU     3840H        ;control key address
00250 OPEN    EQU     4424H        ;open existing disk file
00260 READ    EQU     4436H        ;read a record from disk
00270 CLOSE   EQU     4428H        ;close a disk file
00280 DOSERR  EQU     4409H        ;display DOS error
00290 DOS     EQU     402DH        ;return to operating system
00300 ;
00310         ORG     5800H
00320 ;
00330 BASFMT  LD      SP,BASFMT    ;initialize stack
00340 GETFS   LD      DE,FSMSG     ;point to prompt
00350         LD      HL,DCB       ;point to FCB
00360         LD      B,32         ;maximum chars
00370         CALL    INPUTP       ;get filespec
00380         CP      1            ;<BREAK>?
00390         JP      Z,DOS        ;yes, exit
00400         LD      DE,DCB       ;point to FCB
00410         LD      HL,BUFFER    ;point to buffer
00420         LD      B,0H         ;LRECL=256
00430         CALL    OPEN         ;open the file
00440         JP      NZ,DOSERR    ;error, exit
00450         CALL    GETBYT       ;get a byte
00460         CP      OFFH         ;BASIC file?
00470         JR      Z,GETLL      ;yes, continue
00480         LD      HL,FYLMSG    ;no, point to error msg
00490         CALL    PRINT        ; and display message
00500         JR      GETFS        ;return to get filespec
00510 ;
00520 ;>>==>   Get printer characteristics from operator
00530 ;
00540 GETLL   LD      DE,LLMSG     ;point to prompt
00550         CALL    GETVAL       ;get line length
00560         JR      Z,GETPL      ;got something?
00570         SUB     6            ;adjust line length
00580         LD      (LL),A       ; and store it
00590 GETPL   LD      DE,PLMSG     ;point to prompt
00600         CALL    GETVAL       ;get page length
```

```
00610                 JR       Z,GETPS          ;got something?
00620                 LD       (PL),A           ;store page length
00630 GETPS           LD       DE,PSMSG         ;point to prompt
00640                 CALL     GETVAL           ;get page spacing
00650                 JP       Z,BEGIN          ;got something?
00660                 LD       (PS),A           ;yes, store it
00670                 JP       BEGIN            ; and start
00680 ;
00690 ;>>==>   Obtain a decimal value from the keyboard
00700 ;        and convert it to a binary value.
00710 ;
00720 GETVAL          LD       HL,LINE          ;point to buffer
00730                 LD       B,4              ;maximum chars
00740                 CALL     INPUTP           ;get characters
00750                 CP       1H               ;<BREAK>?
00760                 JR       Z,EXIT           ;yes, exit
00770                 CALL     DECBIN           ;convert to binary
00780                 LD       A,E              ;get binary value
00790                 OR       A                ;set/reset Z flag
00800                 RET
00810 ;
00820 DONE            LD       A,0CH            ;form feed
00830                 CALL     LPBYTE           ;eject page
00840 EXIT            LD       DE,DCB           ;point to FCB
00850                 CALL     CLOSE            ;close the file
00860                 JP       DOS              ;exit, stage left
00870 ;
00880 ;>>==>   Prompt and information messages section
00890 ;
00900 FSMSG           DEFW     1F1CH            ;home & cls
00910                 DEFM     '=*= Basicfmt =*= by Davy * Crockett'
00920                 DEFW     0D0DH            ;couple a <CR>'s
00930                 DEFM     'Enter source filespec'
00940                 NOP
00950 LLMSG           DEFB     0DH
00960                 DEFM     '* File successfully opened'
00970                 DEFW     0D0DH
00980                 DEFM     'Enter line length (<ENTER>=85)'
00990                 NOP
01000 PLMSG           DEFB     0DH
01010                 DEFM     'Enter page length (<ENTER>=60)'
01020                 NOP
01030 PSMSG           DEFB     0DH
01040                 DEFM     'Enter page spacing (<ENTER>=6)'
01050                 NOP
01060 FYLMSG          DEFB     0DH
01070                 DEFM     '>>==> File not compressed BASIC'
01080                 NOP
01090 ;
01100 ;        *********************************
01110 ;        * Main program control section. *
01120 ;        *********************************
01130 ;
01140 BEGIN           CALL     LPSTAT           ;check printer status
01150 NXTLYN          CALL     GETLYN           ;buffer a statement
01160                 CALL     PUTLYN           ;print a statement
01170                 JR       NXTLYN           ;go get another
01180 ;
01190 ;>>==>   Store a statement number in the primary buffer
01200 ;
```

```
01210 GETLYN  CALL    GETBYT          ;get a byte
01220          LD      L,A             ; and store in (L)
01230          CALL    GETBYT          ;get another byte
01240          LD      H,A             ; and store in (H)
01250          LD      DE,0            ;test value
01260          CALL    CPHLDE          ;zero address?
01270          JP      Z,DONE          ;yes, quit
01280          LD      DE,LINE         ;point to buffer
01290          CALL    GETBYT          ;LSB line number
01300          LD      L,A             ; and store in (L)
01310          CALL    GETBYT          ;MSB line number
01320          LD      H,A             ; and store in (H)
01330          CALL    BINDEC          ;convert to decimal
01340          LD      DE,LINE+5       ;end of line #
01350          LD      A,' '           ;get a blank
01360          LD      (DE),A          ; and store after #
01370          XOR     A               ;set the quote
01380          LD      (QUOTE),A       ; indicator
01390          LD      (REMARK),A      ; and rem indicator
01400 ;
01410 ;>>==>  Store the rest of the statement in the buffer
01420 ;
01430 BUILD   CALL    GETBYT          ;get a byte
01440          OR      A               ;end of the line?
01450          RET     Z               ;yes, return
01460          CP      80H             ;token?
01470          JR      NC,TOKEN        ;yes, jump
01480 CHK22   CP      '"'             ;quote?
```

(Continued on next page)

BASICFMT (Listing Continued, text begins on page 60)

```
01490          JR      NZ,CHK10        ;no, continue
01500          PUSH    AF              ;save (AF)
01510          LD      A,(QUOTE)       ;get quote value
01520          XOR     1               ;switch it
01530          LD      (QUOTE),A       ; and put it back
01540          POP     AF              ;restore (AF)
01550          JR      CONT            ; and continue
01560 CHK10    CP      0AH             ;line feed?
01570          JR      NZ,VALID        ;no, check validity
01580          LD      A,0DH           ;replace w/ <CR>
01590          JR      CONT            ; and continue
01600 VALID    CP      ' '             ;valid char?
01610          JR      C,BLANK         ;no, zap it
01620          CP      ':'             ;colon character?
01630          CALL    Z,SPACE         ;yes, add space
01640 CONT     INC     DE              ;bump pointer
01650          LD      (DE),A          ;store character
01660          CP      ','             ;comma?
01670          JR      Z,CONT1         ;yes, insert blank
01680          CP      ';'             ;semi-colon?
01690          JR      Z,CONT1         ;yes, insert blank
01700          CP      ':'             ;colon character?
01710 CONT1    CALL    Z,SPACE         ;yes, add a space
01720 CONT2    JR      BUILD           ;go for next
01730 BLANK    LD      A,' '           ;zap the char
01740          JR      CONT            ; and continue
01750 ;
01760 ;>>==>   Obtain characters represented by basic token
01770 ;        from the token table in the basic ROM.
01780 ;
01790 TOKEN    CP      0FCH            ;garbage char?
01800          JR      NC,BLANK        ;yes, zap it
01810          PUSH    AF              ;save token
01820 CHK0FB   CP      0FBH            ;zap REM token?
01830          JR      NZ,CHK95        ;no, continue
01840          DEC     DE              ;backup
01850          DEC     DE              ; to the
01860          DEC     DE              ;  prior
01870          DEC     DE              ;   colon
01880          LD      A,27H           ;change colon
01890          LD      (DE),A          ; to quote (')
01900          POP     AF              ;clear the stack
01910          JR      BUILD           ;go get next character
01920 CHK95    CP      95H             ;ELSE token?
01930          JR      NZ,CHK93        ;no, continue
01940          DEC     DE              ;yes, delete the prior
01950          DEC     DE              ;  space and colon
01960          JR      OFFSET          ;  and continue
01970 CHK93    CP      93H             ;REM token?
01980          JR      NZ,OFFSET       ;no, compute offset
01990          LD      (REMARK),A      ;set REM indicator
02000 OFFSET   SUB     7FH             ;get position in table
02010          LD      B,A             ; and put in (B)
02020          LD      HL,TOKTAB       ;point to token table
02030 TOKEN1   LD      A,(HL)          ;get table character
02040          INC     HL              ;increment the pointer
02050          CP      7FH             ;first character of a token?
02060          JR      C,TOKEN1        ;no, get next character
02070          DJNZ    TOKEN1          ;yes, loop until token
02080          SUB     80H             ;got it, adjust first char
```

```
02090            CALL    SPACE          ;insert space (maybe)
02100  TOKEN2    INC     DE             ;increment buffer pointer
02110            LD      (DE),A         ;and store in buffer
02120            LD      A,(HL)         ;get next character
02130            INC     HL             ;increment token pointer
02140            CP      7FH            ;end of token?
02150            JR      C,TOKEN2       ;no, continue moving token
02160            POP     AF             ;retrieve token
02170            CP      0BCH           ;TAB( token?
02180            JP      Z,BUILD        ;yes, don't put blank
02190            CP      93H            ;REM token?
02200            JP      Z,BUILD        ;yes, don't put blank
02210            JR      BLANK          ;no, insert space
02220  ;
02230  ;>>==>   Conditional blank insert routine
02240  ;        Inserts a blank in the output buffer if:
02250  ;        1)  we are not processing a REM statement, and
02260  ;        2)  we are not processing within quotes,   and
02270  ;        3)  there is not a blank there all ready.
02280  ;
02290  SPACE     PUSH    AF             ;save (AF)
02300            LD      A,(QUOTE)      ;get quote value
02310            OR      A              ;inside quotes?
02320            JR      NZ,XSPACE      ;yes, no insert
02330            LD      A,(REMARK)     ;get rem value
02340            OR      A              ;inside remark?
02350            JR      NZ,XSPACE      ;yes, no insert
02360            LD      A,(DE)         ;get last char
02370            CP      ' '            ;is it a blank?
02380            JR      Z,XSPACE       ;yes, ok
02390            INC     DE             ;bump pointer
02400            LD      A,' '          ;get a blank
02410            LD      (DE),A         ;and add it
02420  XSPACE    POP     AF             ;retrieve (AF)
02430            RET
02440  ;
02450  ;>>==>   At this point, the entire statement is in the
02460  ;        primary output buffer.  This section moves the
02470  ;        statement to the secondary output buffer and
02480  ;        inserts line breaks and indentation as required.
02490  ;
02500  PUTLYN    XOR     A              ;zero (A)
02510            INC     DE             ;mark the end
02520            LD      (DE),A         ; of the statement
02530            LD      DE,OUTLYN      ;set output pointer
02540            LD      HL,LINE        ;set input pointer
02550            LD      A,(LL)         ;get max line length
02560            ADD     A,6            ;add length of #
02570            LD      B,A            ;and put in (B)
02580  PUTNXT    LD      A,(HL)         ;get a byte
02590            LD      (DE),A         ; and store it
02600            INC     HL             ;bump in pointer
02610            INC     DE             ;bump out pointer
02620            CP      0DH            ;<LF>?
02630            JR      Z,SPACER       ;yes, insert indent
02640            OR      A              ;end of input?
02650            JR      Z,ENDLYN       ;yes, jump out
02660            DJNZ    PUTNXT         ;no, continue
02670  ;
02680  ;>>==>   This section will backup and insert a line break
02690  ;        in the secondary buffer when the maximum line
02700  ;        length has been reached.  If a blank is not found
02710  ;        in the last 10 characters stored, the line break
02720  ;        is inserted at the maximum line length location.
02730  ;
02740  BACKUP    LD      (OLDDE),DE     ;save current (DE)
```

(Continued on next page)

BASICFMT   (Listing Continued, text begins on page 60)

```
02750             LD      (OLDHL),HL      ;and current (HL)
02760             LD      B,0AH           ;limiter value
02770 BACK        DEC     HL              ;unbump pointer
02780             DEC     DE              ;unbump pointer
02790             LD      A,(HL)          ;get character
02800             CP      ' '             ;blank?
02810             JR      Z,BREAK         ;yes, break here
02820             DJNZ    BACK            ;no, try another
02830             LD      HL,(OLDHL)      ;can't break it
02840             LD      DE,(OLDDE)      ; use old breaks
02850             DEC     HL              ;unbump for bump
02860             DEC     DE              ;unbump for bump
02870 BREAK       INC     HL              ;skip blank
02880             INC     DE              ;skip blank
02890             LD      A,0DH           ;insert a <CR>
02900             LD      (DE),A          ; in the line
02910             INC     DE              ;bump pointer
02920 SPACER      LD      A,' '           ;get a blank
02930             LD      B,6             ; and blank count
02940 SPACE2      LD      (DE),A          ;insert a blank
02950             INC     DE              ;bump pointer
02960             DJNZ    SPACE2          ;next?
02970             LD      A,(LL)          ;reset the line
02980             LD      B,A             ; length value
02990             JR      PUTNXT          ;go back for more
03000 OLDHL       DEFW    0               ;storage space
03010 OLDDE       DEFW    0               ;storage space
03020 ;
03030 ;>>==>  The statement has now been stored in the secondary
03040 ;       buffer with line breaks and indentation.  ENDLYN
03050 ;       prints the secondary buffer on the line printer
03060 ;       inserting page eject commands as necessary.
03070 ;
03080 ENDLYN      DEC     DE              ;back up one
03090             LD      A,0DH           ;get a <CR>
03100             LD      (DE),A          ; and insert it
03110             INC     DE              ;bump pointer
03120             XOR     A               ;get a EOL char
03130             LD      (DE),A          ; and mark it
03140             LD      HL,OUTLYN       ;point to output
03150 NEXTLN      CALL    PAGER           ;check page break
03160             CALL    LPRINT          ;print a line
03170             INC     HL              ;bump over <CR>
03180             LD      A,(HL)          ;get next char
03190             OR      A               ;EOL?
03200             JR      NZ,NEXTLN       ;no, get more
03210             RET                     ;all done
03220 ;
03230 ;>>==>  This routine generates carriage returns to
03240 ;       eject a page when the page is full.
03250 ;
03260 PAGER       LD      A,(LINES)       ;get line count
03270             INC     A               ; and add 1
03280             LD      (LINES),A       ; and store back
03290             LD      B,A             ;store in (B)
03300             LD      A,(PL)          ;get page length
03310             CP      B               ;too many lines?
03320             RET     NC              ;no, return
03330             LD      A,1             ;yes, repage
03340             LD      (LINES),A       ;reset counter
03350             LD      A,(PS)          ;get page spacing
```

```
03360          LD      B,A             ; and store in (B)
03370          LD      A,0DH           ;<LF> value
03380 LFEED    CALL    LPBYTE          ;feed a line
03390          DJNZ    LFEED           ;more?
03400          RET                     ;back for more
03410 ;
03420 ;>>==>  This routine returns a single byte from the
03430 ;       disk buffer to the caller in register A.
03440 ;
03450 GETBYT   PUSH    HL              ;save (HL)
03460          PUSH    DE              ;save (DE)
03470          LD      HL,(CHRPTR)     ;get byte pointer
03480          LD      DE,ENDBUF       ;and end of buffer
03490          INC     HL              ;bump pointer
03500          CALL    CPHLDE          ;compare them
03510          JR      Z,GREAD         ;need next sector
03520 GETO2    LD      A,(CTLKEY)      ;get control key
03530          BIT     2,A             ;<BREAK>?
03540          JP      NZ,EXIT         ;yes, stop
03550          LD      A,(HL)          ;get the byte
03560          LD      (CHRPTR),HL     ;store pointer
03570          POP     DE              ;retrieve (DE)
03580          POP     HL              ;retrieve (HL)
03590          RET                     ;send back byte
03600 ;        storage for address of last character returned
03610 CHRPTR   DEFW    ENDBUF-1
03620 ;        read next sector of the basic program
03630 GREAD    LD      DE,DCB          ;point to FCB
03640          CALL    READ            ;read a sector
03650          LD      HL,BUFFER       ;point to buffer
03660          JR      Z,GETO2         ; and return byte
03670          OR      80H             ;set return code
03680          CALL    DOSERR          ;display error
03690          JP      EXIT            ;exit, stage left
03700 ;
03710 ;=================================================
03720 ;  BINDEC -  (HL)=binary -  (DE)-->output buffer
03730 ;   The output buffer must be exactly 5 bytes long.
03740 ;=================================================
03750 ;
03760 BINDEC   LD      (BINA),HL       ;save binary number
03770          LD      (BIND),DE       ;save beginning of buffer
03780          INC     DE              ;increment (DE)
03790          INC     DE              ; to the end
03800          INC     DE              ; of the
03810          INC     DE              ; buffer
03820          LD      (BIN4),DE       ;store end of buffer address
03830          LD      B,5H            ;number of bytes in buffer
03840          XOR     A               ;initial value for buffer
03850          LD      HL,(BIND)       ;beginning of buffer
03860 BD1      LD      (HL),A          ;initialize the
03870          INC     HL              ; buffer to all
03880          DJNZ    BD1             ;  binary zeros
03890          LD      HL,(BINA)       ;retrieve the binary number
03900          LD      A,H             ;get the MSB
03910          CALL    BD4             ; and convert to decimal
03920          LD      A,(BINA)        ;get the LSB
03930          CALL    BD4             ; and convert to decimal
03940          LD      HL,(BIND)       ;buffer beginning address
03950          LD      B,5H            ;buffer length
03960 BD2      LD      A,30H           ;convert to ascii value
03970          ADD     A,(HL)          ;convert digit to ascii
03980          LD      (HL),A          ; and store back in buffer
03990          INC     HL              ;increment pointer
04000          DJNZ    BD2             ;all digits done?
04010          LD      HL,(BIND)       ;beginning of buffer
```

(Continued on next page)

69

```
04020              LD      B,4H            ;buffer length less 1
04030 BD3          LD      A,30H           ;ascii zero
04040              CP      (HL)            ;is buffer char "0"
04050              RET     NZ              ;no, return to caller
04060              LD      A,20H           ;yes, change to a blank
04070              LD      (HL),A          ; and store in buffer
04080              INC     HL              ;increment pointer
04090              DJNZ    BD3             ;loop for all but last char
04100              RET                     ;return to calling routine
04110 BINA         DEFW    0H              ;storage for binary value
04120 BIND         DEFW    0H              ;storage beginning of buffer
04130 BIN4         DEFW    0H              ;storage end of buffer
04140 BD4          LD      B,8H            ;bits per byte
04150 BD5          PUSH    BC              ;save loop counter
04160              PUSH    AF              ;save (AF)
04170              CALL    BD6             ;multiply buffer by 2
04180              POP     AF              ;retrieve the value
04190              RLA                     ;shift bit to carry flag
04200              PUSH    AF              ;save the value again
04210              CALL    C,BD9           ;if bit set, add to buffer
04220              POP     AF              ;retrieve value again
04230              POP     BC              ;restore loop counter
04240              DJNZ    BD5             ;continue for all bits
04250              RET                     ;return to caller
04260 ;
04270 ;>>==>  Multiply the contents of the buffer by 2
04280 ;
04290 BD6          LD      HL,(BIN4)       ;end of buffer address
04300              OR      A               ;reset the carry flag
04310              LD      B,5H            ;length of buffer
04320 BD7          LD      A,(HL)          ;get a digit
04330              ADC     A,A             ;multiply by 2 and add carry
04340              CP      0AH             ;less than 10?
04350              JR      C,BD8           ;yes, skip adjustment
04360              SUB     0AH             ;adjust to less than 10
04370 BD8          LD      (HL),A          ;return to buffer
04380              CCF                     ;switch the carry flag
04390              DEC     HL              ;point to next digit
04400              DJNZ    BD7             ;loop for all digits
04410              RET                     ;return to caller
04420 ;
04430 ;>>==>  Add 1 to the contents of the buffer
04440 ;
04450 BD9          LD      HL,(BIN4)       ;end of buffer address
04460              LD      B,5H            ;length of the buffer
04470 BDA          LD      A,(HL)          ;get a digit
04480              INC     A               ;add 1 to the digit
04490              LD      (HL),A          ; and return it to buffer
04500              CP      0AH             ;digit less than 10?
04510              RET     C               ;yes, return to caller
04520              SUB     0AH             ;no, adjust to decimal
04530              LD      (HL),A          ; and return to buffer
04540              DEC     HL              ;point to next digit
04550              DJNZ    BDA             ;loop to handle carry
04560              RET                     ;return to caller
04570 ;
04580 ;=== DECBIN === Calls: nothing
04590 ;==================================================*
04600 ;======  Convert ascii decimal to binary  ======*
04610 ;==================================================*
04620 ;                                                *
```

```
04630 ;  on entry:  (HL) => character string          *
04640 ;             terminated by first non-decimal    *
04650 ;             character encountered              *
04660 ;                                                *
04670 ;  on exit:   (DE) = binary equivalent          *
04680 ;             (HL) => non-decimal character      *
04690 ;                                                *
04700 ;=================================================*
04710 ;
04720 DECBIN  DEC    HL            ;set initial pointer
04730         LD     DE,0H         ;zero the contents DE
04740 DB1     INC    HL            ;point to next character
04750         LD     A,(HL)        ;get next character
04760         SUB    30H           ;adjust to binary
04770         RET    C             ;return if character < "0"
04780         CP     0AH           ;character greater than 9
04790         RET    NC            ;yes, return to caller
04800         PUSH   HL            ;save character pointer
04810         LD     H,D           ;get current value
04820         LD     L,E           ; from DE register
04830         ADD    HL,HL         ;old value times 2
04840         ADD    HL,HL         ;old value times 4
04850         ADD    HL,DE         ;old value times 5
04860         ADD    HL,HL         ;old value times 10
04870         LD     E,A           ;value of new digit
04880         LD     D,0H          ; to DE register
04890         ADD    HL,DE         ;add to the old value
04900         EX     DE,HL         ; and return value to DE
```

(Continued on next page)

```
04910            POP     HL                  ;retrieve character pointer
04920            JR      DB1                 ;go to process next char
04930 ;
04940 ;=== PRINT === Calls: OUTBYT
04950 ;===============================*
04960 ;   Output a line to the screen   *
04970 ;===============================*
04980 ;
04990 PRINT    PUSH    AF                  ;save (AF) reg
05000 PLOOP    LD      A,(HL)              ;get a byte
05010            OR      A                   ;test or 0H
05020            JR      Z,LAFIN             ;yes, jump
05030            CALL    OUTBYT              ;output byte
05040            INC     HL                  ;bump pointer
05050            JR      PLOOP               ;next byte
05060 LAFIN    POP     AF                  ;restore (AF)
05070            RET                         ;return
05080 ;
05090 ;=== OUTBYT === Calls: $VDCHAR  (0033H)
05100 ;====================================*
05110 ;   Output a character to the screen   *
05120 ;====================================*
05130 ;
05140 OUTBYT   PUSH    DE                  ;save (DE) reg
05150            PUSH    IY                  ;save (IY) reg
05160            CALL    33H                 ;put the byte
05170            POP     IY                  ;restore (IY)
05180            POP     DE                  ;restore (DE)
05190            RET                         ;return
05200 ;
05210 ;=== LPRINT === Calls: LPBYTE
05220 ;====================================*
05230 ;   Output a line to the line printer   *
05240 ;====================================*
05250 ;
05260 LPRINT   PUSH    AF                  ;save (AF) reg
05270 LPNTX    LD      A,(HL)              ;get a byte
05280            OR      A                   ;byte zero?
05290            JR      Z,LAFINX            ;yes, jump
05300            CALL    LPBYTE              ;print the byte
05310            CP      0DH                 ;<CR>?
05320            JR      Z,LAFINX            ;yes, exit
05330            INC     HL                  ;bump pointer
05340            JR      LPNTX               ;go again
05350 LAFINX   POP     AF                  ;restore (AF)
05360            RET                         ;return
05370 ;
05380 ;=== LPBYTE === Calls: $PRCHAR  (003BH)
05390 ;====================================*
05400 ;   Output a byte to the line printer   *
05410 ;====================================*
05420 ;
05430 LPBYTE   EXX                         ;save the regs
05440            LD      HL,37E8H            ;status byte addr
05450            LD      B,A                 ;save character
05460 LP1      LD      A,(HL)              ;load status
05470            AND     0F0H                ;strip off status
05480            CP      30H                 ;printer on?
05490            JR      NZ,LP1              ;no, get stuck
05500            LD      A,B                 ;ret character
05510            CALL    003BH               ;send the byte
```

```
05520            EXX                    ;restore regs
05530            RET                    ;go back
05540 ;
05550 ;=== LPSTAT === Calls: nothing
05560 ;====================================*
05570 ;   Check line printer to see if ready  *
05580 ;====================================*
05590 ;
05600 LPSTAT   LD      A,(LPPORT)       ;get printer status byte
05610          AND     0FOH             ;strip off non-essentials
05620          CP      30H              ;is printer ready?
05630          RET     Z                ;yes, return to caller
05640          LD      HL,LPMSG         ;no, get message address
05650          CALL    PRINT            ;display message on screen
05660 STUCK    LD      A,(CTLKEY)       ;wait until <ENTER> pressed
05670          BIT     0,A              ;<ENTER> pressed?
05680          JR      Z,STUCK          ;no, wait some more
05690          JR      LPSTAT           ;yes, check status again
05700 LPMSG    DEFB    0DH
05710          DEFM    '>>==> Printer not ready, '
05720          DEFM    'press <ENTER> to continue.'
05730          DEFW    0DH
05740 ;
05750 ;=== INPUT (INPUTP) === Calls: INBYTE, OUTBYT, PRINT
05760 ;==========================================================*
05770 ;   Obtain line from the keyboard (with optional prompt)  *
05780 ;==========================================================*
05790 ;                                                         *
05800 ;       entry:  (HL) => input buffer                      *
05810 ;               (B)  =  max. number of characters         *
05820 ;     (INPUTP): (DE) => prompt message                    *
05830 ;                                                         *
05840 ;       exit:   (HL) - unchanged                          *
05850 ;               (DE) => last character                    *
05860 ;               (C)  - original contents of (B)           *
05870 ;               (B)  - actual # of characters             *
05880 ;               (A)  - termination character              *
05890 ;                                                         *
05900 ;       note:   The termination character is              *
05910 ;               included in the character string.         *
05920 ;                                                         *
05930 ;==========================================================*
05940 ;
05950 QUESTM   DEFM    '? '             ;question mark prompt
05960          NOP                      ;prompt terminator
05970 INPUTP   EX      DE,HL            ;shift prompt addr to HL
05980          CALL    PRINT            ;display prompt message
```

```
05990              EX       DE,HL           ;shift buffer address to HL
06000 INPUT        EX       DE,HL           ;shift buffer address to DE
06010              LD       HL,QUESTM       ;point to question mark
06020              CALL     PRINT           ;display question mark
06030              LD       C,B             ;move maximum chars to C
06040              XOR      A               ;initialize the byte counter
06050              LD       B,A             ; contained in register B
06060              PUSH     DE              ;shift buffer address
06070              POP      HL              ; to the HL register
06080              DEC      HL              ;  and decrement it
06090 ILOOP1       CALL     INBYTE          ;get character from keyboard
06100              LD       (SAVEAZ),A      ; and store for later use
06110              CP       08H             ;is character a backspace?
06120              JR       NZ,STOREZ       ;no, jump to store character
06130              LD       A,B             ;get character count
06140              OR       A               ;any characters yet?
06150              JR       Z,ILOOP1        ;no, skip backspace
06160              LD       A,' '           ;yes, load a blank
06170              LD       (HL),A          ; and store in buffer
06180              DEC      HL              ;backup one character
06190              DEC      B               ; and decrement char counter
06200              LD       A,8H            ;load backspace character
06210              CALL     OUTBYT          ; and display on screen
06220              JR       ILOOP1          ;look for more input
06230 STOREZ       CP       0DH             ;is character <ENTER>?
06240              JR       Z,INRETZ        ;yes, go to return
06250              CP       1H              ;is character <BREAK>?
06260              JR       Z,INRETZ        ;yes, go to return
06270              LD       A,C             ;get maximum length
06280              CP       B               ;do we have max. chars?
06290              JR       Z,ILOOP1        ;yes, ignore input
06300              LD       A,(SAVEAZ)      ;no, retrieve character
06310              INC      B               ;increment byte counter
06320              INC      HL              ; and address pointer
06330              LD       (HL),A          ;  and store character
06340              CALL     OUTBYT          ;   and display character
06350              JR       ILOOP1          ;go to look for more
06360 SAVEAZ       NOP                      ;character storage area
06370 INRETZ       INC      HL              ;increment pointer
06380              LD       (HL),A          ; and store last character
06390              EX       DE,HL           ;restore registers
06400              CALL     OUTBYT          ;display last character
06410              RET                      ;return to caller
06420 ;
06430 ;=== INBYTE === Calls: KBSCAN
06440 ;================================*
06450 ;  Wait for a keyboard character  *
06460 ;================================*
06470 ;
06480 INBYTE       CALL     KBSCAN          ;look for a byte
06490              OR       A               ;zero returned?
06500              RET      NZ              ;no, return
06510              JR       INBYTE          ;yes, look again
06520 ;
06530 ;=== KBSCAN === Calls: $KBCHAR (002BH)
06540 ;================================*
06550 ;  Scan for a keyboard character  *
06560 ;================================*
06570 ;
06580 KBSCAN       PUSH     DE              ;save (DE) reg
06590              PUSH     IY              ;save (IY) reg
```

```
06600              CALL    2BH              ;scan keyboard
06610              POP     IY               ;restore (IY)
06620              POP     DE               ;restore (DE)
06630              RET                      ;return
06640  ;
06650  ;=== CPHLDE === Calls: nothing === equivalent to RST 18H
06660  ;===================================================*
06670  ;   Compare the contents of registers DE and HL   *
06680  ;===================================================*
06690  ;
06700  CPHLDE    LD      A,H              ;get MSB number 1
06710            SUB     D                ;compare MSB's
06720            RET     NZ               ;return if not equal
06730            LD      A,L              ;get LSB number 1
06740            SUB     E                ;compare LSB's
06750            RET                      ;return to caller
06760  ;
06770  ;   ****************************************************
06780  ;   *  Data storage areas, buffers, and control blocks  *
06790  ;   ****************************************************
06800  ;
06810  LINES     NOP                      ;number of lines printed value
06820  QUOTE     NOP                      ;quote processing logical variable
06830  REMARK    NOP                      ;remark processing logical variable
06840  LL        DEFB    79               ;line length (less 6)
06850  PL        DEFB    60               ;page length (number of lines)
06860  PS        DEFB    6                ;page spacing (lines between pages)
06870  DCB       DEFS    32               ;File Control Block (FCB)
06880  BUFFER    DEFS    256              ;Input file buffer
06890  ENDBUF    EQU     $                ;end of input file buffer label
06900  OUTLYN    DEFS    512              ;secondary output buffer
06910  LINE      EQU     $                ;primary output buffer
06920  ENDPGM    EQU     $                ;end of program label
06930  CORE      EQU     ENDPGM-BASFMT    ;core used calculation
06940            END     BASFMT
```

End Listing

SOFTWARE REVIEWS

microSUB:MATH

Company: foehn consulting, P. O. Box 5123, Klamath Falls, OR 97601
Computer: Microsoft Fortran-80 and Cromemco Fortran IV (CP/M-80 only), Supersoft SSS Fortran (CP/M-80 and MSDOS), and IBM Fortran 77
Price: $250.00 ($25.00, manual only)
Circle Reader Service No. 133
Reviewed by Elbert F. Hinson

microSUB:MATH is a rich collection of Fortran numerical-methods subroutines for use in the scientific and engineering fields. These routines are meant for serious Fortran programmers. The subroutines in this package provide many of the mainframe computer Fortran library routines for use on a microcomputer.

The package I received was for Microsoft Fortran-80 on a SSSD, 8-inch CP/M 2.2 disk. The routines are also available in almost any popular 5¼-inch disk format and are available for SuperSoft Fortran and Cromemco Fortran. By the time you read this article, a version should be available for the IBM PC.

The software is organized into six general areas. Two libraries, one for single precision and another for double precision, are provided for each area. The six areas are functions, interpolation, integration, matrix/linear systems, polynomials/nonlinear equations, and differential equations. Table 1 (at right) contains a complete description of the contents of each single-precision library. The double-precision libraries contain the same routines but provide the accuracy of double precision. The double-precision library names and subroutine names are prefixed with a D. For example, the library DFUNC is for double-precision functions and function DSINH is the double-precision hyperbolic sine.

Documentation

The manual supplied has a table of contents, an index, and tabs for each chapter. The introductory section gives an overall view of the available routines and how to use them. The method of linking the library routines with your loader is covered in sufficient detail for anyone familiar with Fortran. This is not a tutorial, so some knowledge of Fortran is expected if you wish to use these library routines.

Part of the introductory section is devoted to building custom libraries. This details which of the microSUB:MATH routines call other routines in the library and in what order the routines must be loaded in a custom library. Each of the six library sections has an individual write-up of each routine in the library.

The subroutine documentation is clear and easy to use. The manual provides the Fortran calling sequence, an explanation of the input and output parameters, an explanation of the method used and a sample program for each subroutine in the library. There are appendices for quick reference and test program results.

A set of test programs is provided on the distribution disk that exercises every subroutine in every library. Results of the test programs are provided in the manual for single precision but not double precision.

Performance

I tested a significant number of the subroutines but did not have time to test

Library	Subroutine	Purpose
FUNC (Functions)	SINH, COSH, TANH ASINH, ACOSH, ATANH	Hyperbolic trigonometry
	CM, CD, CEXP, CLOG, CPWR	Complex math – multiply, divide, exponential, natural log, and raise to power
	CSIN, CCOS, CTAN	Complex trigonometry
	PR, RP	Polar to rectangular and rectangular to polar transform
	RJN, RJV	Bessel functions, integer and real orders
	ERF	Error function
	GAMMA	Gamma function
	SFRES, CFRES	Fresnel integrals, $S(x)$ and $C(x)$
	SI, CI	Sine and cosine integrals
	CNST	Constants – PI, Euler's, and computer accuracy
INTP (Interpolation)	LIN	Linear interpolation
	LAGR	Lagrange interpolation
	UDVTB, UDVIN	Table for divided difference and solution for unequally spaced points
	EDVTB, EDVIN	Table for divided difference and solution for equally spaced points
	SPLTB, SPLIN	Second derivative table for cubic spline and cubic spline interpolation
	LAGDB	Double Lagrange interpolation
INTG (Integration)	SIMP	Simpson's
	ASMP	Adaptive Simpson's
	ROMB	Romberg
	GAUS1	Low order Gauss-Legendre
	GAUS2	High order Gauss-Legendre
MATX (Matrix/ Linear Systems)	RLU, SLU	LU decomposition and solution by Gauss elimination
	SOER, SOES	Gauss elimination with scaling and partial pivoting and solution from elimination
	GJOR	Gauss Jordan elimination
	MTM	Matrix multiplication

Table 1.

every one of them. I compared the results of most of the functions for random values against the CRC math tables. Function results appeared good for six digits of precision. Double precision should provide more accuracy if it is needed.

I ran parallel tests of many of the Matrix, Integration and Polynomial subroutines against the programs in the book *Fortran Programs for Scientists and Engineers* by Alan R. Miller. The results were the same for at least five digits in the single-precision versions. Although I did not specifically benchmark the results, the routines appeared to execute at about the same speed. I found it very easy to substitute the subroutines from the libraries for the subroutines in Mr. Miller's book. It took an awful lot of my time and typing to build the same thing that is contained in the library.

Support

I talked to Mr. Hugh Currin of foehn consulting about the package and some of my wish list. The results of the double-precision test programs should be included in later releases, but this present omission should provide no real problems to the user. I asked Mr. Currin if he had considered including the transcendental functions tangent, arcsin, and arccos that Microsoft left out of their library. He indicated that they could be included in a later release.

Conclusions

I highly recommend this package for any scientist or engineer who would like the power of mainframe subroutine packages on a microcomputer. There is a significant amount of scientific work that can be done on a dedicated micro-computer instead of fighting a busy mainframe. This package is very easy to use and contains a very powerful set of numerical subroutines for microcomputers.

METAFILE Application Development System

The vendor's release describes META-FILE as follows:

"The METAFILE Software System creates a development and operational environment that includes integrated facilities for high-level programming, word processing, data base management, modeling, report generation, and communications."

This does indeed reflect the capabilities contained within METAFILE. The dictionary lists a variety of meanings for the prefix "meta," but clearly the one intended here is that of "beyond" or "transcending." And METAFILE certainly does extend the usual concepts associated with filing in its most general sense.

The potential users of this package could easily include a commercial enterprise requiring sophisticated management support, an applications programmer writing in a high-level language (META-FILE, in this case), and a stand-alone individual who likes to have a lot of capability at his or her fingertips. An example, which demonstrates both the program's integration and its power, is the ability to incorporate extensive numerical or logical calculations within a word processing environment and to store those calculations in either an evaluated or an unevaluated form so that the input can be changed, rather like a spreadsheet (though this is certainly not a spreadsheet), to obtain an output reflecting the change.

In order best to understand META-FILE's structure, we have characterized it as an operating system, called META-FILE ASSISTANT, that controls access to a series of integrated and interacting

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(Continued from previous page)

Library	Subroutine	Purpose
	MINV	Matrix inversion
	JCBI	Jacobi, eigenvalues and -vectors of real symmetric matrix
	HQR1	HQRI eigenvalues of a real symmetric matrix
	INVIT	Eigenvector of a real symmetric matrix
	HQR2	HQRI eigenvalues of a general real matrix
	HQR3	HGRI eigenvalues and -vectors of general real matrix
POLY (Polynomials/ Nonlinear Equations)		Real root of an equation:
	HINT	Half interval method
	BRENT	Brent's method
	NEWT	Newton's method
	NEWTS	Nonlinear system of equations (Newton's method)
	JACM	Jacobian matrix (used in NEWTS)
		Orthogonal Polynomials:
	CHB1, CHB2	Chebyshev, first and second kind
	HERM	Hermite
	RJACB	Jacobi
	RLAGR	Generalized Laguerre
	RLEGN	Legendre
		Roots of Polynomials:
	PLYR1	Second, third, and fourth order
	PLYR2	One polynomial root (Muller's method)
	POLY	Polynomial evaluation
	PLYR3	Roots of polynomial (Muller's method)
DIFEQ (Differential Equations)	RGK	Runge-Kutta, fourth order
	ADMLT	Adams-Moulton, fourth order
	RKF	Runge-Kutta-Fehlberg variable step size

Table 1.

programs. These programs include: Record Management; Text Management; Text Print Control; and METAFILE Command Language. These terms are from the manual. In more familar language the programs are, respectively: Relational Data Base; Full Screen Editor; Text Formatter; and METAFILE Interpreter.

METAFILE employs two distinct operational screens. One, which we will refer to as the COMMAND mode screen, can be considered as METAFILE's operating system screen, and the other, which we will call the EDIT mode screen, is that of the text editor.

COMMAND Mode

When you call METAFILE from DOS, the first screen presented prompts for a user ID (optional) and the date/time if not entered via a clock/calendar on boot to DOS; it then asks you to press Enter. The user ID is a METAFILE built-in function and could be easily adapted, using their programming language, to control operator file access. The next screen displays the COMMAND mode of META-FILE (Figure 1, below).

The screen is 80 columns wide by 24 rows deep. The failure to use the 25th row, available on the IBM PC, is probably related to the program's origins from the "Vector 3." These origins are reflected also in the way the function keys on the keyboard are ignored, using instead various Ctr1-letter combinations. These aspects show a weakness in design that

Sensor-Based Systems states will be remedied in their next release.

The 23rd row is used for status information, and the 24th row is used for command area entry and for messages. The 22nd row is held blank to allow separation between the display area and the status lines, thus leaving a usable display area of 21 lines. Both the status area and that containing the METAFILE ASSISTANT are in reverse video.

The METAFILE ASSISTANT controls access to a set of integrated programs. METAFILE is essentially a menu-driven program, choices being made by positioning the cursor on an option and pressing the Return key; depending on the option selected, further menus may be presented for additional selection. Some commands can be entered directly by typing the appropriate command, thus bypassing some intermediate menu structure.

Perhaps the two most-used keys with METAFILE are Esc and ~, the latter not requiring the Shift key. ~ is referred to as the ATTN key, and when pressed it presents in the display area the relevant menu for the program in operation or the METAFILE ASSISTANT menu. The Esc key is implemented in the same way as in VisiCalc, Lotus 1-2-3, and many other programs: it takes you back one level in your calling sequence.

CALCULATOR

Of the five principal options, CALCULATOR offers the most immediate

access. This option positions you directly in a screen insert, in reverse video and of the same contour as METAFILE ASSISTANT, from which you may perform a variety of numerical and logical calculations using the METAFILE Interpreter.

In the CALCULATOR mode, you can carry out any of the allowed METAFILE operations using numbers, strings, defined variables, and built-in functions for its operations. Calculations are carried out in floating decimal to 20 places. Included in the operations are add, subtract, multiply, divide, raise to power (**), operations on "dates," and the standard set of relational operations, for both numbers and strings, yielding a 1 if true or a 0 if false. The logical AND, OR conjunctions are included.

Built-in functions include natural logarithms, absolute value, random numbers, a variety of "date/time" functions, a function returning a one-byte binary equivalent of the value of its argument (this is used for printer control addressing), and a host of others related principally to data base structures. Arithmetic is performed strictly from left to right, and a generous use of ( ) is suggested to ensure that CALCULATOR does what is intended.

ACTIVE CONTEXT

The ACTIVE CONTEXT option offers a menu enabling selection from a listing of submenus. These summarize the active resources in METAFILE. They include: inter alia; the status of parameters established in the text editor, such as margins, search and translate values, page formats, printer status, etc.; structure of record files; listing of variables established with memory used; name of current drive; active files; and a host of other parameters defining the current status of the system. ACTIVE CONTEXT is presented in the same format as META-FILE ASSISTANT.

LOOKASIDE

The LOOKASIDE option is a particularly useful utility. Essentially, LOOKASIDE is a split-screen facility, but in METAFILE's realization, the alternate view is presented in the same reverse video setaside contour as the METAFILE ASSISTANT. Selecting LOOKASIDE presents you with a submenu of three choices: EDIT, UPDATE, and VIEW, as well as the option of typing in a command, as is offered in METAFILE ASSISTANT.

It is useful to think of LOOKASIDE as a second (albeit somewhat less enriched) command screen from which to exercise METAFILE's capabilities. Thus, selecting one of the options can lead to further menu choices, as from META-FILE ASSISTANT, or entering a command directly on LOOKASIDE's com-


```
                    Display Area


                        ///////////////////////
                        /                     /
                        / METAFILE ASSISTANT  /
                        /                     /
                        / Directory           /
                        / Lookaside           /
                        / Active Context      /
                        / Calculator          /
                        / Facilities          /
                        /                     /
                        / End of Session      /
                        /                     /
                        / Select an Option    /
                        / (or type a command) /
                        /                     /
                        ///////////////////////

    :Proc:    :A:Text      :Data: Inactive: Rec: Ø of Ø :

<---------------------- Command Mode Screen ---------------------->
```

Figure 1.

mand line can enable you to move more directly to the task.

Selecting the EDIT option allows you to select a text for editing. Choosing the VIEW option allows a passive interaction with a file or record (like DOS's "type"). Choosing the UPDATE option allows operations to be performed with the data base. Other options may be entered directly. In LOOKASIDE, you have the same full power of the utilities available as if they had been called from the COMMAND screen.

Used in conjunction with the regular EDIT mode, this option allows the equivalent of "copying" from the LOOKASIDE file into the current text buffer by utilizing the include facility of METAFILE's editor. This allows you to look at what you may want to include in LOOKASIDE, edit it, write it out to file, and then read it back in, into the active text buffer.

DIRECTORY

The DIRECTORY option is concerned both with the management of the storage devices and with the active text buffer. Since METAFILE considers certain file types as reserved (e.g., those dealing with procedures, descriptions, and forms) and distinct from text files, this option offers the viewer the choice of a file listing by drive and type (and text is considered a type) or, under General Directory, by drive without restriction. Optionally, the current text buffer can be chosen. Once a file is selected, the user is presented the choices of EDIT, VIEW, or PRINT. Finally, data is presented on storage space availability and entry utilization. Some of the submenus of DIRECTORY are also available from other menus.

FACILITIES

The FACILITIES menu offers the largest number of submenus: nine, plus the usual choice of typing in a command. Each of these options in turn references further submenus. If METAFILE has an operational core, it is contained here. Included is access to basic operating system commands dealing with file management, such as storing, activating, erasing, renaming, transferring from one storage device to another, and establishing communications with other METAFILE systems via an "on line" (local network). Access to the editor is also available here.

Of particular importance is entry to the large resource of commands dealing with the data base; these include commands relating to its definition, maintenance, and interrogation, as well as the structures relating to form preparation and report generation. An additional option available relates to the procedures

to "import" data base files from "foreign" formats into METAFILE.

EDIT Mode

The EDIT mode contains the (full) screen editor. The display area is 80 columns wide and 21 rows deep. Row 22 is blank, and rows 23 and 24 are used for status and prompts, respectively (row 25 is not utilized). The status line displays (in reverse video) the filename; line and column number of the cursor with respect to the upper left corner of the complete text; line and column number of the cursor with respect to the current screen (frame); margin status; mode status (wordwrap, insert or overtype, etc.); and percent of text space filled. Horizontal scrolling up to 250 columns occurs if wordwrap is not enabled at a column less than 80.

METAFILE's EDIT buffer is maintained entirely in memory and has a maximum capacity of about 63K. Due caution is required in movement between the EDIT and COMMAND modes so that the active EDIT file is not lost; METAFILE has inserted prompts to ensure that the buffer is not inadvertently lost.

The EDIT mode is used for a variety of purposes in addition to straight text editing. On one level it can be viewed as

the editor for the METAFILE command language. In EDIT you can write a METAFILE procedure (program) and run it (note the similarity to BASIC). The editor is also used to create forms for data entry and to format reports.

Access to the various editing commands is via the ATTN (~) key where a menu is presented. Selection can be made by positioning the cursor on the command and pressing Enter or by typing the first letter of the command. A help facility is available. The editor is crisp in its response to commands. In fact, the only slow response was that of its type-ahead buffer, principally when scrolling beyond column 80. In that case, it was easily fillable, to which it responds with the standard IBM "beep."

As a word processor, EDIT is quite good, comparing quite favorably with many stand-alone programs. It has some very good features that many do not have and lacks some that would make it even better. One of its nicer features is the automatic reformatting (if wordwrap is enabled) of text in the insertion mode. A negative feature is the lack of any overall reformatting capability. Once the margins are set, you cannot return, change them, and reformat. It has a reasonable set of cursor and text controls: move, copy,

delete line and text block, search and replace, and the ability to insert files into text, to print blocks of text, or to save them to disk.

Text Formatting

Extensive formatting commands allow headers, footers, centering, protection of text block on printing, multiple print copies, printing of selected pages, automatic page numbering and date inclusion, right justifying, and others. Practically any METAFILE command can be included in the print formatting, and printing can be done very easily from within METAFILE. The print control commands are included in the text (on separate lines) using *%command* at column 1, where *command* is any valid print command or any METAFILE command. Comments may be included by following the % symbol with a semicolon, followed by the comment.

Special Features

It is worth considering some of the unusual facilities of METAFILE. One of these is the ability to include operations within a text file. Any allowed META-FILE command or active variable can be accessed in the EDIT mode by placing the variable name or the METAFILE command within { }. The editor has a command named EVALUATE that displays the value of the variable or of the operation contained within { }. EVALUATE toggles between evaluated and unevaluated displays.

For printing or saving text to file, one can choose either the evaluated or the unevaluated mode. Since the variables defining record parameters are included here, it is clear that this offers an exceptionally rich tool in the generation of forms, reports, and letters, utilizing elements from the data base. It is this feature that makes mail merge and the printing of personalized letters extremely easy from METAFILE.

Consider a simple example: an array variable X(i), where each element is a name. A given name is referenced as {X(i)}. On calling EVALUATE, a name is displayed. For the case where the X(i) are numbers, if {X(2) * X(3) – X(4)} is entered in text, on calling the EVALUATE command, the substitutions are made and the formula evaluated.

As with many things, there is a price to pay. The METAFILE editor counts characters in the unevaluated mode. Thus, if {A} (containing on evaluation "Mr. Jones") is within the wordwrap margin in the unevaluated mode but extends beyond it on evaluation, wordwrap does not occur; on printing the string, "Mr. Jones" will extend beyond the wordwrap margin. However, the failure to correctly determine the character count on the unevaluated { } has more serious consequences for embedding printer controls in text, which may be required since META-FILE's current printer support is rather minimal.

The command to embed a control is {Code (#)}, where # is the decimal value of a number in the print control sequence. Here the embedded format dominates in the determination of the column number for wordwrap, and it will always occur before the evaluated location. You are also offered the option of implementing printer control codes by text formatting; the format is: "%CONTROL *formula*" where *formula* is as above without the { }. This will not affect wordwrap, but it does not allow in text embedded controls.

It does not require too much imagination to see that it is possible to produce formatted spreadsheets using the { }

utility. It is not a substitute for a true spreadsheet, since to change the value of a variable requires that the METAFILE "SET" command be used (equivalent to BASIC's LET VARX = xxx), but it does offer a very useful tool.

METAFILE's procedures for creating menus are particularly easy to implement. This is done by using [ ]'s to contain METAFILE commands or the names of executable METAFILE procedures as the options in the menu. An active text, with these choices enclosed within [ ]'s, is the basis for user-defined menus.

The menu and its options are formatted within the EDIT mode under a user-selectable filename with type MEN. A METAFILE procedure of the same filename but of type PRO is written (the manual lists the procedure; it is a mere seven short lines). The menu is called from a command line by typing in its filename. Under control of this program, the cursor keys control movement between options. Return selects an option, and when it is concluded the menu is presented again for further selection; pressing the Esc key concludes the menu session.

The options, as noted, can include METAFILE procedures so that the menu allows selection of user-written procedures to accomplish specified tasks. Since METAFILE allows a procedure to call others as subprocedures, it is clear that the language offers powerful capabilities for user definition and employment.

System Parameters

METAFILE has the implementation limitations listed in Table I (below) METAFILE admits three types of data elements: character (C), numeric (N), and data (D), with implementation restrictions noted in the table.

In addition to restrictions on implementation, another very important aspect of any data base is its operational performance in interrogating files. How long does it take to sort a file of 10,000 records on a single key? With qualifications? And so forth. Unfortunately, these questions are not addressed in the manual, and this reviewer has certainly not had the time to create such a base and test its performance. METAFILE, however, is not unique in failing to address these problems in its documentation.

It is not really meaningful to carry out a test on a small data base since it does not yield meaningful results for the potential user with a data base of significant proportions. This aspect poses a serious problem for any prospective user of any data base program, and it is not one that has been appropriately addressed. This is an area that calls for clarification for all such programs.

1.	Maximum number of characters in a filename	8
2.	Maximum number of characters in a file type	3
3.	Maximum length of a record (characters)	1000
4.	Maximum number of records in a file	32000
5.	Maximum size of a text (approx. pages)	15
6.	Maximum length of a text line (characters)	250
7.	Maximum length of an item (characters)	235
8.	Maximum precision of a number (digits)	20
9.	Double quote (") not allowed in record item.	
10.	Backslash (\) not allowed in text line.	
11.	File and item names may contain only alpha numeric characters and underscores (_).	
12.	Date must lie between 01/01/1776 and 12/31/9999	

Table 1.

METAFILE Restrictions

Installation

The package contains three Reference Guides, an Installation Guide, a "Configuration Device," and two double-sided, double-density disks. The program disk has 317,440 bytes of operational files, and the demo disk has 284,672 bytes of files. The program disk is so packed that if you are operating under DOS 1.1 you cannot get any DOS system files on it. Under DOS 2.0 you can just barely make a bootable disk with the system files, but when you go to enter your printer configuration (even as small a file as that), you obtain a "disk full" message.

Thus, for practical purposes a separate disk to boot the system is required, and you can then include such useful utilities as a clock calendar (whose input METAFILE automatically accepts); ProKey, which works quite well; and a RAM disk, which is quite useful. In exploring METAFILE, I used a boot disk formatted as indicated above, with an autoexec.bat file that, among other things, copied the program files in drive b: to the RAM disk (after a prompt to install in drive b: a disk with program files). I then put a data disk in drive b: (again on prompt) and was set to go.

METAFILE is copy protected via hardware. The disks per se are not protected, but to utilize them effectively you must connect the METAFILE "Configuration Device" to the COM 1 serial port. This device is a short cable having female and male connectors with a "black box" in between. The connector choices allow you to reconnect your modem (or any other device that may have been using the port) onto the cable and operate normally.

I tried my modem in that configuration and it worked without problems. Without the special cable connection, METAFILE can function but with only minimal memory (about 12 percent of system capability) and as such is useful only for demonstration. The device operates to restrict the system disk or its copies to a single machine for full utilization. There are the usual warranties to replace either the system disk or the "Configuration Device" within 90 days if either malfunctions.

Since METAFILE will operate with 128K RAM and the program files take up almost 320K, it is clear that many calls are going to be made to the program disk. If you do not have a hard disk, operating from a floppy can be rather slow; you could, however, use a RAM disk, which will afford some increase of speed even with a hard disk. METAFILE's instructions must be followed very closely for operation with a RAM disk. To operate with RAM, METAFILE is called from DOS with the command: metafile , memory:*xxx*K, where *xxx* is the amount of memory being reserved for METAFILE in units of 1000.

Since METAFILE's capabilities and capacities are affected by the amount of memory available (delineated for the user in a table in the Installation Guide), reasonable care should be exercised in utilizing a RAM disk. Since at least 320K is needed for the RAM disk, a user contemplating that approach should, by and large, have a totally configured system with respect to memory. Of course if you operate from a hard disk, you have more options available.

In any event, METAFILE demonstrates some peculiarities with respect to RAM disks. I use two different programs. One utilizes the RAM disk driver contained in DOS 2.0, modified for higher capacity, and the other is Tall Trees' JFORMAT. METAFILE seems to accept the former with great tolerance, behaving well even when I forget to call it as indicated above. On the other hand, with JFORMAT, unless it is called in precisely the specified format, the system hangs up; if called correctly, it performs flawlessly.

There is one operational oddity, at least in the system I have, that is independent of the use of RAM disks or of any other peripheral hardware or software. After once calling METAFILE and closing a session to return to DOS, the DOS underlined blinking cursor is lost and does not return until the system is rebooted (warm does it OK); otherwise, DOS performs normally. I must admit that loss is a little disconcerting, but it is not disfunctional.
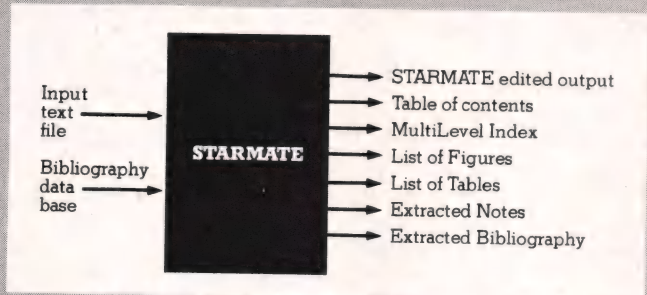
Documentation

The Installation Guide tells you how to use the demonstration facility (the demo disk). I find demos hard to evalu-

ate. It is my belief the reasonably educated user is probably going to have more difficulty than a novice, principally because he or she just does not follow them with sufficient attention. On occasion I have asked my children or my wife to react to a demo disk of some program, and they have invariably been more satisfied with it than I was.

With that caveat, I was not overly impressed with the demo disk. It did introduce you to the various aspects of what the system can do, asked you to create some records, to enter some text, to correct it, to print it, to do some finds and replaces, to sort a file, etc. But I do not really believe that someone who does not have some understanding of what files, records, and fields mean can really use the demo and get very much out of it.

METAFILE is supported by three separate reference documents: spiral bound, 8.5 x 11 inches in size, with heavy paper covers.

Volumes 1 and 2 are labelled as User's Guides. Volume 1 is dedicated to A. Word/Text Processing and B. Record Management; Volume 2 to C. Introduction to Programming, D. Applications Design, and E. Glossary/Index. Volume 3, about twice the thickness of either 1 or 2, is labelled Reference Manual. These manuals have many faults. One of the most nettlesome is that they are "dirty." I found after using them even briefly that my work area was full of lint generated by the manuals and I was constantly vacuuming it.

The manuals are by far the weakest part of the package and do not, in fact, do it justice. Operationally, they are not organized for ready reference: there are no index tabs to allow quick selection of text sections, and while there are tables of contents for each of the major subject areas, the indexes are associated only with the glossary and with the Reference Manual. The latter two indexes are reasonably well presented, each organized alphabetically with a description and a page reference for each item, but there is a fair degree of redundancy as well as missing items. Included with the reference documents are a METAFILE User Reference Card (8.5 x 3.5-inch folder) and a Function Key Specification card of the same dimension.

It is very cumbersome, particularly when finding your way through a new and rather complicated package, to have to shuffle three separate documents, none particularly well designed for reference. The vendors do indicate that new documentation and a new format (the IBM notebook and slipcase style) is in preparation for future release. It is sorely needed.

The first two User's Guides are hands-on tutorials. Volume 1 takes you through a subset of the capabilities of the word processor and includes a tutorial on record management using a hypothetical personnel office as a case study. Volume 2 includes a tutorial on the programming facilities of METAFILE, again using a case study based on a set of employee records, and a final set of case studies dealing with the use of advanced commands and facilities, application guidelines, prototyping techniques, testing, debugging, and documentation.

The last three tutorials contain many useful examples and METAFILE programs that can be employed by the user, demonstrating such diverse aspects as mail merge, label addressing, form letter preparation, report documentation, and "converting" documents to METAFILE from foreign formats. The conversion presented is for record formats only, and I was unable to find any discussion dealing with the "importation" of text files. Some experimentation indicated that the METAFILE editor will accept only the purest of ASCII text files. Since a reformat capability is not included, the ability to import text files is relatively moot.

Displaying the hierarchy of a command structure, particularly with a program so richly structured as METAFILE, in a readily comprehensible fashion is important. This is an area where the documentation is also quite weak. Although the structure is presented, it is not done systematically nor in a way that gives an overall perspective; as a result, it is not useful for orientation or reference. The user is unable to discern the relationships between various units without an undue and uncalled for effort on his or her part. A tree diagram would have been most useful.

Printer Support

A few comments on printer support are in order. The documentation states that the IBM Epson is supported and it means just that: the bare-bones Epson without Graftrax. This is not made very clear, and certain printer commands indicated in the manuals are not realizable with any Epson other than the bare-bones IBM. It turns out that it is possible to embed printer controls in text, as well as in other formats, due to the enormous capabilities of the program, but it is almost impossible to discern from the manuals how to do this. I had to resort to the vendor for information.

Sensor-Based Systems is very responsive to requests for assistance and was able to head me in the right direction. As a result, I found it possible to utilize the full text print capabilities of the Epson MX 100 (and, I would judge, any other printer) since decimal sequences can be sent to the printer from within text files (with care, as noted above) or from formatting commands. Printer support is certainly one of the weakest operational aspects of the program.

Any serious user of METAFILE in substantive proportions will almost certainly need print spool capabilities. A print spooler is not included, and given the memory requirements of METAFILE, the most efficient way to obtain that would be with a stand-alone spooler.

Use of the Keyboard

Although a card indicates the META-FILE function key specifications, the word "function" is used operationally. The IBM function keypad is not used at all. In addition to a reasonably standard use of the cursor pad, 16 METAFILE-dedicated commands are effected by Ctrl-letter combinations. Included in the 16 are two cursor-control commands.

Some of the letters in the Ctrl-letter combinations have a relation to the command but many do not; it seems to me a significant improvement would be to assign to the function keys their related combinations. In fact, I reassigned them with ProKey so that they coincided, insofar as possible, with similar assignments I use on other programs (one of the great advantages of products like ProKey is cross-program standardization). When I commented to the vendors on their failure to use the function keys, their response indicated that future releases would include a built-in macro facility that would allow the user that option, among others. This would be an enormous enhancement indeed.

METAFILE Programming Language

METAFILE's programming language encompasses a wide range of capabilities. It is primarily a nonprocedural language, although it contains a significant number of procedural commands. In the latter respect, it is structurally similar to BASIC, though the COMMAND names are different. In METAFILE, the command WRITE 5*45 will yield 225, just as in BASIC; PRINT 5*45 will do the same.

METAFILE procedures (programs) are identified by the file type (extension) PRO, just as BASIC programs are identified by the extension BAS, and as in interpreted BASIC, they can be run from the editor. All METAFILE commands must be entered in upper case, which given the IBM's lack of a CAPS LOCK indicator is somewhat of a nuisance.

An appreciation of METAFILE's capability can only be obtained by working with it, but it seems reasonable at least to list the classes of commands available and their functions (see Table II, page 83). Brief explanatory comments are noted in some cases.

Variables can be named and multi-dimensional subscripting is allowed as long as the total length of a data element

is less than 235 characters. Variables are initialized by the command SET (with multiple initialization on one line allowed) and undefined with the command UNSET. Output is routed to SCREEN, PRINTER, or TEXT by use of the DEST command; default is SCREEN.

It seems useful to list the memory-dependent capacities of METAFILE as given in the Guide to Operations, since system performance is dependent on it. See Table III (page 83).

Summary

METAFILE is an excellent data base manager (record management, in META-FILE's lexicon) that utilizes a relational format. Each data record file has a data description specifying the attributes of the file and the items in its records. These are submitted in response to a command called DESCRIBE. Each item is capable of quite extensive description, including a comment line. Five files may be activated simultaneously.

Records may be updated, descriptions changed, new records inserted, and old ones deleted. Changes may be permanent or temporary, and the files may be searched and sorted with flexible querying commands. Fast lookup is implemented by a KEYING command. Columnar reports are produced with a LIST command and a row format with a ROW command. ITEMIZE will produce qualified records having non-null values.

The FORMAT commands allow extensive text management and preparation facilities, and flexible form and report generation is quite accessible. The ability to write METAFILE procedures to incorporate user-defined command sequences makes it possible to produce a combination of menus and procedures to permit system use by relatively inexperienced operators.

METAFILE goes well beyond data base management and word processing. The support from its programming language and the tight integration of the system offer the user enormous flexibility and suggest many interesting uses beyond those usually encompassed within those terms. If METAFILE gains acceptance in the user community, I would expect to see significant offerings of applications programs utilizing its facilities.

I am sure we have all heard the statement from high-level management: "I want to be able to press a button and see the status of XYZ." This is literally realizable with METAFILE. It is very easy to envisage writing a procedure with minimal prompts (whose implementation META-FILE makes easy) to load the appropriate file(s) and display the appropriate variables. Alternatively, the display could be done from the editor, exhibiting data variables within { } whose evaluation is

---

PRINT CONTROL COMMANDS — 14

SESSION FACILITIES — 8
(start, stop, HELP, print, etc.)

PROCEDURAL CONTROL

Clauses	8 commands
Commands	9 commands

FORMULA FACILITIES

Calculations	6 commands
Conjunctions	2 commands
Formatting	7 commands
Character Manipulation	5 commands
Built-In Functions	29 commands

All of the above are executable from a METAFILE procedure (program) or, as appropriate, from a command line or within the editor.

TEXT PREPARATION FACILITIES — 30
(in the EDIT mode only)

FILE COMMANDS

Record Files	9 commands
Text Files	6 commands
All Files	4 commands (erase, rename, transfer, connect)

RECORD COMMANDS

Update	8 commands
Qualification	3 commands (criteria)
Navigation	7 commands (moving within files)

TEXT LINE COMMANDS

Update	4 commands
Navigation	8 commands (locations)

DATA ELEMENT COMMANDS — 3

FRAME COMMANDS — 9
(these relate to the generation of forms and reports, including user interaction)

REPORTING COMMANDS — 18

Table 2.
METAFILE Commands

Total Memory[1]	Text Area	Separate Text/Sort	Sort Area[2]	Variables[3]	Items[4]
126K	26K	no[5]	26K	2175	3859
192K	63K	no[5]	63K	2943	4352
256K	63K	yes	63K	2943	4352
>256K[6]	63K	yes	63K	4100	6400

[1] Total memory means that memory assigned to METAFILE and does not include any assigned to a RAM disk.

[2] ((KeyLength)+3) * 10 (#Records to be sorted) must be less than 10 times the sort area (in bytes).

[3] Total number of characters in variables' names, values, and control information must not exceed value in table.

[4] Within active data file descriptions, the length of all items' names plus 10 characters per name must not exceed number in the table.

[5] These systems share the same area for text and sort. If you edit a text then sort a data file, text just edited is cleared from memory and must be recalled for re-edit.

[6] >256K means 320K to 640K in steps of 64K.

Table 3.
METAFILE Memory-Dependent Capacities

accessed by hitting the Evaluation key. METAFILE allows scrolling (using the Tab key) through the variable sequence. The existence of a macro facility would enhance this even further.

LOOKASIDE makes a user's "calendar" readily available for viewing or for updating without leaving the current work session. The enterprising programmer will, I am reasonably sure, come up with additional applications. The only aspect really lacking is a graphics capability; the lack of this capability is certainly not an oversight since METAFILE is records and text oriented, not number oriented, though it does more than a creditable job in that area.

The integration achieved in METAFILE and the ease both of utilizing the facilities in the editor and of transferring data, accompanied by the capabilities of the ASSISTANT and the LOOKASIDE facility, certainly make this a strong competitor in the integrated program sweepstakes. All this without the need of a hard disk, a mouse, or an enormous memory capability.

METAFILE's manners are impeccable. The prompts are to the point and understandable. Its error handling facilities are clean and clear as to correction. As noted, the only hang up that occurred was when basic operational instructions were not faithfully followed. The only real quirk was that relating to the disappearance of the DOS cursor after leaving a session.

The more I have used METAFILE in the process of preparing this review, the more I have been impressed both by its capabilities and by its depth. Alas, I have also been more and more depressed by the poor documentation and the failure of the implementation to utilize the IBM PC screen and keyboard features and to respond to its weaknesses (here I am referring to the lack of any flags for CAPS LOCK or NUM LOCK status and the concomitant requirement of COMMAND entry in upper case).

By using the term "depth," I imply that the user will continually find the program opening up new opportunities for application and use. When I initially ascertained the price of the package and of the major updates, I was somewhat taken aback. After exploring its potential, and considering the price of other DBM systems currently in the market place, I am not as taken aback as I was. I would expect, however, that a program as costly as this would not be as lax in dealing with details of screen and keyboard usage nor exhibit the deficiency of printer support. The question of performance speed in the crucial area of data base application is just not answerable in the context of a single review.

This is not, in my opinion, a program for the novice, publicity claims notwithstanding. It may be true that the novice could do a few trivial things right off, so to speak, but to really utilize the capabilities of this program, users must take the time and effort to learn something about it. I believe they will be rewarded for that effort. To be able to effectively utilize the program's resources will require some programming skills, again claims to the contrary notwithstanding.

What do I conclude? If the new releases contain a significant improvement in documentation, more responsiveness to the IBM's screen and keyboard features (with due respect for strengths and weaknesses therein), and enhanced printer support, the METAFILE must be seriously considered by those requiring an effective data base and management support system.

DDJ

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16-BIT SOFTWARE TOOLBOX

by Ray Duncan

More on MS-DOS EXEC Function

Morton F. Kaplon, of Pomona, NY, writes: "I read with some interest your article in the December 1983 issue of *DDJ* . . . . I certainly do agree with you that the documentation on the DOS function 4BH is less than transparent, though my prize for opacity in DOS is still easily won by Chapter 13, 'Using Extended Screen and Keyboard Control.'

"In any event, my approach was somewhat different than yours and the program I had created offers a somewhat more extensive demonstration of the utility of that call. It allows the user to enter from the program anything which can be entered from the command line, albeit in two steps if there is a tail to the command line.

"Given this program, I believe it is fairly easy to see how one can construct a 'DOS MANAGER' for the user who wants to be prompted through DOS." Morton was kind enough to provide his program on a diskette, and we are printing a slightly modified version of it in Listing One (page 88).

PC-DOS Manual Typo

PC-DOS 2.0's manual contains a small typographical error in Chapter 13 which may confuse programmers who are attempting to use the Extended Screen Control. The command for "Erase in Line" (usually known as Clear to End of Line) is "ESC [K", not "ESC [k" as given in the text.

IBM PC Junior

We have now received a PC Junior to play with and are making many interesting discoveries. The graphics support on the Junior is considerably enhanced over the standard IBM PC, with programmable palettes and the ability to display sixteen colors in medium-resolution graphics mode or four colors in high-resolution mode.

One issue that has been completely ignored in the popular magazines is the raw performance of the Junior. Apparently, everyone assumes that because it uses a 8088 microprocessor at 4.77 MHz, it will execute programs at a speed comparable to a normal PC. Well, friends, it just ain't so. just ain't so.

The PC Junior does not have dual ported memory for the video refresh buf-fer. Consequently, the video controller must "steal cycles" from the CPU while it is updating the screen. It turns out that two out of three memory cycles for the main machine RAM are allocated to the 6845 controller, and only one out of three for the 8088. Running several standard benchmarks, I found that the execution speed of a computation-bound program on the Junior is exactly 45% of its speed on a normal PC. Programs which perform a great deal of screen I/O appear to suffer even more, though I haven't measured this exactly yet.

The literature on the Junior leads one to believe that programs executing out of a ROM cartridge will execute much faster, since the video controller will not be accessing that memory. I'll check this out and provide more information next month.

The performance of PC-DOS 2.1, which is provided with the Junior, has also been degraded. Apparently the disk drives provided on the Junior have much longer seek and head settle times, so the BIOS disk driver's delay factors have been lengthened accordingly. Access to files by a disk-based program is painfully slow — just the process of opening a file seems to take over a second.

MS-DOS Volume Labels

I had surmised from reading the MS-DOS 2.0 manual that I could change the volume label on a disk by simply performing a search with an extended file-control-block for the first directory entry with a "volume" attribute byte, requesting a "delete" on any filename found in that manner, then requesting a "create" function with the new name. However, this path is fraught with peril. The delete appears to work OK and returns a successful status code, but it clobbers the file allocation table unmercifully (see Figures 1 and 2, page 87).

I was unable to figure out exactly what was going on here. Examining a dump of the FAT and directory sectors, it is readily seen that a volume label is simply a normal appearing directory entry with the attribute byte set to 8, date and time fields filled in, and the "starting cluster" and "file size" fields zeroed. Except for the attribute byte, this is just the same as a directory entry for a file that was created but never had any data written into it; however, a delete of such an empty file works fine. Evidently, finding the volume attribute bit turned on during a delete function sidetracks PC-DOS into some nutty logic or other.

By experimentation, I found that the following sequence will safely add or change a volume label under PC-DOS or MS-DOS 2.0:

1. Set the disk transfer address to a scratch buffer 128 bytes long.

2. Using an extended file control block with the format shown in Listing Two (page 92), perform a "search for first" function (#17). If the register AL is returned as 0FFH, then the disk has no label — go to step 5.

3. If function 17 returned register AL = zero, the buffer contains a simulated extended file control block with the volume name in bytes 8-19. Move your new volume name (11 characters) to buffer + 24.

4. Passing the address of the scratch buffer in registers DS:DX, request a "rename" function (#23). If AL returns 00, the volume name was successfully modified and you are finished, otherwise something is horribly wrong with your system.

5. (Come here if the disk has no previous label.) Replace the wild-card characters shown in the extended file control block with the desired 11-character volume name. Passing the address of the extended fcb in DS:DX, request a PC-DOS "create" function (#22). AL returned as 0 implies success; AL returned as 0FFH means failure (usually due to a full directory).

I have provided an extract of the final working program as Listing Two (page 92) to get you started. Interestingly, this procedure *always* adds the volume label to the root directory no matter what your "current" subdirectory is.

I made another surprising discovery. Even though volume names were added in MS-DOS 2.0, the new version 2.0 extended file and record handling functions don't support adding or changing the volume labels. When I first tried this, I thought that I could probably use the new

CREAT call (3CH) and then use function 43H (CHMOD) to set the attribute on the new file to "volume"; however, the CHMOD call always returns an "access denied"error code.

Finding Size of TPA under MS-DOS 2.0

On the IBM PC, the Wang Professional Computer and presumably most other MS-DOS machines, there is a special BIOS call or some other method to get the amount of machine memory. Although it would be nice if RAM size could be obtained in a machine-independent way, there is no obvious MS-DOS function call that will do the trick. However, it can be done, although circuitously!

The method hinges on the "Modify Memory Allocation" call, MS-DOS function 4AH. It is defined as follows in the MS-DOS manual: "on entry, ES contains the segment of the block, BX contains the new requested block size in paragraphs. DOS will attempt to 'shrink' or 'grow' the specified block. If the call fails on a grow request, then on return BX contains the maximum block size possible . . . ."

Knowing that part of the MS-DOS procedure for loading an application program includes allocating a memory block to it, we can simply request an increase of our program's area to the maximum possible, inspect the amount of RAM that is actually given to us, and then add that to the segment address our program is based at in order to find the actual size of machine memory. See Listing Three (page 93) for an example subroutine that returns RAM size in Kbytes.

CP/M Discrepancies

While porting a program that performed ramdom file I/O from CP/M-86 to CP/M-68K, I discovered that they are not quite symmetric with respect to file control block structure. Specifically, bytes 31-34 (r0, r1, r2) have different meanings under the two operating systems (see Figure 3, at right). This was obviously done so that the lower 16 bits of the random record number would fall on a word boundary for the 68000, since 16-bit accesses on off addresses are verboten. However, it would have been nice if DRI had printed a warning about this in the manual somewhere.

■■J

(Listings begin on page 88)

```
Drive:B    Sector:1        Cluster:FAT#1    Dir:ROOT.DIR


           0  1  2  3  4  5  6  7  8  9  A  B  C  D  E  F

000000:   FC FF FF 03 40 00 05 60 00 07 80 00 09 A0 00 0B
000010:   C0 00 0D E0 00 0F 00 01 11 20 01 13 40 01 15 60
000020:   01 17 80 01 19 A0 01 1B C0 01 1D E0 01 1F 00 02
000030:   21 20 02 23 40 02 25 60 02 27 80 02 29 A0 02 2B
000040:   F0 FF 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000050:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000060:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000070:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

Drive:B    Sector:5        Cluster:        Dir:ROOT.DIR


           0  1  2  3  4  5  6  7  8  9  A  B  C  D  E  F

000000:   57 4F 52 4B 20 20 20 20 20 20 20 08 00 00 00 00
000010:   00 00 00 00 00 00 6F 14 32 08 00 00 00 00 00 00
000020:   46 4F 52 54 48 20 20 20 43 4F 4D 20 00 00 00 00
000030:   00 00 00 00 00 00 D9 08 6F 07 02 00 04 53 00 00
000040:   00 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6
000050:   F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6
000060:   00 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6
000070:   F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6 F6
```

Figure 1.

First block of File Allocation Table and Disk Directory before Delete performed on volume label using extended file control block. This disk contains only two directory entries and one file.

```
Drive:B    Sector:1        Cluster:FAT#1    Dir:ROOT.DIR


           0  1  2  3  4  5  6  7  8  9  A  B  C  D  E  F

000000:   FC FF FF 03 40 00 05 60 00 07 80 00 00 00 00 00
000010:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000020:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000030:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000040:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000050:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000060:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
000070:   00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
```

Figure 2.

First block of File Allocation Table after Delete function call performed on volume label. Note that many of the clusters allocated to the normal file on the disk are now missing.

Byte name	Offset	CP/M-86	CP/M-68K
r0	33	1sb *	msb+
r1	34	2sb	2sb
r2	35	msb	1sb

*1sb = least significant byte
+msb = most significant byte

Figure 3.

Discrepancy in file control block structure between CP/M-86 and CP/M-68K.

Listing One

```
         name    docom
         title   'DOCOM - Load and execute programs'
         page    55,132
;
; DO-COM.ASM --- Load and execute COM programs
; demonstrates use of PC-DOS function 4BH
; Morton F. Kaplan, 11 White Birch Drive, Pomona, NY 10970
;
; Definition of Function 4BH call:
; DS:DX points to ASCIIZ filespec
; ES:BX points to parameter block for load
; AL = function type:
;         0 is load and execute. Program must save registers
;             and must restore SS and SP.
;         3 is load only (may be used for overlays)
; Parameter Block
; AL= 0: Word segment address of environment string to be passed,0 value
;             for address will cause parent's environment to be inherited
;         DWORD pointer to command line to be placed at PSP+80h
;         DWORD pointer to default FCB to be passed at PSP+5ch
;         DWORD pointer to default FCB to be passed at PSP+6ch
;
; AL= 3: WORD segment address where file will be loaded
;         WORD relocation factor to be applied to the image.
;
;
; In this program, the filespec has to contain the filename and extension,
; a path may also be included.  Thus, to load and execute the program
; testmac.com on b: from drive a:, do as follows:
;
; A>docom
; Enter filespec to be loaded and run          (this prompt appears on screen)
; >b:testmac.com<enter>
; Enter command tail, if none then <enter>     (this prompt appears on screen)
; >
; Program then executes.
;
;
; Build this program from file DOCOM.ASM as follows:
; MASM DOCOM,DOCOM,DOCOM,DOCOM
; LINK DOCOM,DOCOM,,,
; EXE2BIN DOCOM.EXE DOCOM.COM
; ERASE DOCOM.EXE
;
;
; define necessary MACROs

dos      macro   fn                      ;fn is dos function to be called
         mov     ah,fn                   ;fn passed to ah
         int     21h                     ;dos interrupt
         endm                            ;end of macro
kb_in    macro                           ;chr typed at KB returned in al
         dos     1                       ;function 1
         endm
```

```
crlf      macro
          mov       dl,10                          ;line feed
          dos       2
          mov       dl,13                          ;carriage return
          dos       2
          endm

cue       macro
          crlf                                     ;perform CR-LF sequence
          mov       dl,'>'                         ;then display CUE mark
          dos       2
          endm

bfr_kbin  macro     bfr_addr                       ;buffered keyboard input
          mov       dx,offset bfr_addr             ;dx points to input buffer
          dos       0ah                            ;dos function call
          endm                                     ;end of macro

print_string macro bfr_locn                        ;print string at bfr_locn
          mov       dx,offset bfr_locn             ;point dx
          dos       9                              ;function call 9 of dos
          endm                                     ;end of macro

; END of Macro definitions

          codeseg segment para
```

(Continued on next page)

16-Bit Toolbox (Listing Continued, text begins on page 86)
Listing One

```
        assume  cs:codeseg,ds:codeseg,es:codeseg

        org     100h                            ;required for COM

start:  jmp     docom                           ;jump over data section


error1  db      'Memory Shrink Failure','$'
error2  db      'Error in Loading Program','$'
prompt_0 db     'Enter filespec to be loaded and run ','$'
prompt_1 db     'Enter command tail, if none then <enter> ','$'

stk_ptr dw      0                               ;save contents of SP
stk_seg dw      0                               ;save contents of SS
                                                ;PARAMETER BLOCK
                                                ;
parm_blk dw     0                               ;segment address of environment
                                                ;
                                                ;string passed,use parent
                                                ;pointer to command line at
        dw      offset psp80+1                  ; at psp+80h, +1 needed
        dw      ?                               ;to allow for structure of
                                                ;psp80 using buffered
                                                ;input from keyboard
                                                ;
        dw      offset fcb1                     ;pointer to default FCB1 at psp+5ch
        dw      ?
        dw      offset fcb2                     ;pointer to default FCB2 at psp+6ch
        dw      ?


                                                ;DATA FOR PARAMETER BLOCK
                                                ;
                                                ;buffer of 20 char for buffered
                                                ;line input, can be increased
filespec db     20,?,20 dup (?)
psp80   db      32,?,32 dup (?)                 ;32 chars in buffer,increasable
fcb1    db      1                               ;drive id
        db      11 dup ('?')                    ;filename and extension
        db      25 dup (0)                      ;rest of fcb
fcb2    db      1                               ;drive id
        db      11 dup ('?')                    ;filename and extension
        db      25 dup(0)                       ;rest of fcb


docom:                                          ;start of effective code
                                                ;initialize DS and ES
        push    cs
        pop     ds
        push    cs
        pop     es                              ;initialization finished
        mov     bx,50                           ;save 50*16 bytes for program
        dos     04ah                            ;es already set to code segment
```

```
                    crlf
                    jc          not_set             ;check for memory shrink
                    print_string prompt_0           ;display prompt
                    cue
                    jmp         over_erl

not_set:                                            ;display error message
                    print_string error1
                    mov         dx,ax               ;print error code in ax
                    add         dl,48               ;convert to ascii
                    dos         2                   ;display on kb
                    jmp         quit

over_erl:
                    bfr_kbin filespec               ;get filespec
                    crlf
                    sub         ax,ax               ;clear to 0
                    mov         al,filespec+1       ;get length of string input
                    mov         di,offset filespec+2
                    add         di,ax               ;to end of string
                    mov         byte ptr [di],0     ;terminate with 0

                    print_string  prompt_1          ;prompt for command tail
                    cue
                    bfr_kbin psp80                  ;get command tail
                    crlf
                    sub         ax,ax               ;clear to 0
```

(Continued on next page)

Listing One

```
                mov      al,psp80+1          ;get length of string input
                mov      di,offset psp80+2
                add      di,ax               ;to end of string
                                             ;terminate with carriage return
                                             ;and null byte
                mov      byte ptr [di],13
                mov      byte ptr [di+1],0
                mov      cs:stk_ptr,sp       ;save SS and SP in
                mov      cs:stk_seg,ss       ;variables addressable by CS
                mov      di,offset parm_blk
                mov      [di+4],es           ;seg address to psp80
                mov      [di+8],es           ;seg address to fcbl
                mov      [di+12],es          ;seg address to fcb2
                                             ;point to parameter block
                mov      bx,offset parm_blk
                                             ;point to ASCIIZ filename
                mov      dx,offset filespec+2
                mov      al,0                ;load and execute
                dos      4bh                 ;dos function to execute program
                mov      ss,cs:stk_seg       ;restore SS:SP pair
                mov      sp,cs:stk_ptr
                jnc      quit                ;jump if no load error
                print_string error2          ;print error message
quit:           dos      4ch                 ;back to dos,
                                             ;easiest way to terminate, can also
                                             ;pass code back with this call

codeseg         ends
                end      start
```

End Listing One

Listing Two

```
                .
                .
                .
        mov dx,offset buffer
        mov ah,26               ;set disk transfer address
        int 21h                 ;to scratch area (DS:DX) for
                                ;dictionary search.
        mov dx,offset xfcb      ;DS:DX=addr extended fcb
        mov ah,17               ;"search for first match"
        int 21h
        cmp al,0ffh             ;successful?
        je  no_label            ;no label on disk, jump.
        jmp label_found         ;label found, jump.
                .
                .
                .
xfcb    db  0ffh                ;flag signifying extended fcb
        db  5 dup (0)           ;reserved (should be zeroes)
        db  8                   ;volume attribute byte
        db  0                   ;drive code (set by program)
        db  11 dup ('?')        ;wild card filename & ext.
        db  25 dup (0)          ;remainder of fcb (not used)

buffer  db  128 dup (?)         ;buffer for directory search
```

End Listing Two

Listing Three

```
;
; Return memory size (Kbytes) in AX.
; Calling arguments:  none.
: Destroys AX, BX, CX.
; This routine assumes that ES is pointing
; to the Program Segment Prefix (true
; on original entry from MSDOS for either
; an EXE or a COM type program)
;
ramsize     proc        near
            mov         bx,-1       ;request 65535 paragraphs!
            mov         ah,4ah      ;MSDOS modify block function
            int         21h         ;returns paragraphs available
                                    ;in register BX
            mov         ax,cs       ;add that to the segment of
            add         ax,bx       ;the base of our program
            mov         cx,6        ;then shift paragraphs right
            shr         ax,cx       ;six places to get Kbytes
            ret                     ;back to caller
ramsize     endp
;
```

End Listing

C/UNIX PROGRAMMER'S NOTEBOOK

by Anthony Skjellum

This month's Notebook includes reader response to my December 1983 column, where I discussed perceived deficiencies in the Unix operating system, as well as a follow-up discussion on runtime libraries and link format incompatibilities.

It was also my intention to discuss two C compilers that support the large memory concept on the 8086. But since I have yet to receive one of the compilers, this comparative review will have to be included in a future column.

Several more letters have come in concerning C layout standards. While we don't have room to include their comments at present, they will undoubtedly lead to several columns on related topics of interest.

Comments about Unix

In the December column, I mentioned several aspects of Unix that I believe constitute weaknesses. Tony LiCausi of Canoga Park, California, sent me a letter with his dissenting opinions. He begins by remarking:

"Unix was designed for the programming environment, for use by folk already familiar with computers. Unnecessary I/O slows down a program's performance and hence should be eliminated. Thankfully, verbosity is often eliminated in Unix as the default option. However, many of the programs do have command line switches to increase the progress reporting of the program."

Unix was designed to circumvent the flaws of existing operating systems. Under Unix, modular programs are coupled via pipelines to create a convenient working environment and to reduce repetition of effort. Furthermore, Unix offers unequalled portability between large and small computer systems. Since it is successful in these respects, Unix's popularity is growing; hence, more and more computer users are gaining access to machines that use Unix. Many of these users are familiar with computers and programming, while others are novices to both computing and Unix.

Let's assume that a user is well versed in programming but is new to Unix; according to Mr. LiCausi, this user belongs to the group of people for whom Unix was designed. I'll assume that the user has an application in mind, since the average user does not write programs as ends unto themselves. To succeed in implementing an application, the user must learn about various aspects of Unix, including various shell commands, system calls, and runtime libraries.

I submit that this type of user will be slowed down significantly by the overall poor quality of the Unix documentation. Since examples are typically both terse and complicated, the typical user could spend an enormous amount of time circumventing a single snag. Imagine linking the mathematical function library as part of a C compilation. Two seemingly possible ways for doing this are:

```
cc –o test test.c –1m
            or
cc –1m –o test test.c
```

It turns out that the former method is correct, while the latter produces undefined function error messages at the linker stage. However, the documentation does not indicate that the "–1m" must be near the end of the line.

Our hypothetical user will also be slowed down by the cryptic nature of the error messages. True, the more sophisticated users require little or no prompting, and only gentle nudges when they make an incorrect entry, but the vast majority of users probably don't belong to this category.

Finally, the user will be slowed down by inexplicable and undocumented bugs in systems such as **troff** and **eqn**, which I mentioned before. Furthermore, circumventing bugs requires highly specific knowledge, and the solutions, even if available, may not be widely known. On the Unix machine I use, only a small group of people is aware of the procedure for avoiding some common but annoying bugs in **eqn**.

Mr. LiCausi makes a very valid point concerning this:

"Software bugs are not acceptable in any form at any time. Unix from Bell Labs is sold as unsupported software. Their attitude is, 'OK, we'll sell it to you, but we don't want to hear from you.' There is no mechanism to report or correct bugs at the source. Source code is not typically available to the average user. Source code is guarded on Unix like gold, just like any other commercial software. Having done debugging of the f77 compiler on one system, I found access to other software (for examples and integration) severely limited."

Thus, while Mr. LiCausi contends

that Unix was designed for people with programming experience, he sets his standard for "unnecessary I/O" and "program performance" based on the most sophisticated level of Unix user. Furthermore, it is still not clear to me why the sophisticated users should be more willing to stand for poor quality documentation and user-unfriendly software than everyday users. Finally, it is not at all apparent to me from practical experience that "many" of the standard shell commands support extended progress reporting.

Mr. LiCausi goes on to state:

"The secret society of Unix is partially the result of the paranoia of 'unauthorized access.' The remainder is due to the lack of commercial incentive. The incentive has obviously increased, as can be evidenced by the proliferation of books and articles on Unix and C."

I stated in my previous column that some users perceive Unix as a secret society because they cannot grasp certain aspects of its operation based on the written and on-line documentation alone. Only through initiated members can the knowledge (and undocumented tricks, patches, etc.) be obtained. However, other popular operating systems do not have the same aura as Unix. For example, CP/M users are not considered a "secret society" even though CP/M has nontrivial aspects. Nor do minicomputer operating systems such as VAX/VMS (DEC) and AOS/VS (Data General) have the mysterious quality of Unix. That's because they have much more documentation to explain the nitty-gritty of using the system.

Mr. LiCausi finishes by stating:

"Unix is not user-unfriendly. Admittedly it is beginner-hostile. Changes will be required to adopt Unix en masse (i.e., for the micro), but I hope that it does not spell the end of 'silent' software."

My final counter-remark is that most users never become sufficiently sophisticated in enough diverse aspects of Unix to escape the feeling of user-hostility. I personally still run into difficulties, even though I've used Unix for almost four years. As soon as I use aspects of Unix that are not part of my daily routine, "gottchas" appear.

Other Points of View

H. T. Gordon of Berkeley, California, sent in a letter concerning Unix. He writes:

"It was a pleasure to read the Skjellum critique of [Unix] in *DDJ* No. 86. Since C/Unix is becoming the sacred cow of the computer Brahmins, few have the courage to point out its flaws. It was designed by one of the elite, for elite users. To them it's concise and elegant, even 'obvious.' Only rarely is any explanation needed. If it is, a terse one will do. Prophets inspired by the high gods cannot waste their time spelling things out for mere mortals in words of one syllable."

I don't think that the state of C/Unix is as bad as Mr. Gordon suggests. While there are problems, I think that Unix has many strengths and these strengths will be the basis for future operating systems. It is the purpose of this column to examine strengths and weaknesses of C and Unix, and to propose solutions as well as to point out the problems. I encourage readers to pose solutions to these and other problems they encounter.

Low Level Input-Output in C

Patrick Cawood of Los Angeles writes:

"I read with great interest your last article on tendencies in Unix to produce poor operator interaction programs. I seem to have met some of the same in C.

"In order to provide a secure operator interface, one does not echo a keyboard character to the screen until it has been examined and approved by the program. But **getch( )** function automatically echoes — even line feeds, up arrows, etc.! Function **putch( )** provides an automatic line feed after printing on the screen!"

He goes on to state:

"I simply cannot believe that anyone would wittingly design these functions as they are, or not provide any alternative hardware interfaces. Especially considering some of the tasks I've heard were written in C. But perhaps these people writing serious software were all forced to write their own hardware interfaces."

The problem that Mr. Cawood is referring to exists in a number of non-Unix C compilers. To begin with, let's review the problem in the Unix environment. Under Unix, **putc( )** and **getc( )** acquire and return a single character, respectively. To offload the host system, however, many terminal interface boards programmatically handle the user input-output in lines. Thus, before any input is received by the host, a whole line must be entered. Naturally, the characters are echoed by the terminal interface hardware/firmware, and only limited line editing is permitted. On output, a whole line is buffered up before transmission to the terminal.

This process can be overcome by use of the "raw" terminal mode (raw implies no host character processing). In this mode, the program is completely responsible for input-output. This mode is much more expensive in terms of input-output cost, since the host must handle an inter-

rupt for each input character and perform an output request for each printed character. However, as I mentioned in the previous column, this is the only way for a program to get full control of what is entered and appears on the display device.

With the introduction of C to microcomputers, compiler implementors often based the behavior of the runtime libraries on their Unix experiences rather than on *The C Programming Language*. Whole lines are prebuffered by typical C runtime libraries before a single character is received by getc(); conversely, a whole line of output is internally buffered before it is printed on the display. Thus, the runtime libraries of microcomputer C compilers often emulate the terminal hardware found on real Unix systems.

However, this is not what putc() and getc() are supposed to do. In the truest sense, raw mode is the fundamental terminal mode. These functions should really work on a character-by-character basis. C libraries should permit selection between the "raw" and "cooked" modes (and echo and no-echo modes), thus permitting input-output flexibility without resorting to assembly language routines.

One C compiler that correctly handles low level input-output is the Q/C compiler from The Code Works. Single characters, for example, are required only when a getchar() call is made. While it doesn't seem possible to turn off the echo, this can be effected under Q/C since The Code Works is kind enough to provide runtime and compiler source code.

Link Formats

In the October 1983 column I discussed incompatibility between link formats. Readers had quite a bit to say about this and their comments are summarized here.

Guy Scharf of Mountain View, California, writes:

"I find the incompatibility of linkage editor formats to be a real problem. For example, I want to use Digital Research's Access Manager and Display Manager 86

with C. I would prefer CI-86 (because I am used to it) but have to switch to DRI's C because of the link format. Another compiler to learn and idiosyncrasies to surmount."

He concludes: "I'm not sure what to do about this problem (except complain)."

David D. Clark of State College, Pennsylvania, writes:

"The big problem isn't really the format of linkable files. Even compilers that use Microsoft's M80 and L80 will not allow linkage to code produced by different compilers. The function-calling protocols vary tremendously from compiler to compiler. BDS C and Q/C have fairly straightforward function-calling protocols. Eco-C, on the other hand, has a tortuous function-calling sequence. And even though Q/C and Eco-C use the Microsoft assembler and linker, the code files produced by them are not compatible."

The problem that Mr. Clark mentions is also present in the 8086 (MS-DOS, CP/M-86) world. Code from different compilers cannot be mixed because:

1. Each requires its own main function.
2. A wide variety of link formats exist.
3. Calling conventions differ between compilers.

The first two points are essentially insurmountable problems from the end user's point of view. However, the third point can be overcome by adding dummy routines to convert calling conventions.

Mr. Clark makes some additional points concerning the deficiencies of the 8080 Microsoft .REL (relocatable) format. These are of interest since many 8080 C compilers rely on this format. His comments are summarized in Table I (below).

It is obvious that an enhanced standard is necessary for the CP/M-80 world. Even under MS-DOS, where Microsoft has enhanced the linkage editor format, problems still exist.

Microsoft's MS-DOS linkage format supports long symbols (31 unique characters) and has case sensitivity. Thus, it

overcomes the objections that Mr. Clark posed for the 8080 .REL system. However, one nagging problem still remains: the dichotomy between object modules and libraries. Because of this dichotomy, the MS-DOS linker always includes the full contents of an .OBJ module during linking.

However, libraries are searched. To make a library efficient, each function must be compiled into its own .OBJ file. Each .OBJ file then becomes a single subblock in the library, and all the functions of that sub-block file are included at program linkage. If libraries and object modules were equivalent, this problem would be overcome since the functions of an object module would be separable during linking.

Runtime Libraries

Nonstandard runtime libraries are a plague to programmers. They inhibit portability and introduce bugs when software is transported between different compilers. Charles Brady of New South Wales, Australia, writes:

"An enormous contribution to standardization of C programs would be the publication of a standard I/O library for BDS C, with, if necessary, a modified runtime package. The very fast and efficient compiler, particularly with the symbolic debugger tool, provides a very inviting environment for software development. It is a great pity that this means abundant nonstandard C. As there is no inherent reason why this should be the case . . . someone should be able to . . . produce a Unix compatible I/O library."

This point is especially well taken in view of the large amount of BDS C software available through the C User's Group of Yates Center, Kansas.

In the case of the runtime library itself, a clear standard exists. This standard is spelled out in *The C Programming Language*. All compilers have to do is support a proper subset appropriate to the environment in which they work.

Conclusion

In this column, I have included the follow-up discussion on Unix and on link formats/runtime libraries. The discussion has been mainly of a critical nature. However, this is not intended to be a condemnation of the systems but rather an impetus for readers to suggest new ideas to improve what we already have. Unix is a worthwhile standard and should be supported. C is a valuable tool but can be improved. This column is a forum where we can discuss, develop, and nurture new ideas about C and Unix.   ∎DDJ

Table 1.

Deficiencies in the Microsoft 8080 .REL Format and Related Software

1. M80 and DRI's RMAC assemblers support only six unique characters (all upper case). This is awkward for many purposes.

2. While the .REL format apparently will handle seven unique characters, neither M80 nor RMAC supports this.

3. Apparently the .REL format, M80, and L80 were designed to work with the Fortran-80 compiler, which permits only six character symbols. This is an old standard and does not reflect the needs of today's compilers. The absence of case sensitivity in symbols is especially limiting.

OF INTEREST

by Michael Wiesenberg

TEXnical Information

If Stanford's Donald Knuth has his way, typesetting will no longer be an esoteric trade practiced only by a privileged and knowledgeable few. Computerized typesetting with its almost infinite flexibility permits individuals to set their own copy. No longer need the wasteful method be followed of one person producing copy on a word processor and printing that copy out on a daisy-wheel printer for another individual to reenter on the front end of a phototypesetter. That, of course, introduced new mistakes by the person doing the typesetting, an individual often untrained in the discipline addressed by the original author and thus unlikely to recognize many typographical errors.

Also the typesetting equipment, even if computerized, does not have error-checking facilities like the spelling and grammar checkers available on microcomputers. Sometimes the only interface between the supposedly computerized phototypesetting front end and the phototypesetter itself is not even a direct link but an antiquated paper tape. A completely new tape has to be punched each time a particular document is updated.

Some people predicted the death of the typesetting industry, or at least danger to members of the typesetting unions, when Knuth released **TEX** (pronounced *tek*) in 1978; but in fact what will happen is that typesetting will become more universally accessible, creating more jobs. When word processing was introduced, it did not — as some doomsayers had predicted — render secretaries obsolete; it just moved some of them from the typewriter to the terminal. What *did* happen was that more documents were produced. With the proliferation of computerized typesetting, we'll just see more good-looking documents.

Now, how can you get TEX onto your TRS-80 or Apple? Well, that's the bad news. You can't. Until recently, TEX was available only for mainframes. But here's the good news. You can get it on a system that has, depending on configuration, a minimum of 500K of main memory. Such systems range from the IBM PC to Hewlett-Packard computers. Since the

TEX software is in the public domain, you could theoretically put it onto a virtual memory system with as little as 128K memory, but you'd have to adapt the software yourself and write the drivers for whatever output device you needed. That shouldn't be too hard for you Pascal wizards, right?

If you plan to implement TEX yourself and want the public domain software, you can get it from Maria Code by special arrangement with the Computer Science Department of Stanford University. 1200-foot tapes in various formats are $82 per tape if you supply the tape, or add $10 if you

don't. The generic tape contains TEX82, WEB, fonts, and other miscellany; you need a Pascal compiler. **Reader Service No. 101.**

The TEXbook

Knuth's first version, TEX78, was written in SAIL. TEX82 is written in WEB, a system of structured documentation that combines the features of both a programming language (Pascal) and a document formatter (TEX). To illustrate Dr. Knuth's marvelous sense of humor (also found, by the

```
If the generator is now tuned to $f_0$, where the receiver power gain
is $G(f_0)$, the total power output will be
$$P=FkT\int^\infty_0 G(f)\, df+SG(f_0)$$
where $S$ is the available signal power. If $S$ is now adjusted so that
$P=2N$, i.e., so that the reading of the square law detector doubles, then
$$FkT\int^\infty_0 G(f)\, df=SG(f_0)$$
$$F={S\over kT{1\over G(f_0)}\int^\infty_0 G(f)\, df}={S\over kTB}
   \eqno{(3)}$$
where
$$B={1\over G(f_0)}\int^\infty_0 G(f)\, df=\hbox{effective noise bandwidth}$$
In order to find $B$ it is usually necessary to measure the power gain
as a function of frequency and integrate the curve graphically. This
method has major disadvantages, however, since it is time-consuming
and since $G(f)$ may vary with different conditions of adjustment.
A further disadvantage is the difficulty of determining $S$ with the required
accuracy at the low levels involved.
```

Figure 1.

Source code for the HPTEX output shown in Figure 2, below.

If the generator is now tuned to f_0, where the receiver power gain is $G(f_0)$, the total power output will be

$$P = FkT \int_0^\infty G(f)\, df + SG(f_0)$$

where S is the available signal power. If S is now adjusted so that $P = 2N$, i.e., so that the reading of the square law detector doubles, then

$$FkT \int_0^\infty G(f)\, df = SG(f_0)$$

$$F = \frac{S}{kT \frac{1}{G(f_0)} \int_0^\infty G(f)\, df} = \frac{S}{kTB} \tag{3}$$

where

$$B = \frac{1}{G(f_0)} \int_0^\infty G(f)\, df = \text{effective noise bandwidth}$$

In order to find B it is usually necessary to measure the power gain as a function of frequency and integrate the curve graphically. This method has major disadvantages, however, since it is time-consuming and since $G(f)$ may vary with different conditions of adjustment. A further disadvantage is the difficulty of determining S with the required accuracy at the low levels involved.

Figure 2.

Actual output of HP2688A Laser Printer using HPTEX.

Source and output courtesy of Boise Division of Hewlett-Packard.

way, in his *T_EXbook*): you pass a .WEB source file through the TANGLE program to produce a .PAS Pascal file or through the WEAVE program to produce a .TEX file (that can now be processed by T_EX). ("Oh what a tangled web we weave. . . .")

You cannot use T_EX without a copy of *The T_EXbook*. It is a delight to read, though really meant only for actual users of the software because it has a lot of exercises for you to try out. If you don't do them, you won't be able to follow what he's talking about. If you are familiar with his monumental series, *The Art of Computer Programming* (considered by most programmers to be the best thing written about programming), you have an idea of the level of expertise required for *The T_EXbook*. Published jointly by Addison-Wesley and the American Mathematical Society, the book may be ordered from the latter for $15 (including p. and h., or add $2 for first class, $3 for airmail outside the U.S.). For 438 pages, most of it in nine- and ten-point type, that's a bargain! **Reader Service No. 103.**

HPT_EX

The version I have been using is HPT_EX, Hewlett-Packard's implementation of T_EX82, on a 68000-based HP series 200 Pascal workstation, model 36. The software costs $4000. A typical configuration includes 9836A (the workstation with 500K RAM, Pascal, BASIC, UCSD editor, and several other utilities), $12,110, three more 256K memory boards at $1060 each, a disk interface for $605, DMA card for $500, 65Mb hard disk, $17,350, and the HP2688A desktop laser printer for $26,000. (You could save some money by substituting a model 26: you need 1.25Mb main memory, not necessarily the 1.75 I've been using; you can "get by" with 15Mb storage — although 65 is recommended — and you don't necessarily need the high-speed memory access, which also eliminates the need for the disk interface card.)

With this system, I can literally typeset an entire book by myself. I can choose from any of over 340 fonts, which come in sizes ranging from five to 40 points. A powerful macro set allows easy specification of any format. Dr. Knuth originally developed T_EX so that complicated mathematical formulas in his books would be set exactly as he wanted, rather than as the typesetter felt it should look. T_EX really shines at this, but it also excels at easily setting tables and other complicated formats.

HPT_EX comes with an extensive set of macros that permit, among other things, multicolumn mode, with each column using a different font and a different measurement between lines, automatic generation of table of contents, and automatic indexing. I can define macros myself that permit juggling figures and tables around in a manual without worrying about changing the numbers (what happens to Figure 3-1, for example, and all the figures thereafter, when you have to insert two figures at the front of a manual?) of each. Many other computerized typesetting systems do not allow complete page layout, but this one does. Figures automatically float to a page with enough room for them, footnotes fall on the page on which they are referenced, and if material is added, the rest of the text automatically reformats itself.

Those familiar with typesetting know that any time anything is changed in a book, the effect ripples through the rest of the book, often requiring intervention by the typesetter on each page. Because things move around, technical writers have learned never to refer to other sections by page number; with T_EX, this is no longer the case.

The heart of the HPT_EX system is the HP2688A laser printer, which is no bigger than a computer terminal. The resolution produced by this little marvel is 300 dots per inch. While that does not approach the 3000 dots per inch of the Autologic APS-5, it is considerably better than most dot matrix or daisy-wheel printers. And quieter. And faster. **Reader Service No. 105.**

And For The 1000

Third-party vendors provide T_EX for other HP computers, including the 3000 and the 1000. On the 1000 minis and micros with RTE-6/VM or RTE-A operating systems, J_DJ Wordware offers **T_EX/1000**. You need at least 512K main memory and 2Mb disk storage. Output devices include the Toshiba P1350 and Epson MX80, with HP2688A soon to be supported. T_EX/1000 formats pages at 5 to 30 seconds each, depending on processor and memory. The output driver builds a page in memory, and then dumps the graphics to a printer. $2500. **Reader Service No. 107.**

Tick Tock

TYX Corporation offers an implementation of T_EX, written in C rather than the usual Pascal (or SAIL), on

DEC Rainbow, Victor 9000, IBM PC, and their own integrated **T_YXSet 1000** system. The latter consists of a minicomputer that supports up to 24 users, a Canon Laser printer, a Mergenthaler Omnitech laser typesetter, terminals, and a typesetting interface. In fact, this system could provide the entire front end for a typesetting system suitable for a magazine or the publications department of a medium-sized company. No longer would a magazine have to wait several days to look at galleys and several more for the corrections.

TYX also offers smaller versions of their system configured on DEC Rainbow, Victor 9000, or IBM PC. All come with software and a friendlier interface than normally comes with T_EX. Various function keys insert commands as needed, including the right number of beginning and ending group delimiters. All systems include a screen preview mode. While this implementation is T_EX78, TYX expects to have T_EX82 within a year. TYX also provides the service of taking T_EX output files from customers who do not have phototypesetters and providing camera-ready copy from their own Autologic. **Reader Service No. 109.**

Get on the TUG

If you plan to use T_EX, or attempt to implement it on your system, you will want to talk to others who have already done it. T_EX has been put onto everything from Z80 systems to IBM 370s. The **T_EX Users' Group** puts people in touch with "site coordinators," those in your area who have already implemented T_EX with a system like yours. They also publish *TUGboat*, have yearly meetings, sell T_EX tee shirts, etc. They'll send you an information packet containing *TUGboat* 1 #1, "T_EX and Metafont: Errata and Changes," the T_EX Users' Group Membership List (with the over 800 members sorted by name, computer, and output device), and forms for ordering T_EX82 and joining TUG. You can also get information on videotaped instruction and special in-house courses on T_EX. **Reader Service No. 111.**

Lost of Storage

Optical disk storage may be available within a few years for most micros, and at the prices now being paid for hard disks. The advantage is storage not in megabytes, but in *gigabytes*! Companies like Optimem, Thomson CSF, Panasonic, and 3M are

working in this field. Who needs that kind of storage, you ask? Hospitals, storing thousands of digitized x-rays for immediate display, each of which can require 512 by 512 pixels by 16 bits deep, CAD systems, and movie production companies, who need to store millions of frames with virtually instantaneous recall. Did you know that credit cards are *now* available that can hold megabytes of information? And here's an ominous development: the military is very interested in optical drives because they feel such drives to be the only medium that could survive a nuclear attack. You can read all about all of this in *Optical Memory News*, $295 per year for six issues. **Reader Service No. 117.**

Improving the Compatibles

The **STM Personal Computer**, from STM Electronics (the folks who brought you the Pied Piper portable), a 17-pound, IBM PC-compatible, transportable computer based on the 80186 true 16-bit CPU, with integrated 16-line, 84-character, LCD, backlit by an electroluminescent panel for total darkness use, 40-column thermal printer, 94-key (including numeric pad and 10 function keys) keyboard, programmable 300- to 1200-baud auto dial/auto answer direct connect data communications modem, 256K RAM (expandable to 512), dual 1Mb (800K each formatted) floppies, standard video monitor support with windowing software and high resolution graphics, parallel port, two programmable serial ports, hard disk interface, edge connector to send I/O to an external expansion box, and integrated software, running under MS-DOS 2.0, with word processing, spreadsheet, database, graphics functions and communications software all standard, costs $3000, with a desktop version for $2500. **Reader Service No. 115.**

Cross Assembling 68K

Quelo wants us to know that their cross assembler for the MC68000 runs not only on CP/M-80, but also on CP/M-86, CP/M-68K, and PC-DOS, and they offer portable C source code. You get up to 31 characters in symbols, lowercase distinction, byte object linking, XOR, MOD, and NOT, *break* and *elseif* for control structures, error messages in English, a linker with limitless program size and number of modules that can be linked, and linker resolution of complex expressions. **Reader Service No. 119.**

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DDJ

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